University of Arkansas, Fayetteville ScholarWorks@UARK

Graduate Theses and Dissertations

5-2020

# Cybersecurity Methods for Grid-Connected Power Electronics

Stephen Joe Moquin University of Arkansas, Fayetteville

Follow this and additional works at: https://scholarworks.uark.edu/etd

Part of the Electrical and Electronics Commons, Power and Energy Commons, and the Systems Architecture Commons

#### Citation

Moquin, S. J. (2020). Cybersecurity Methods for Grid-Connected Power Electronics. *Graduate Theses and Dissertations* Retrieved from https://scholarworks.uark.edu/etd/3680

This Thesis is brought to you for free and open access by ScholarWorks@UARK. It has been accepted for inclusion in Graduate Theses and Dissertations by an authorized administrator of ScholarWorks@UARK. For more information, please contact scholar@uark.edu, uarepos@uark.edu.

Cybersecurity Methods for Grid-Connected Power Electronics

A thesis submitted in partial fulfillment of the requirements for the degree of Master of Science in Electrical Engineering

by

Stephen Joe Moquin Auburn University Bachelor of Science in Philosophy, 2009 University of Arkansas Bachelor of Science in Electrical Engineering, 2017

> May 2020 University of Arkansas

This thesis is approved for recommendation to the Graduate Council.

H. Alan Mantooth, Ph.D. Thesis Director

Jia Di, Ph.D. Committee Member Roy McCann, Ph.D. Committee Member

Yue Zhao, Ph.D. Committee Member Chris Farnell Committee Member

#### Abstract

The present work shows a secure-by-design process, defense-in-depth method, and security techniques for a secure distributed energy resource. The distributed energy resource is a cybersecure, solar inverter and battery energy storage system prototype, collectively called the Cybersecure Power Router. Consideration is given to the use of the Smart Green Power Node for a foundation of the present work. Metrics for controller security are investigated to evaluate firmware security techniques. The prototype's ability to mitigate, respond to, and recover from firmware integrity degradation is examined. The prototype shows many working security techniques within the context of a grid-connected, distributed energy resource. Further work is expected in the Cybersecure Power Router project. Consideration is also provided for the migration of the present research and the Smart Green Power Node to realize a pre-production prototype.

©2020 by Stephen Joe Moquin All Rights Reserved

#### Acknowledgements

I would like to thank my advisor, Dr. Alan Mantooth, for taking a chance on me, affording many lessons on leadership, and the opportunity to participate in the environment he has so diligently cultivated over the years. It is an honor to work in the Power Mixed-Signal Computer Aided Design research group. Dr. Jia Di, Dr. Roy McCann, and Dr. Yue Zhao are owed much thanks for their advising and contributions to the success of my academic career. I owe a great deal to Chris Farnell, and have benefitted greatly from the years of mentoring and friendship he has provided me. The present work would not be possible without the edifice of work Chris Farnell has constructed, and the foundations laid by Dr. Mantooth and the fellow research faculty of the University of Arkansas. I am thankful for my teammates SangYun Kim and Nicholas Blair. Dr. Shannon Davis, Beth Benham, Karin Alvarado, Jamie Stafford, Cindy Pickney, Sharon Brasko, Connie Howard, and Tracey Long are owed thanks for their care and administration. Eric den Boer, Jeff Knox, and Mike Steger are owed thanks for their practical wisdom in life and in research. I am thankful for the support of my past and present fellow graduate students: Haider Mhiesan, Hamdi Albunashee, Haoyan Liu, Yuzhi Zhang, Janviere Umuhoza, Shuang Zhao, Yusi Liu, Sayan Seal, Andrea Wallace, and Audrey Dearien. Finally, I'd like to thank the Department of Energy and industrial partners who contributed to the Secure Evolvable Energy Delivery Systems research center for the funding of both the work and researchers of the Cybersecure Power Router.

Chapter 1 - Introduction	1
Chapter 2 - Technical Background	5
Chapter 3 - Cybersecure Power Router	
3.1 System Description	
3.2 Power Electronics	
3.3 Digital Signal Processor Board	
3.4 Complex Programmable Logic Device Board	
3.5 Signal Splitter	
3.6 Hardware Authentication Module	
3.7 BeagleBone Black	
3.8 Test Bed	
3.9 Power Flow	
3.10 Data Flow	
3.11 Control Multiplexing	
3.12 Firmware and Boot Management	
3.13 Hardware Authentication	
3.14 Submodule Encrypted Communication	
3.15 Hardware Protections	
3.16 Display	
Chapter 4 - Results	
Chapter 5 - Future Work	
5.1 Multi-Mission Controls	
5.2 SGPN and CSPR Integration and Completion	
Chapter 6 - Conclusion	
References	
Appendix	61
Appendix A: Hardware and Software Design Details	61
Appendix B: EEPROM_WRITE_PASSWORD	
Appendix C: CSPR_V7.lpf	
Appendix D: hardware_protections.vhd	
Appendix E: CSPR_MODULES.vhdl	
Appendix F: top.vhdl	

## **Table of Contents**

# List of Figures

Fig. 1. Design Inventory for Distributed Energy Resource
Fig. 2: Cost/Performance regions (left) for various security solutions (right)
Fig. 3. Cybersecure Power Router prototype
Fig. 4. Block diagram (left) and figure (right) of Cybersecure Power Router prototype
Fig. 5. Minimal configuration of Cybersecure Power Router prototype
Fig. 6. Asynchronous buck converter schematic of Power Electronics Evaluation Board of the UCB project
Fig. 7. Asynchronous boost converter schematic of Power Electronics Evaluation Board of the UCB project
Fig. 8. 3-Phase inverter filter and current sensing schematic of Power Electronics Evaluation Board of the UCB project
Fig. 9. 3-Phase inverter/rectifier switching stage schematic of Power Electronics Evaluation Board of UCB project
Fig. 10. Power Electronics Evaluation Unified Controller Board, with Signal Splitter and Hardware Authentication Module, used in Cybersecure Power Router prototype
Fig. 11. USB to UART schematic of Digital Signal Processor board of the UCB project 20
Fig. 12. DIMM pinout of Digital Signal Processor board of the UCB project
Fig. 13. Majority of analog and digital I/O used by Digital Signal Processor board of UCB project
Fig. 14. IDCs in Complex Programmable Logic Device Unified Controller Board
Fig. 15. PCB layout of Complex Programmable Logic Device board of UCB project
Fig. 16. Fabricated Complex Programmable Logic Device board of UCB project used in Cybersecure Power Router prototype
Fig. 17. Complete schematic of Signal Splitter
Fig. 18. Complete schematic of Hardware Authentication Module
Fig. 19. Hardware Authentication Module used in Cybersecure Power Router prototype

Fig. 20. BeagleBone Black
Fig. 21. Block diagram of testbed used in Cybersecure Power Router project
Fig. 22. Simplified block diagram of grid-connected power flow capabilities of UCB hardware and Cybersecure Power Router
Fig. 23. Block diagram of clock generation within Complex Programmable Logic Device of the Cybersecure Power Router
Fig. 24. Block diagram of data bus controller within Complex Programmable Logic Device of the Cybersecure Power Router
Fig. 25 .Block diagram of serial interface within Complex Programmable Logic Device of the Cybersecure Power Router
Fig. 26. Firmware code snippet to generate heartbeat from Digital Signal Processors
Fig. 27. Heartbeats of Controllers 1 and 2 while running identical firmware
Fig. 28. Block diagram of control multiplexing using Digital Signal Processor signals and Hardware Authentication Module within the Complex Programmable Logic Device of the Cybersecure Power Router
Fig. 29. Block diagram of hot patching process
Fig. 30. Diagram of CSPR components involved in hardware authentication
Fig. 31. Excerpt from EEPROM datasheet, with erroneous information noted
Fig. 32: Hardware Authentication Module communication beside oscilloscope capture
Fig. 33. Block diagram of encrypted serial communication within Complex Programmable Logic Device used in the Cybersecure Power Router
Fig. 34. Shoot-Through hardware protection
Fig. 35. LED display diagram for Cybersecure Power Router
Fig. 36. Time between execution cycles of controller firmware vs. switching frequency
Fig. 37. Detail of Figure 36
Fig. 38. Inverter output voltage at 114, 118, and 128 kHz switching frequencies

Fig. 39. Inverter output during controller transition	. 48
Fig. 40: Radar chart of missions for controls	. 51
Fig. 41: CSPR and SGPN migration	. 52
Fig. 42. Digital Signal Processor Unified Controller Board schematic	. 61
Fig. 43. Power Electronics Evaluation Unified Controller Board schematic	. 62
Fig. 44. Complex Programmable Logic Device Unified Controller Board schematic	. 63
Fig. 45. BeagleBone Black schematic	. 64
Fig. 46. Hardware Authentication Module schematic	. 64
Fig. 47. Analog Splitter schematic	. 65
Fig. 48. Crontab configuration on BeagleBone Black to run CPLD UCB LED control script or startup	
Fig. 49. Content of LED control Python script running on the BeagleBone Black	. 66

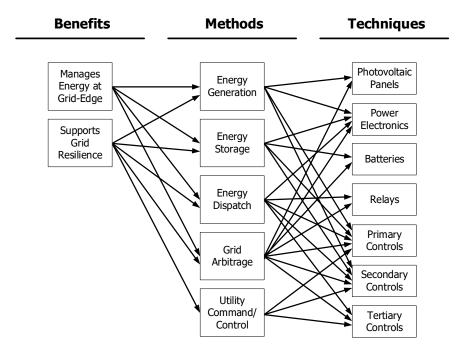
# List of Tables

TABLE I: Security-by-Design cycles	5
TABLE II: Defense-in-Depth Dependencies	6
TABLE III: Cyber-Physical Attack Matrix	. 15
TABLE IV: MachXO2 Family Features	. 24
TABLE V: Available MCLK frequencies	. 24

#### **Chapter 1 - Introduction**

The reliability and safety of the electrical grid face challenges. These challenges include aging infrastructure, tight regulatory environments, and the integration of new technologies. Power electronics are some of these new technologies, and provide a wide range of assets and liabilities to the electrical grid, its operation, and its evolution. Reactive power compensation, phase load balancing, battery energy storage systems, solar power, and flexible ac transmission are all potential assets. These devices may pose serious threats to both the electrical grid and interdependent critical infrastructure [1]. As these grid-connected power electronics permeate more of the electrical grid, the need for their security becomes greater [2].

A series of questions can begin this investigation of security for grid-connected power electronics, and establish appropriate security measures [3]. The first question is "what benefit does the device provide?" In the present case, the distributed energy device manages energy at the edge of the electrical grid, and supports grid resilience.



#### Fig. 1. Design Inventory for Distributed Energy Resource

The next question is "how does the device provide those benefits?" A distinction in method and technique is useful, here. The methods to realize energy management at the grid-edge are energy generation, storage, dispatch, and arbitrage (see "Methods" in Figure 1). These methods identify a particular process or type of system. The techniques to realize energy management at the grid-edge include the use of photovoltaic panels, dc converters, inverters, batteries, and hierarchical controls. These techniques identify the practical elements of a method, that is, of a process or system.

After arriving at a set of techniques, the next question is "what can happen to prevent this technique from working?" A boost converter or more complicated topology can be used to provide maximum power point tracking, current control, and hardware protection for a string of photovoltaic panels. In short, answering the question is difficult [4]. The hardware can fail, the switches can cause excessive electromagnetic interference, and noisy environments can corrupt transmitted data. Poor design can also cause various failures. The list of possible points and modes of failure is extensive for a complex system like a converter. Putting aside the incomplete answer, the next question is "what can be done to protect this technique?" In the case of a dc converter, more robust hardware can be used, filters can reduce electromagnetic interference, and error detection in communication can limit the use of corrupted data. This is not an exhaustive list of points and modes of failures or security measures. To better organize and address these questions, product lifecycle management and design dependency can be used (and will be discussed in greater depth in the next section). Product lifecycle management can be used to consider security of a device from specification all the way to end-of-life, and illustrates the security-by-design process. Design dependency can be used to build layers of security to protect an asset or the system as a whole, and illustrates the defense-in-depth method. A cyber-physical threat matrix is built on the

design dependencies listed and the implementation of security techniques to address the potential threats.

The Cybersecure Power Router (CSPR) uses the security-by-design process and defense-indepth method to realize a cybersecure distributed energy resource. The CSPR operates as a solar inverter and battery energy storage system. It has hierarchical controls and network communication, allowing a utility operator to control the device, its operation, and its power flow. Finally, it employs a number of security features for a wide range of functionalities. These security features include AES-128 encryption for network communication, hardware-assisted monitoring for improved firmware integrity during runtime, hardware protection during nominal operation, and more.

The Smart Green Power Node (SGPN) is a device developed at the University of Arkansas to manage energy resources at the grid edge, specifically in residential applications. The SGPN predicts and optimizes power flow of a solar inverter and battery energy storage system. Hierarchical controls and network communication are also included within the design of the SGPN. The system optimizes power flow of the energy resources through powerful predictive algorithms and weather data collection. The system is rated for 2 kW operation.

The Unified Controller Board (UCB) project is a set of hardware, firmware, software, and instructional material also developed at the University of Arkansas. The devices within the UCB project include a DSP docking station; a complex programmable logic device (CPLD) PCB; buck and boost converter and inverter PCB; and several expansion boards. The Unified Controller Boards are designed around flexible controls and modular hardware.

A real-world system of sophisticated, grid-tied power electronics is needed to show the practical demands and limits of security. The use of the Smart Green Power Node and the Unified

3

Controller Board devices as a prototype for the Cybersecure Power Router was used for such a real-world system. The SGPN is a distributed energy resource with many sophisticated assets. The 2 kW power rating of the system provides an appreciable power flow for grid-connected applications. The UCB devices are modular, allowing for rapid configurability of hardware and controls. The UCB software and firmware is extensible, allowing for the integration of security features and changes in control hierarchy. The combination of these two systems provides the necessary power flows, complexity, and direct results necessary for this cyber-security investigation.

The present work shows the security-by-design process and defense-in-depth method for a grid-connected, distributed energy resource prototype. The security features chosen are developed and tested within a grid-connected power electronics context. Security features to protect firmware integrity at runtime are specifically investigated. The ability for the CSPR prototype to quantify firmware integrity degradation and respond to firmware integrity failure is shown. This ability is provided by the CSPR prototype monitoring and maintaining liveness of controllers through control multiplexing. Future research into greater flexibility and resiliency of controls is discussed. Finally, the necessary work to migrate research from the Smart Green Power Node and the Cybersecure Power Router into a pre-production prototype is presented.

### **Chapter 2 - Technical Background**

Security-by-design ensures greater security of a device by considering both the processes

behind the development and life of a device, and the device itself. Product lifecycle management

[5] serves as the framework for security-by design. This process stands over and above the IEEE

standard for system, software, and hardware verification and validation [6].

Lifecycle Stage	Security Feature(s)			
Hardware				
Specification	IEEE Standards (e.g., 1547)			
Simulation	Accurate Modeling, Thermal Co-Simulation			
Design	IEEE Standards (e.g., 3001, 3003), Thermal Co-Design			
Verification	Electrical Rule Checking, Design Review, Hardware-In-the-Loop			
Firmware				
Development	Restricted Access, Version Control, Standard Protocols, Standard Libraries			
Distribution	Restricted Access, Message Digest, Server Authentication			
Installation	Message Digest, Error Detection and Correction			
Run-Time	Side Channel Analysis, Challenge-Response Authentication			
Manufacture				
Fabrication	Trusted Supplier, ISO 9001 Certification, Hardware Authentication			
Quality Control	Burn-In Testing, Fuzz Testing, Standard Metrics			
<b>Design</b> Iteration	Restricted Source Code and Design Files			
Operation				
Installation	Certified Installers, Standard Connections, Lockout-Tagout			
Use	Key Management, Challenge-Response Authentication, Behavior Analysis			
Aging	Hardware Health Diagnostics			
Attack	User Authentication, Command Whitelisting, Asset Segmentation			
Failure	Fails Safe, Hardware Protection, Resilient Communication			
Recovery	Startup Sequence, Sanity Check, Firmware Integrity Check			
Maintenance				
Update	Patching, Maintained Uptime			
Replacement	Modular Design			
Upgrade	Modular Design, Reconfigurable Architecture, Flexible Controls			
End of Life				
Removal	Lockout-Tagout			
Documentation	Failure Modes, Effects, and Diagnostic Analysis (FMEDA)			
Reiteration	Restricted Access to Source Code, Specification and Quality Control			
Disposal	Certified eWaste Recycling and Disposal			

 TABLE I: Security-by-Design cycles

Any device, a distributed energy resource in this case, has a lifecycle. It is specified, designed, fabricated, tested, installed, operated, uninstalled, and disposed of during that lifecycle. Each step in the lifecycle of a device serves some purpose. For instance, hardware specification creates the exhaustive list of design requirements for a device. The result of the hardware specification step is a list. How could security be applied to this step? The use of standards (in this case, IEEE 1547 for the design of utility electric power systems and distributed energy resources) provides greater assurance that the list created in the hardware specification step is exhaustive. Stated another way, the IEEE standards secures the intended result of hardware specification. The product lifecycle security approach also requires a designer to consider the full lifecycle of a device, not just the useful life. In the context of the electric grid, people install and remove distributed energy resources. By considering the lifecycle of a device, the safe installation and removal of a device is considered and included in the hardware and firmware design stages. For instance, a Lockout/Tagout technique can be designed for a solar inverter to keep both people and hardware safe during installation and removal [7].

Defense-in-depth provides layered security for assets of a device [8]. Returning to the "Design Inventory" from the introductory section, each technique in the design inventory process has dependencies. These dependencies arise from the techniques chosen to realize a device. An inverter depends on various switches, gate drivers, feedback signal chains, capacitors, other hardware components, and firmware to operate. Examples of these dependencies are listed in Table II, along with possible methods of security.

Design Category	Security Feature			
Component Health	Hardware Authentication, Hardware Protections, Safety Factor			
Feedback Signal Chain	Galvanic Isolation, Buffered I/O			
Temperature Control	Thermal Management, Current Limitation			
Current Control	Controller Current Limiting, Body Diodes, Fuses			

 TABLE II: Defense-in-Depth Dependencies

These dependencies are not spread across the lifecycle of a device, as illustrated in the securityby-design process. Rather, these dependencies are logical constituents of the design of a device, and are typically part of a complex, cyber-physical set of interdependencies [9]. For instance, how might one protect the current flow into the batteries from potentially damaging commands? Secure network communication [10] protects the battery energy storage system from noise and remote adversaries. If the secure network communication is defeated, current controls prevent harmful behavior of the device [11]. If the current controls are defeated, various design features (like galvanic isolation, buffered I/O, fuses, and over-design) allow the device to withstand or limit the harmful behavior [12]. In this case, the energy storage assets of the BESS are protected by layers of security.

Security features have a range of costs and performance gains [13]. Any change to the design of a system, including those to increase security, comes at a cost. This cost may include the price of more sophisticated integrated circuits, hardware to support increased power consumption, time to develop new firmware, or expertise to identify and execute security strategies. Improvements to security may come at a low cost. An existing system may extensively benefit from simple firmware management [14]. Such management, including revision, could greatly increase system security without incurring costs from additional hardware and hardware development. Beyond firmware management, an example of a firmware security feature is a checksum for network communication [15]. The inclusion of this security feature is lightweight: incurring a small increase in firmware size, computational load, and communication overhead. Checksums can prevent electromagnetic noise from corrupting communication, and weakly protect against a malicious actor tampering with communicated data. A checksum is an instance of a common security feature implemented in firmware, but is far from the full benefits of improved firmware management and security.

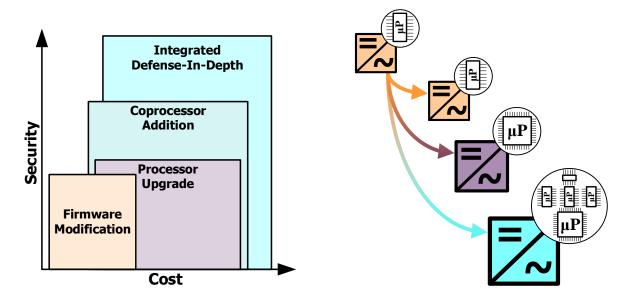


Fig. 2: Cost/Performance regions (left) for various security solutions (right)

Not all development leads to increased security. Consider the use of encrypted communication for a controller using a digital signal processor. The controller is responsible for the power flow through a solar inverter. Processing resources are required to encrypt and decrypt communication with the controller. The increased demands on the DSP during communication may cause overrun conditions [16] and degrade power processing. The system has more confidential communication, but at the cost of lower integrity of power processing. The general effect may be a less secure system, despite the addition of a security feature and the costs of development.

The cost/performance analysis for many security features used in the Cybersecure Power Router are readily available. The Advanced Encryption Standard 128-bit key (AES-128) is well defined [17],[18] and researched in various applications [19], [20], [21]. AES-128 is used in communication between hardware modules, and between outside controllers and the CSPR. The MD5 message-digest algorithm is also well researched [22], [23], [24]. The MD5 algorithm is used to protect the integrity of firmware as it passes from an outside controller, to the Cybersecure Power Router, and into any on-board controller. Hardware protections against overcurrent [25] and shoot-through conditions [26] are also well understood. Overcurrent and shoot-through conditions are used in both controller firmware and within the hardware logic of the Digital Signal Processors. These techniques are robust and well researched. Security techniques are still needed for other essential functions of any grid-connected power electronics.

Firmware security at runtime for power electronics is less researched and widely implemented [27]. Grid-connected power electronics typically use microprocessors, microcontrollers, or processors to run firmware. Microprocessors and microcontrollers often use watchdog timers to reset the device in the case firmware execution faults. If the controls of the power electronics run on the reset device, the power electronics will stop. The use of a watchdog timer may, therefore, be inappropriate for grid-connected power electronics, where downtime is to be minimized or eliminated [28]. Processors can provide more sophisticated techniques than microprocessors and microcontrollers to prevent, detect, identify, and recover from firmware execution faults, namely through using an embedded operating system [29], [30]. The choice of processors in grid-connected power electronics at runtime. Is there an option between a simple watchdog timer and a sophisticated embedded operating system? If so, how can the cost and performance of that runtime security be evaluated?

The two considered threats to CSPR firmware integrity during runtime are task overrun and firmware patching. Other threats are relevant [31], but fall outside the present scope of grid-connected power electronics. A task overrun condition occurs if a controller is not able to finish

the various tasks before another set of tasks are started [32]. The result is a degradation in the power flow of the power electronics, as shown later. Unlike task overrun, firmware patching is more likely to halt power flow altogether. The patching process requires the rebooting of the DSP running a controller, halting the controller during the process. The power flow through the electronics is, therefore, also halted as the DSP reboots.

Task overrun can be described as a controller's loss of liveness. Formally, liveness can be expressed as

$$\forall \alpha: \alpha \in S^*: (\exists \beta: \beta \in S^{\omega}: \alpha \beta \models P), \tag{1}$$

where  $\alpha$  and  $\beta$  are a sequence of states,  $S^*$  is a set of finite sequence states,  $S^{\omega}$  is a set of infinite sequence states, and P is a property (executability, in this case) [33]. Formally, this definition reads as there exists a state within an infinite set of states that satisfies a given property (executable), such that it does so given any arbitrarily sized sequence of compossible states. Or, simply, liveness means a task will be executed, even if there are many more tasks for a controller to complete. The tasks to be executed are part of interrupt service routines on the DSP, and in the present work are initiated every switching period. Each interrupt request (IRQ) initiating an interrupt service routine (ISR) is assigned a priority. The DSP resources handling an ISR are a critical section to other ISRs, especially those at the same priority. A critical section is a set of resources accessed or used by multiple processes [34]. For the present architecture, an IRQ can cause the interruption of an ISR in progress. This is also true if the ISR in progress and raised IRQ have the same priority. Here is an example from the Cybersecure Power Router. Assume the switching frequency is set at 30 kHz. Every 33.33 µs, a number of interrupt requests will be raised within the DSP. These IRQs signal the DSP to read various voltages and currents, perform mathematical operations, look up values, and set the pulse width modulation of several switches. Let's assume the DSP requires 80 µs to handle all of these tasks. While processing the last round of ISRs, new ISRs are created. The previous ISRs are interrupted and started again by the new IRQs. The tasks never complete, given their interruption and restart during processing. Theoretically, the result is the controller losing liveness.

A simple consideration of timing can maintain liveness. The controller has a maximum duration of time to complete its tasks:

$$T_{d,max} = \frac{1}{f_s},\tag{2}$$

where  $f_s$  is the switching frequency. In the example above, the 30 kHz switching frequency provides a maximum duration of 33.33 µs for the controller to complete its tasks. This is the theoretical maximum amount of time the controller can take to process the tasks of one switching period before being interrupted by the next switching period. The time the controller requires to complete the switching period tasks is not dependent on the switching frequency, however. The time required dependents more on the firmware, speed of the DSP, and competing interrupts (such as those from communication). This required time can be measured. This duration of time can vary, even if the firmware and processor remain the same. A mean time for the completion of tasks can be empirically found, and used to quantify available processing resources. The expected available resources can be articulated as

Available Resources (%) = 
$$\left(1 - \left(T_{d,mean} * f_s\right)\right) * 100\%,$$
 (3)

where  $T_{d,mean}$  is the mean of the duration of time the controller requires to complete switching period tasks, and  $f_s$  is the switching frequency. As this percentage approaches 0%, tasks are more likely to be interrupted, and liveness of the controller is more likely to be compromised.

The Cybersecure Power Router uses a signal sensitive to controller liveness and hardwareassisted monitoring to protect controller liveness. The next section details the design of the Cybersecure Power Router, especially the security features protecting controller liveness. Later sections provide the results and examination of the operation and performance of those security designs.

#### **Chapter 3 - Cybersecure Power Router**

#### 3.1 System Description

The Cybersecure Power Router is a set of switch-mode power electronics, controllers, and processors. The majority of controller devices used are from the Unified Controller Board project, developed by Chris Farnell at the University of Arkansas. These devices include the Power Electronics Evaluation Unified Controller Board (PE Eval UCB), Complex Programmable Logic Device Unified Controller Board (CPLD UCB), and the Digital Signal Processor Unified Controller Board (DSP UCB).

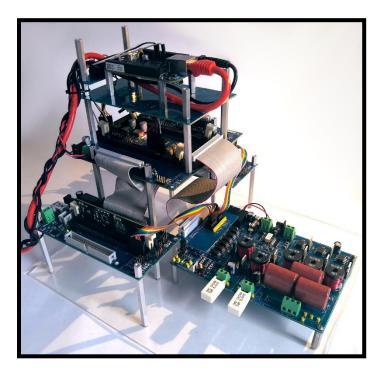
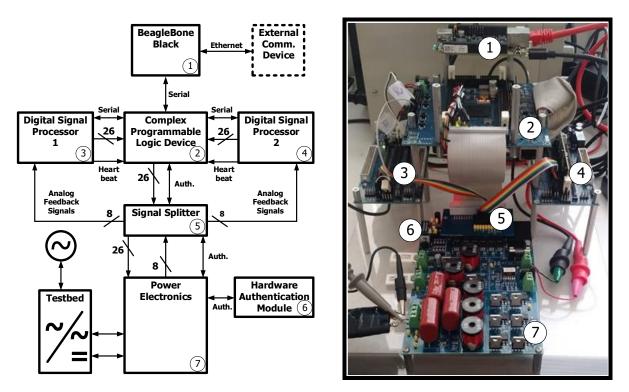


Fig. 3. Cybersecure Power Router prototype

The Digital Signal Processor Unified Controller Board uses many design features of the Texas Instruments' TMDSDOCK28335 Digital Signal Processor docking station. The power electronics of the PE Eval UCB are a buck converter, boost converter, and an inverter/rectifier. A testbed provides dc and ac power flow to provide safe and reliable conditions for testing. Two Digital Signal Processors are used in the prototype. Each DSP is capable of controlling the power electronics. A Complex Programmable Logic Device UCB provides many security features, and multiplexes the control signals of the DSPs. The Hardware Authentication Module is used to authenticate the power electronics and enable power flow. The Signal Splitter routes analog signals to both DSPs from the power electronics. Finally, a BeagleBone Black provides high-level control and Ethernet communication to the CPLD UCB.



*Fig. 4. Block diagram (left) and figure (right) of Cybersecure Power Router prototype* The CSPR prototype can be reconfigured. The modular design allows flexibility in both hardware and control. A more simplified configuration of the prototype could use one or two DSPs on the CPLD UCB, as shown below.



Fig. 5. Minimal configuration of Cybersecure Power Router prototype

The following table lists the security threats and mitigations chosen for the Cybersecure Power

Router.

Asset	Threat	Mitigation			
Communication					
Confidentiality	System Surveillance [35]	AES-128 Encryption [36]			
Integrity	Corrupted Firmware [37]	Error Detection			
Availability	Unauthorized User Access [38]	HW-Asst. Monitor, Key Mgmt.			
Firmware					
Distribution	Tampered Firmware [39]	Encryption, Error-Detection			
Installation	Reduced Integrity	MD5 Hash Check			
Loading	System Downtime [40]	Control Multiplexing			
Runtime	Operation Outside Parameters [41]	Heartbeat, HW-Asst. Monitor			
Hardware					
Authenticity	Counterfeit Hardware [42]	Hardware Authentication			
<b>Power Processing</b>					
Quality	Harmonic Distortion [43]	Robust Hardware/Controller Design			
Availability	System Downtime [44]	Control Multiplexing			
Response	Non-Recoverable State [45]	Robust Controller Design			

TABLE III: Cyber-Physical Attack Matrix

Chris Farnell's contributions to the Cybersecure Power Router project

The asset inventory, threats, and mitigations are not exhaustive. The above table shows the current

work done and where the security is implemented.

### **3.2 Power Electronics**

The power electronics of the Power Electronics Evaluation Unified Controller Board consist of switch-mode power supplies, signal chains for controls and feedback, isolated power supplies, voltage and current sensors, filters, ports, ancillary circuits, and human-machine interfaces. Control signals routed from the Complex Programmable Logic Device enter the Power Electronics Evaluation Unified Controller Board through a 40 pin insulation-displacement connector. The control signals then pass through 120  $\Omega$  resistors to trigger HCPL-3120-300E optocouplers. The optocouplers then drive the STGP15H60DF insulated-gate bipolar transistors that act as switches. An isolated, flyback regulator using the LT3748EMS integrated circuit energizes the optocouplers to drive switching. The switch-mode power supplies of the board include an asynchronous buck converter, an asynchronous boost converter, and a three-phase inverter/rectifier. The buck converter provides current sensing before and after the 560 $\mu$ H, 100 $\mu$ F LC filter; and output voltage sensing. The buck converter is rated for an input voltage of 9 to 50 Vdc, and an output voltage of 0 to 50 Vdc. The voltage ratings can be increased if higher voltage rated capacitors are used on the input and output of the buck converter.

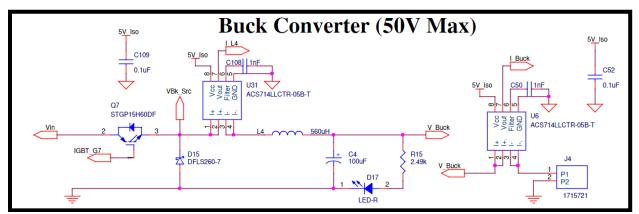


Fig. 6. Asynchronous buck converter schematic of Power Electronics Evaluation Board of the UCB project

The asynchronous boost converter provides current sensing at the input and output; and output voltage sensing. The boost converter is rated for an input voltage of 9 to 50 Vdc, and an output voltage of 9 to 50 Vdc. The voltage ratings can be increased if higher voltage rated capacitors are used on the input and output of the boost converter.

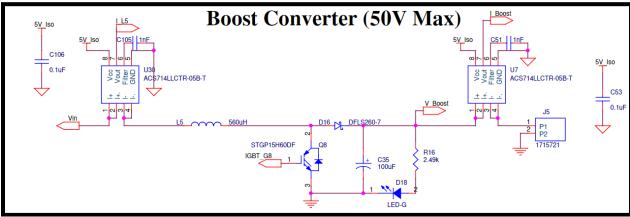


Fig. 7. Asynchronous boost converter schematic of Power Electronics Evaluation Board of the UCB project

The three-phase inverter/rectifier can operate bi-directionally. When acting as an inverter, it can output a 1 Hz to +1000 Hz sinusoid from 0 to 50 Vac. The inverter is rated for 0 to 50 Vdc input. Current sensing is available on all three phases after the inductive filtering, and after capacitive filtering on phase A. Voltage sensing is available on the output of all three phases.

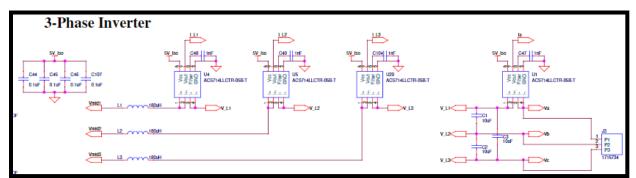


Fig. 8. 3-Phase inverter filter and current sensing schematic of Power Electronics Evaluation Board of the UCB project

Three switching legs using STGP15H60DF insulated-gate bipolar transistors are the switching stage for the inverter/rectifier. As stated earlier, isolated power is provided to the gate drivers. The

low switches share the same isolated power rail. Jumper 1 (JP1) can be used to connect or disconnect the inverter/rectifier with the dc input used on the rest of the Power Electronics Evaluation board. Current and voltage sensing is provided between the dc bus and the capacitive filtering of the inverter/rectifier.

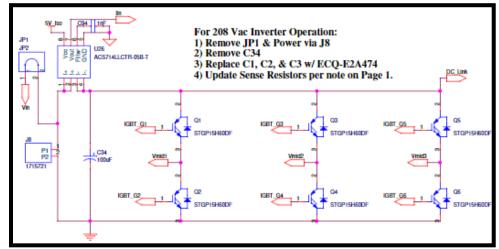


Fig. 9. 3-Phase inverter/rectifier switching stage schematic of Power Electronics Evaluation Board of UCB project

The PCB layout of the Power Electronics Evaluation board is shown below. Considerations were made for mixed analog/digital signals within the board. The isolated dc/dc power supplies and ACPL-C87B-000E optical isolation amplifiers provide isolation between analog signals in the mixed environment and the feedback signals supplied to the DSPs and CPLD.



Fig. 10. Power Electronics Evaluation Unified Controller Board, with Signal Splitter and Hardware Authentication Module, used in Cybersecure Power Router prototype

The total system provides buffered, isolated control and feedback signal chains; and common types of switch-mode power conversion capable of various power flows, with switching frequencies exceeding 100 kHz.

#### 3.3 Digital Signal Processor Board

Two Digital Signal Processor Unified Controller Boards are used in the Cybersecure Power Router prototype. The modular designs of the Unified Controller Board allow two DSPs to slot into the Complex Programmable Logic Device Unified Controller Board. However, having two, standalone DSP boards provide advantages to the testing of the prototype. One advantage is greater accessibility of tools and probes to the I/O of the boards. Another advantage is more direct access to the serial communication ports of the DSPs. A third advantage to using the standalone DSP UCBs is the ability to apply power independent of the CPLD UCB. This allows DSPs to have controlled power disruptions while still allowing the rest of the CSPR prototype to continue operation. Serial communication with the DSPs in the DSP UCB is provided through a Future Technology Devices International (FTDI) Universal Serial Bus (USB) to Universal Asynchronous Receiver/Transmitter (UART) bridge. This USB to UART bridge is the FT2232D FTDI chip. The receive and transmit lines between the FTDI chip and the DSP are also accessible from the General Purpose Input/Output (GPIO)-28 and GPIO-29 pins. This configuration allows USB connectivity with a computer and with a UART device. For the Cybersecure Power Router prototype, the USB connectivity is used to control the DSP from Code Composer Studio and a LabVIEW script; and the UART connectivity is used to control the DSP from the CPLD UCB.

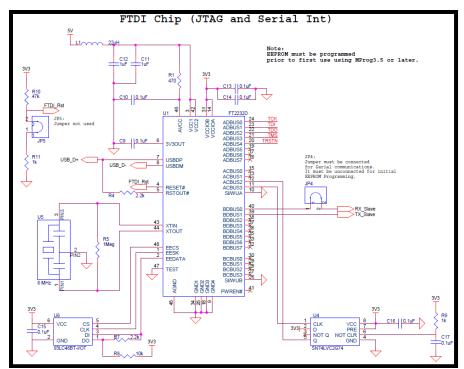


Fig. 11. USB to UART schematic of Digital Signal Processor board of the UCB project

The analog and digital I/O of the DSP within the DSP UCB is readily available to various tools and probes. A connection not pictured is the reset manually added to the TPS3828  $\overline{RESET}$  pin on

U6 of the F28335 DSP controller card. This reset connection cycles power to the DSP during firmware patching.

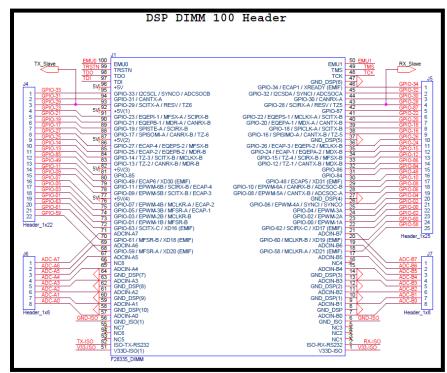


Fig. 12. DIMM pinout of Digital Signal Processor board of the UCB project

Two, 40-pin insulation-displacement connectors (IDCs) are used to bus analog and digital signals between the other Unified Controller Boards. The 28 digital signals use GPIO-00 through GPIO-27. The eight analog feedback signals use ADC-A0 through ADC-A7 for IDC A, and ADC-B0 through ADC-B7 for IDC B. Two 5 Vdc and Ground signals are provided within the 40-pin IDC.

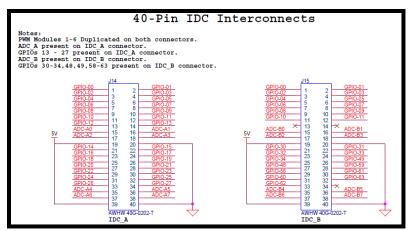


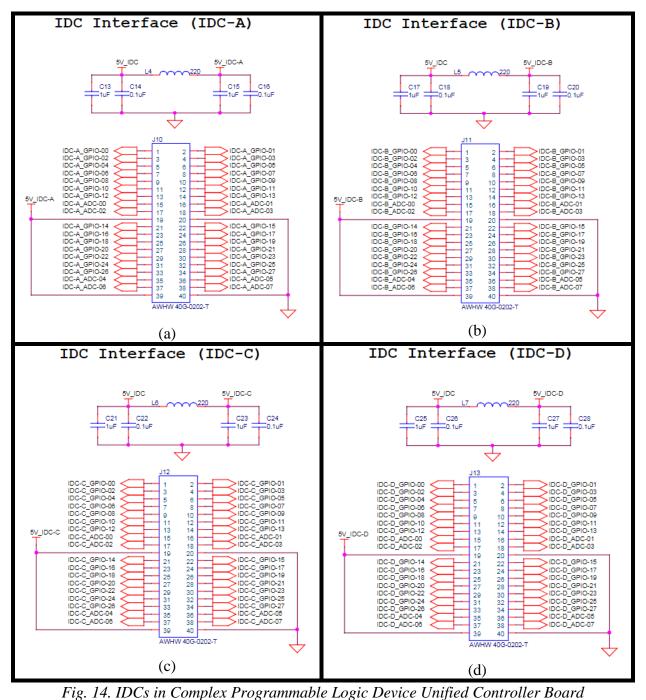
Fig. 13. Majority of analog and digital I/O used by Digital Signal Processor board of UCB project

The Digital Signal Processor Unified Controller Board differs from the Texas Instruments C2000 DAF docking station source design. Mechanical holes are included to allow the use of standoffs. An isolated power supply provides steady power and noise rejection. The isolated power also overcomes ground loops that may be very problematic in noisy environments with long communication cables (such as those used in this prototype).

### 3.4 Complex Programmable Logic Device Board

The Complex Programmable Logic Device Unified Controller Board is the core of the Cybersecure Power Router. It routes digital control and communication signals between other Unified Controller Boards. It also instantiates the VHDL modules that provide functionality essential to the CSPR project. Four IDCs are provided to allow up to four converters or controllers to interface with each other. Presently, the IDCs allow converters or controllers to interface with one another and the CPLD. Each IDC provides 28 digital, GPIO channels; eight analog channels (typically as part of a feedback signal chain); two 5 Vdc output pins; and two Ground pins. High

frequency filtering is provided between the 5 Vdc power of the CPLD UCB and the 5 Vdc outputs of the IDCs.



The Complex Programmable Logic Device is Lattice's LCMXO2-7000HC. The IC is packaged as a TQFP with 144 pins. The XO2-7000 has the greatest resources of the MachXO2 family [46].

		XO2-256	XO2-640	XO2-640U <sup>1</sup>	XO2-1200	XO2-1200U1	XO2-2000	XO2-2000U1	XO2-4000	XO2-7000
LUTs		256	640	640	1280	1280	2112	2112	4320	6864
Distributed RAM (k	bits)	2	5	5	10	10	16	16	34	54
EBR SRAM (kbits)		0	18	64	64	74	74	92	92	240
Number of EBR SF kbits/block)	AM Blocks (9	0	2	7	7	8	8	10	10	26
UFM (kbits)		0	24	64	64	80	80	96	96	256
Device Options:	HC <sup>2</sup>	Yes	Yes	Yes	Yes	Yes	Yes	Yes	Yes	Yes
	HE <sup>3</sup>						Yes	Yes	Yes	Yes
	ZE <sup>4</sup>	Yes	Yes		Yes		Yes		Yes	Yes
Number of PLLs	ł	0	0	1	1	1	1	2	2	2
Hardened	I2C	2	2	2	2	2	2	2	2	2
Functions:	SPI	1	1	1	1	1	1	1	1	1
	Timer/Coun- ter	1	1	1	1	1	1	1	1	1

TABLE IV: MachXO2 Family Features

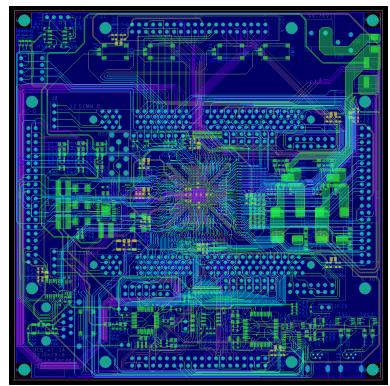
The CPLDs of the MachXO2 family also provide a number of internal clock frequencies for use, in addition to the phase lock loops provided.

MCLK (MHz, Nominal)	MCLK (MHz, Nominal)	MCLK (MHz, Nominal)		
2.08 (default)	9.17	33.25		
2.46	10.23	38		
3.17	13.3	44.33		
4.29	14.78	53.2		
5.54	20.46	66.5		
7	26.6	88.67		
8.31	29.56	133		

TABLE V: Available MCLK frequencies

The extensive I/O of the MachXO2-7000HC is predominantly utilized by the GPIO of IDC A, B, C, and D; connections with the slots designed for DSP cards; serial peripheral interface with an analog to digital converter; Flash programmer/JTAG interface; button inputs; and an LED display. IDC A GPIO, dual in-line package (DIP) switches, push buttons, and four wire communication with the Lattice Flash Programmer connect to Bank 0 of the MachXO2-7000HC. Dual In-Line Memory Module (DIMM) slots connect DSPs cards to the MachXO2-7000HC. The connections between the DIMM-B slot and the MachXO2-7000HC are through Bank 1. The connections between the DIMM-C slot and the MachXO2-7000HC are through Bank 2. IDC C also connects with the MachXO2-7000HC through Bank 2. LED1 through LED8 interface with the MachXO2-

7000HC through Bank 1. The XPort connections interface with MachXO2-7000HC with the Lantronix Ethernet adapter. This interface is not currently developed. Serial communication interface (SCI) receive and transmit port are connected to Bank 3. The two, serial peripheral interface (SPI) analog to digital converters (ADCs) are also connected to Bank 3. The serial communication with the two FTDI chips (one for the DSPs, the other for the CPLD) interface at Bank 3. Bank 4 of the MachXO2-7000HC provides an external clock and reset signal. Finally, Bank 5 provides connections to the IDC D port.



*Fig. 15. PCB layout of Complex Programmable Logic Device board of UCB project* The components for the Complex Programmable Logic Device Unified Controller Board were reflowed using the Sikama 5/C reflow furnace in the Assembly Laboratory in the High Density Electronics Center at the University of Arkansas. A tin, silver, and copper alloy (SAC305) solder paste was applied using a stencil and squeegee. No solder paste was applied for the MachXO2-7000HC. A thin layer of no clean flux was applied to the pads of the TQFP footprint on the

Complex Programmable Logic Device Unified Controller Board PCB. After the application of solder paste, flux, and the population of surface mount components, the boards were flowed in the Sikama 5/C reflow furnace.

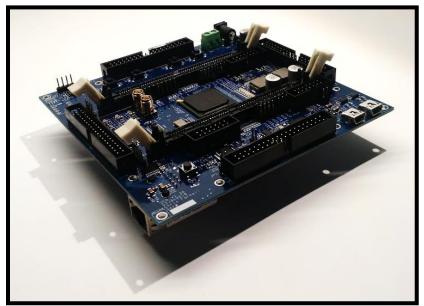


Fig. 16. Fabricated Complex Programmable Logic Device board of UCB project used in Cybersecure Power Router prototype

After inspection and debugging, the through-hole components were populated. Another round of visual, manual, and electrical inspection tested for shorts, unconnected legs, high integrity solder joints, and correct part orientation. Finally, a debug utility was flashed onto the CPLD and tested the I/O.

### 3.5 Signal Splitter

Two DSPs are used in the Cybersecure Power Router prototype. Two DSP Unified Controller Boards were used to interface the DSPs to the rest of the prototype. To route the analog feedback signals, eight Y connections needed to be created from the output of the Power Electronics Evaluation Unified Controller Board to the ADC input of the two DSP UCBs. A simple printed circuit board was designed and fabricated to provide this split analog feedback signal. No buffering is provided on the Signal Splitter on account of existing buffered output from the PE Eval UCB.

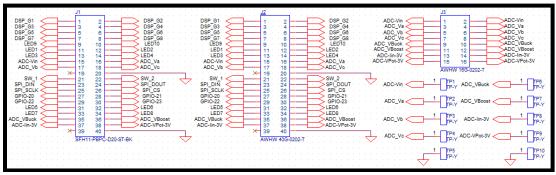


Fig. 17. Complete schematic of Signal Splitter

Test points are provided for each analog signal. Two ground test points are provided to reduce the size of the ground loop of a probe. Female/Female jumper wires connect the analog output of the Signal Splitter PCB to the ADC inputs of the two DSP UCBs.

# **3.6 Hardware Authentication Module**

An embedded, 1 kilobyte password is used to authenticate the power electronics of the Cybersecure Power Router. This password is continually checked by the hardware-assisted monitor instantiated in the CPLD UCB. The password is stored in both the Microchip 93LC46BT-I/OT EEPROM and within the memory of the hardware-assisted monitor.

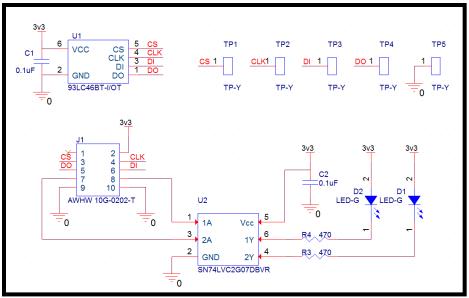


Fig. 18. Complete schematic of Hardware Authentication Module

The Hardware Authentication Module and the CPLD UCB communicate with a four-wire protocol, through the IDC D port. A Texas Instruments SN74LVC2G0 is included to drive two indicator LEDs. The pins for power, ground, the four-wire protocol, and the two indicator LEDs are routed through a 10-pin header.

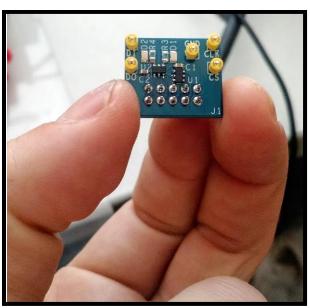


Fig. 19. Hardware Authentication Module used in Cybersecure Power Router prototype

This assembly plugs into the J7 header on the Power Electronics Evaluation Unified Controller Board. Test points are included for the four-wire protocol and a ground to minimize the ground loop during oscilloscope data collection.

# **3.7 BeagleBone Black**

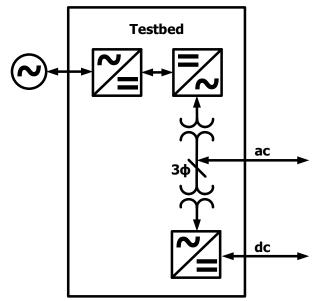
Texas Instruments' BeagleBone Black provides an embedded Linux environment for sophisticated, high-level operations.



Fig. 20. BeagleBone Black

The BeagleBone Black uses a 1 GHz ARM Cortex-A8 processor, and runs Debian Linux. The CPLD UCB and the BeagleBone Black communicate serially. The embedded operating system provides a platform for Python scripts and powerful utilities for automation, analysis, and debugging.

### 3.8 Test Bed



A testbed provides controlled ac and dc power flow to and from the Cybersecure Power Router.

Fig. 21. Block diagram of testbed used in Cybersecure Power Router project

The testbed functions like a microgrid, allowing power assets to be added in a variety of ways while maintaining controlled power flow and a utility frequency (independent of the utility frequency of the electrical grid). The ac power is three phase, and galvanically isolated through low frequency transformers.

#### **3.9 Power Flow**

The Cybersecure Power Router shows a security-by-design process and defense-in-depth methods for a Distributed Energy Resource (DER). The security-by-design process outlines both the assets and their dependencies to be secured. To create this inventory of assets and dependencies, a specific device or system must be chosen. Presently, a modular, grid-tied inverter/rectifier distributed energy resource is chosen. The power assets of this DER include onsite energy generation, energy storage, and bi-directional power flow with the grid.

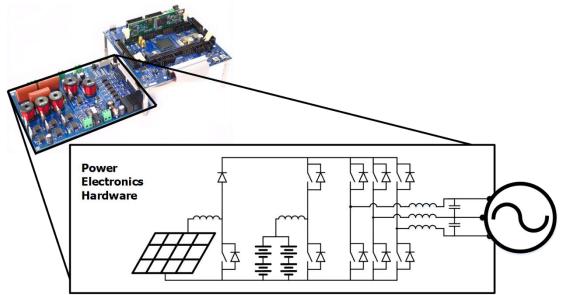


Fig. 22. Simplified block diagram of grid-connected power flow capabilities of UCB hardware and Cybersecure Power Router

Other resources can be integrated within the system. The diagram shows a three-phase inverter/rectifier working in tandem with a photovoltaic panel and battery energy storage system. The PE Eval UCB includes an asynchronous buck and boost converter that may be used to interface dc energy resources, in addition to the three-phase inverter/rectifier provided.

# 3.10 Data Flow

A hardware-assisted monitor and other utilities are instantiated within the complex programmable logic device on the CPLD Unified Controller Board. The monitor and utilities are developed in the Lattice Diamond integrated development environment. The monitor and utilities are developed using the Very high speed integrated circuit Hardware Description Language (VHDL). The complete source code is included in Appendix B. Not included in the appendix are the Intellectual Property (IP) cores used within Lattice Diamond, such as the phase-locked loop (PLL) or the digital memory. An oscillator is instantiated within the CPLD to provide an internal clock at 53.2 MHz. This clock is supplied to an internal phase-locked loop. A 53.2 ( $1 \cdot f_{clk}$ ), 24.93 ( $1/2 \cdot f_{clk}$ ), and 1.5 ( $1/32 \cdot f_{clk}$ ) MHz clock signal is derived from the original 53.2 MHz clock signal from the oscillator. The 53.2 MHz clock is used for the hardware-assisted monitor and other high speed applications. The 24.93 MHz clock is used for serial communication and the data bus. The 1.5 MHz clock is used for the four-wire communication with the Hardware Authentication Module.

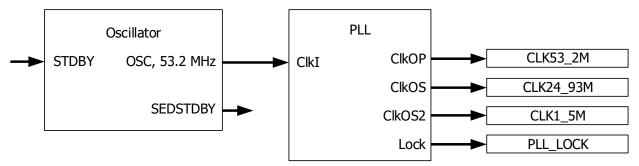


Fig. 23. Block diagram of clock generation within Complex Programmable Logic Device of the Cybersecure Power Router

A data bus is instantiated to allow data flow between various modules. The bus uses 16-bit addresses and 16-bit data. Prioritized access to the bus is given to modules. Currently, ten modules can be prioritized according to access privileges. In addition to controlling the data bus, the Bus Master contains Randomly Accessible Memory (RAM). This memory is used as registers for various controls and functions. A separate memory allocation is used for the booting partition of the digital signal processors.

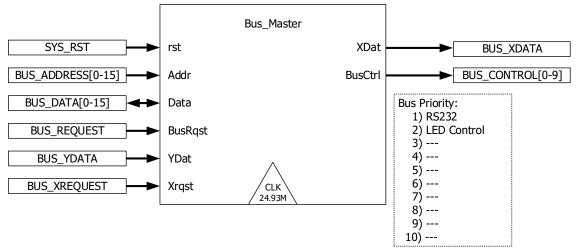


Fig. 24. Block diagram of data bus controller within Complex Programmable Logic Device of the Cybersecure Power Router

The serial communication with the modules instantiated in the CPLD uses a 9600 baud rate, and connects through Bank 3 of the MachXO2-7000HC CPLD. The current serial communication is designed for fixed packet lengths.

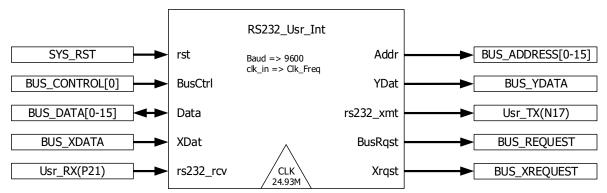


Fig. 25 .Block diagram of serial interface within Complex Programmable Logic Device of the Cybersecure Power Router

# 3.11 Control Multiplexing

A defense-in-depth approach to controller security is explored in the Cybersecure Power Router. Control multiplexing strengthens the availability and integrity of the hardware controller, and the entire system by extension. The concept of multiplexing is common in telecommunication [47], computer networks [48], and various signal conditioning and sampling [49] contexts. The concept is extended to an entire bus of control signals for the Cybersecure Power Router. The controllers running on the DSPs toggle a bit on GPIO-24 every time a switching cycle is completed. The firmware snippet is provided below.

783	
784	<pre>if (HeartBeat_High){</pre>
785	<pre>GpioDataRegs.GPASET.bit.GPI024 = 1;</pre>
786	HeartBeat_High = false;
787	}else{
788	<pre>GpioDataRegs.GPACLEAR.bit.GPI024 = 1;</pre>
789	HeartBeat_High = true;
790	}
791	

Fig. 26. Firmware code snippet to generate heartbeat from Digital Signal Processors

The toggled GPIO-24 pin creates a clock signal visible to external devices.

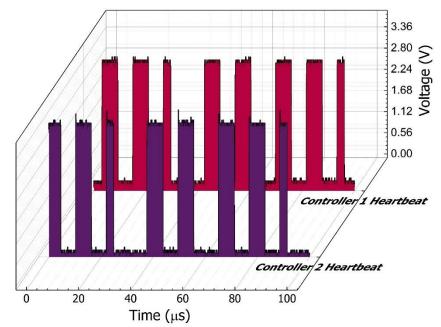


Fig. 27. Heartbeats of Controllers 1 and 2 while running identical firmware

This clock signal is the heartbeat of the controller, and is used by the Hardware Assisted Monitor to assess the liveness of the controller. Presently, if the heartbeat of a controller beats more often than 75  $\mu$ s, it is considered to maintain liveness. If the heartbeat takes longer than the given 75  $\mu$ s to toggle, the Hardware Assisted Monitor considers the controller to have lost liveness. The period

of the heartbeat is a function of the controller's clock rate and execution cycles of the firmware. A slower processor or a longer execution cycle would require a longer period between heartbeats.

The security features instantiated in the CPLD communicates with the Hardware Authentication Module on the UCB PE Eval board to authenticate the power electronic hardware. The password stored in the EEPROM of the UCB Hardware Authentication Module is checked against the password stored in the memory of the CPLD Hardware Authentication Module.

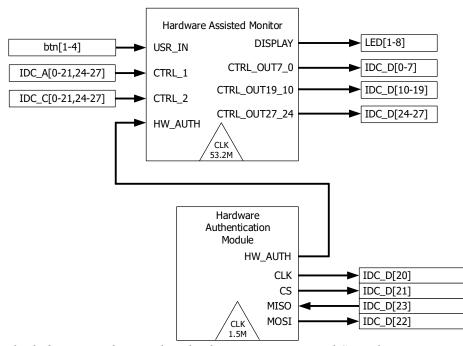


Fig. 28. Block diagram of control multiplexing using Digital Signal Processor signals and Hardware Authentication Module within the Complex Programmable Logic Device of the Cybersecure Power Router

If the two passwords match, then the "hardware is authorized" (HW\_AUTH) signal goes HIGH (TRUE). If there is a mismatch between the symmetric keys, then the HW\_AUTH signal goes LOW (FALSE). This mismatch will occur when the key stored in the power electronics (i.e., the Hardware Authentication Module) differs from the key stored in the controller (i.e., the Hardware Assisted Monitor).

The Hardware Assisted Monitor uses the liveness of the controllers and the authentication of the power electronics to decide the routing of control signals. When the hardware is authenticated, control signals from Controller 1 or Controller 2 are routed to the power electronics. The Hardware Assisted Monitor assigns priority to Controllers 1 over Controller 2 if both controllers have liveness. If only one controller has liveness, that controller's control signals are routed to the power electronics. If no controller has liveness, the hardware is held in a lockout state. User inputs from buttons 1 through 4 can override this logic to manually set the routing of control signals.

#### 3.12 Firmware and Boot Management

The control multiplexing behavior of the Hardware Assisted Monitor prevents downtime during firmware updates. The firmware is loaded into memory instantiated in the CPLD UCB allocated for boot loading. The firmware is loaded through encrypted serial communication and the internal data bus.

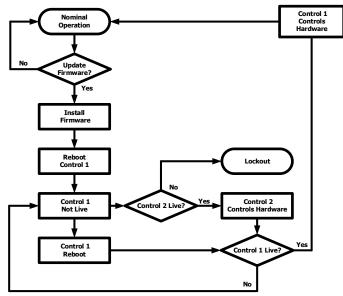


Fig. 29. Block diagram of hot patching process

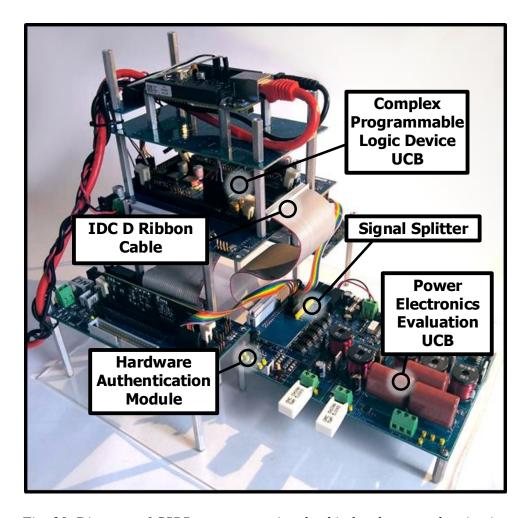
When a command is given, the designated DSP is power cycled by the CPLD. The power cycled DSP boots from the allocated memory hosting the new firmware. The DSP boots using this new firmware. While the DSP is booting, hardware control is passed to the second DSP. The hardware

continues operating while the first DSP boots with the new firmware. When the first DSP resumes operation and provides a heartbeat with a period less than 75  $\mu$ s, it is either passed control of the hardware, or remains on standby. The above block diagram illustrates this process for Controller 1 being updated. A similar process is used for updating the firmware of Controller 2. As of the time of this writing, this uptime during update process is being developed.

Hot patching refers to modifying currently used data in system memory. Hot patching, strictly speaking, refers to a process that only applies to software. The present process is similar, but works at the firmware and hardware level. Here, the process modifies currently used data flow (like control signals) in a running system. The result of both techniques is the same: a running system while patches, updates, and other fixes are applied.

#### 3.13 Hardware Authentication

Authentication of the Power Electronics Evaluation Unified Controller Board requires the Hardware Authentication Module PCB, the PE Eval UCB, and the CPLD UCB. The process begins when power is applied to the PE Eval UCB. The on-board power is used to energize the Hardware Authentication Module PCB. The IDC D port is used to connect the PE Eval UCB and the CPLD UCB. In the CSPR prototype, the IDC D ribbon cable plugs into the Signal Splitter board, which plugs into the PE Eval UCB.

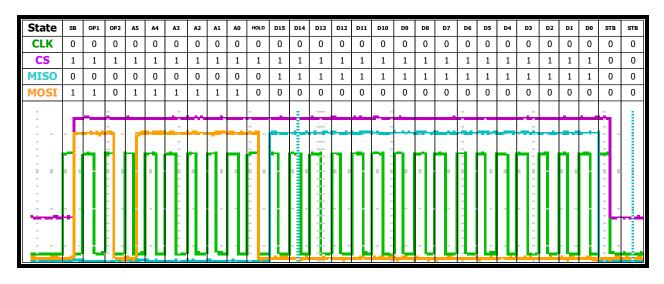


*Fig. 30. Diagram of CSPR components involved in hardware authentication* The CPLD UCB interfaces with the Hardware Authentication Module PCB through four connections in the IDC D connection. These four connections provide the chip select (CS), clock (CLK), master in slave out (MISO), and master out slave in (MOSI) signals between the EEPROM of the Hardware Authentication Module on the PE Eval UCB and the VHDL Hardware Authentication Module instantiated in the CPLD UCB. The clock signal provided is 1.5 MHz. Only the read command for the EEPROM is used in the authentication process. To read a specific memory address from the 93LC46BT-I/OT EEPROM [50], the following steps are required.

Unaccounted clock cycle here.         TABLE 1-3: INSTRUCTION SET FOR X16 ORGANIZATION (93XX46B OF 93XX46C WITH ORG = 1)										
SB	Opcode	Address			Data In	Data Out	Req. CLK Cycles			
1	11	A5	A4	A3	A2	A1	A0	- /	(RDY/BSY)	9
1	00	1	0	Х	Х	Х	Х	_	(RDY/BSY)	9
1	00	0	0	Х	Х	Х	Х		High-Z	9
1	00	1	1	Х	Х	Х	Х	_	High-Z	9
1	10	A5	A4	A3	A2	A1	A0	_	D15 - D0	25 <b>26</b>
1	01	A5	A4	A3	A2	A1	A0	D15 - D0	(RDY/BSY)	25
1	00	0	1	Х	Х	Х	Х	D15 - D0	(RDY/BSY)	25
_		B         Opcode           1         11           1         00           1         00           1         00           1         00           1         00           1         00           1         00           1         00           1         00           1         01	B         Opcode           1         11         A5           1         00         1           1         00         0           1         00         1           1         00         1           1         00         1           1         00         1           1         00         1           1         00         1           1         01         A5	B         Opcode           1         11         A5         A4           1         00         1         0           1         00         0         0           1         00         1         1           1         00         1         1           1         00         1         4           1         00         1         1           1         10         A5         A4           1         01         A5         A4	B         Opcode         Add           1         11         A5         A4         A3           1         00         1         0         X           1         00         0         0         X           1         00         1         1         X           1         00         1         1         X           1         10         A5         A4         A3           1         00         1         1         X           1         00         1         4         A3           1         01         A5         A4         A3	B         Opcode         Address           1         11         A5         A4         A3         A2           1         00         1         0         X         X           1         00         0         0         X         X           1         00         1         1         X         X           1         00         1         1         X         X           1         00         1         1         X         X           1         10         A5         A4         A3         A2           1         01         A5         A4         A3         A2	B         Opcode         Address           1         11         A5         A4         A3         A2         A1           1         00         1         0         X         X         X           1         00         1         0         X         X         X           1         00         1         1         X         X         X           1         00         1         1         X         X         X           1         00         1         1         X         X         X           1         10         A5         A4         A3         A2         A1           1         00         1         4         A3         A2         A1           1         01         A5         A4         A3         A2         A1	B         Opcode         Address           1         11         A5         A4         A3         A2         A1         A0           1         00         1         0         X         X         X         X           1         00         0         0         X         X         X         X           1         00         1         1         X         X         X         X           1         00         1         1         X         X         X         X           1         00         1         1         X         X         X         X           1         10         A5         A4         A3         A2         A1         A0           1         10         A5         A4         A3         A2         A1         A0           1         01         A5         A4         A3         A2         A1         A0	B         Opcode         Address         Data In           1         11         A5         A4         A3         A2         A1         A0         —           1         00         1         0         X         X         X         X         —           1         00         0         0         X         X         X         X         —           1         00         1         1         X         X         X         —         —           1         00         1         1         X         X         X         —         —           1         00         1         1         X         X         X         —         —           1         00         1         1         X         X         X         —         —           1         10         A5         A4         A3         A2         A1         A0         —           1         01         A5         A4         A3         A2         A1         A0         D15 - D0	STRUCTION SET FOR X16 ORGANIZATION (93XX46 OR 93XX46           B         Opcode         Address         Data In         Data Out           1         11         A5         A4         A3         A2         A1         A0         —         (RDY/BSY)           1         00         1         0         X         X         X         —         (RDY/BSY)           1         000         1         1         X         X         X         —         High-Z           1         00         1         1         X         X         X         —         High-Z           1         00         1         1         X         X         X         —         High-Z           1         00         1         1         X         X         X         —         High-Z           1         10         A5         A4         A3         A2         A1         A0         —         D15 - D0           1         01         A5         A4         A3         A2         A1         A0         D15 - D0         (RDY/BSY)

Fig. 31. Excerpt from EEPROM datasheet, with erroneous information noted

First, the chip select must go high and remain high for the duration of the instruction. Sent data on the MOSI signal is read at the falling edge of the clock. The starting bit of "1" is given on the MOSI signal. The operational code for read, "10", is given over the next two clock cycles. A 6-bit address is then supplied. A hold cycle, not present in the datasheet, but present in the captured waveform below, is provided to transition between the final bit of the address (read on the falling edge of the clock) and the first data out (sent on the rising edge of the clock).



*Fig. 32: Hardware Authentication Module communication beside oscilloscope capture* A 16-bit value is read from each address. After transmission, MOSI and CS are set low. The system is now in standby, and ready for another instruction cycle. The read values are parts of a 1 kilobit

password. There are 64 values in total. These values are flashed onto the EEPROM with an Arduino running the "EEPROM\_Write\_PASSWORD" program (source code provided in Appendix B). A matching password is included in the VHDL Hardware Authentication Module instantiated in the CPLD UCB. These values are polled by the Hardware Authentication Module in the CPLD UCB approximately five times a second. While the passwords match, a HW\_AUTH ("hardware is authorized") signal is provided to the Hardware Assisted Monitor to allow control signals to pass to the hardware. The result is that the CSPR hardware will only run if the Hardware Authentication Module is operational.

## **3.14 Submodule Encrypted Communication**

Communication between components of the CSPR prototype is designed to be encrypted. This encrypted communication is the serial communication between the two digital signal processors and the complex programmable logic device, and serial communication between the BeagleBone Black and the complex programmable logic device.

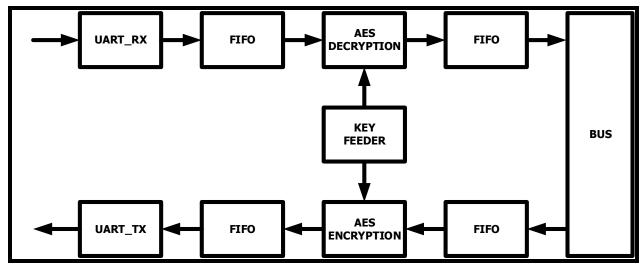


Fig. 33. Block diagram of encrypted serial communication within Complex Programmable Logic Device used in the Cybersecure Power Router

The Advanced Encryption Standard 128-bit (AES-128) is used for this encryption. A key feeder supplies the 128-bit key to the two AES-128 instantiations used to encrypt and decrypt serial communication. This allows a rolling key to be used in later development of the CSPR project. Additionally, different keys can be used for each communication port, and have separate keys for receiving and transmitting serial information. First In, First Out (FIFO) buffers are used to queue data until the appropriate block length is reached for encryption or decryption. Serial communication is made available to the data bus instantiated within the CPLD UCB. Presently, superficial, bidirectional, encrypted exchanges are available between the CSPR and the test bed. This sets the groundwork for an encrypted, MD5 secured communication pipeline for later development of the CSPR project.

#### **3.15 Hardware Protections**

Simple hardware protections are available for instantiation in the Cybersecure Power Router prototype. These hardware protections prevent shoot-through faults in the switching legs of the PE Eval UCB. These faults result when the control signals of both the high and low switch positions are set to HIGH. This creates a direct connection between the rails of the dc bus of the PE Eval UCB. The digital signal processors used in the CSPR prototype have shoot-through protections built into the PWM modules that generate the control signals of the switching legs. These built in protections can be reinforced by the Hardware Assisted Monitor of the CPLD UCB. The logic of the shoot-through protection is to allow only one switching position to be ON at a time, and both to be OFF in all other cases.

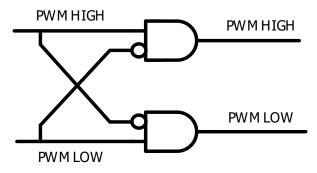


Fig. 34. Shoot-Through hardware protection

A delay or jitter of this protection on the order of 18.80 µs could be introduced by the sampling resolution of the CPLD UCB's Hardware Assisted Monitor. The resolution can be found with the following simple equation:

$$\frac{1}{f_{clock}} = T_{period} \tag{2}$$

Given the 52.3 MHz clock supplied to the Hardware Assisted Monitor, the resolution is limited to:

$$\frac{1}{53.2 \, MHz} = 18.7969 \, \mu s \tag{3}$$

The resolution may be increased with a faster clock, up to the limitation of the hardware. CPLD cost, power demands, parasitic capacitance from layout, and EMI are all likely to limit the clock frequency at which the Hardware Assisted Monitor operates.

### 3.16 Display

A bank of LEDs are used to display various system information of the CPLD UCB.

LED(8)	LED(7)	LED(6)	LED(5)	LED(4)	LED(3)	LED(2)	LED(1)
Hardware IS NOT Authorized	Hardware IS Authorized	BBB Serial Comm. Operational	DSP Serial Comm. Operational	HARDWARE LOCKOUT	NOMINAL CONTROL	ONLY CONTROLLER 1	ONLY CONTROLLER 2

Fig. 35. LED display diagram for Cybersecure Power Router

LED(8) and LED(7) are dedicated to displaying the status of hardware authentication. The red LED(8) is ON and the greed LED (7) is OFF if the hardware is not authenticated. The red LED(8) is OFF and the greed LED (7) is ON if the hardware is authenticated. The yellow LED(6) is dedicated to displaying the serial connectivity between the BeagleBone Black and communication modules instantiated in the CPLD UCB. While serial connectivity is active, LED(6) will fade ON and OFF. LED(5) is dedicated to displaying the serial connectivity between the digital signal processors and the communication modules instantiated in the CPLD UCB. Finally, LED(4-1) are dedicated to display the state of the Hardware Assisted Monitor. LED(4) indicates a Lockout state. LED(3) indicates a nominal state of both Controller 1 and Controller 2 being live. LED(2) indicates that only Controller 1 is live. LED (1) indicates that only Controller 2 is live.

## **Chapter 4 - Results**

The maximum duration of time between firmware execution cycles is sampled across switching frequencies. Specifically, GPIO-24 of Controller 1 is probed, and measured for the maximum pulse width. Ideally, this maximum pulse width corresponds to the duration of time between interrupt service routines triggered at every switching cycle. An ideal trend line of duration between switching periods is included in teal, below. An ideal trend line of available controller resources is included in purple. These samples are plotted in orange as switching frequency is swept from 1 kHz to 160 kHz.

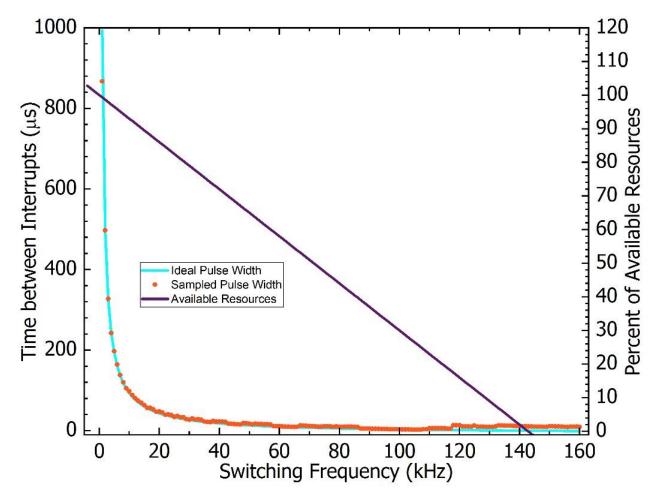


Fig. 36. Time between execution cycles of controller firmware vs. switching frequency

The Cybersecure Power Router can operate up to 2050 kHz, but the samples produce a monotonic trend beyond 120 kHz. The detail of Fig. 36 shows this transition in trends.

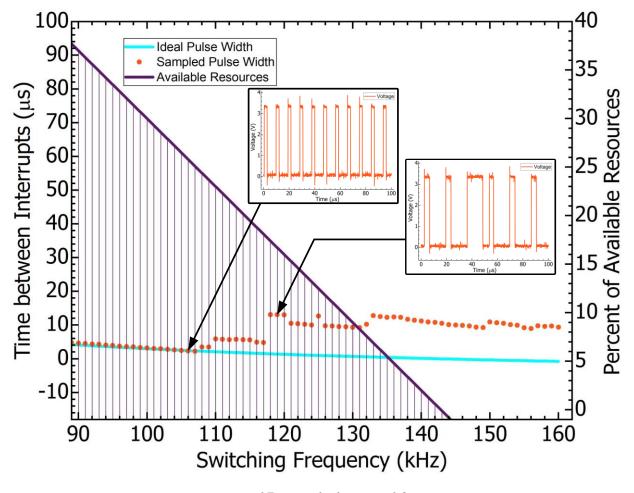


Fig. 37. Detail of Figure 36

The inverter is set to a 60 Hz output. The voltage waveform of the inverter output is captured as the switching frequency is set to 114, 118, and 128 kHz. Compare these switching frequencies to those in the figure above (Fig. 37) to see the threshold of an increasing switching frequency to degrading and decreasing the inverter output frequency. These trends are intrinsic to a controller's hardware. The operate at higher switching frequencies, all things being equal, a controller needs to do more processing in less time. This usually requires are more powerful controller.

The duration between switching period interrupts behaves ideally, until switching frequency is increased beyond 106 kHz. Increasing the switching frequency reduces the time a controller has to complete tasks required for controlling the power electronics. For the current controller hardware and firmware, a deterioration in power flow occurs when switching frequency is increased beyond 114 kHz. The output frequency of the inverter is set to 60 Hz. Yet, as switching frequency is increased, inverter output falls to 45.05 and finally 30.07 Hz.

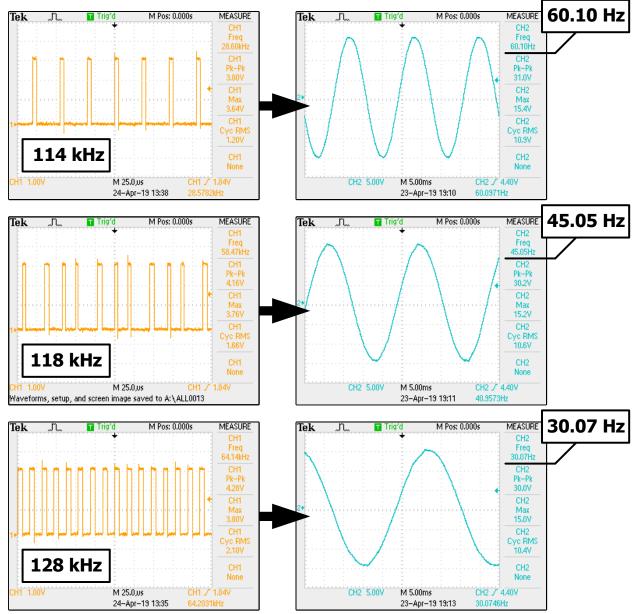


Fig. 38. Inverter output voltage at 114, 118, and 128 kHz switching frequencies.

The interruption of ISRs before completion accounts for this. At 128 kHz switching frequency, only 10% of processing resources are expected to be available for a particular interrupt. This amount of resource utilization provides little insulation between the resources used of one task from another. A possible result is a metastable condition resulting in the inverter output frequency one half of the set output frequency. While an ISR is running, it is interrupted by another ISR. When the new ISR completes, the previous ISR is able to finish. The results of the new ISR are overwritten by the previous ISR.

To safeguard against such overrun conditions, the sensitivity of the Hardware Assisted Monitor to loss of controller liveness can be adjusted. This can be done by adjusting the timing requirements of the controller heartbeat against the expected pulse width (as pictured in figures 36 and 37). The Hardware Assisted Monitor can reroute control away from the deteriorating controller, such as one causing 45.05 Hz or 30.07 Hz inverter output (as pictured in figure 38). If the Hardware Assisted Monitor and heartbeat features are thus employed, the power flow of the power electronics will be protected in case of firmware loss of liveness.

This security feature may result in disruptions and phase shifts if the controllers are not synchronized. Figure 39 shows the voltage waveform of the inverter output while control is rerouted.

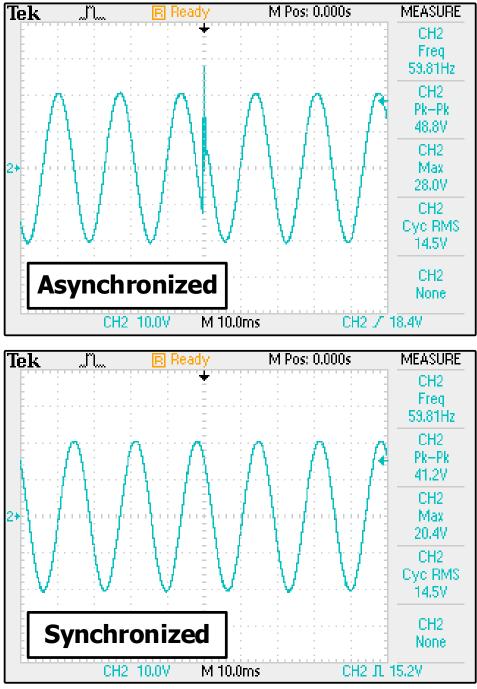


Fig. 39. Inverter output during controller transition

The inset labeled "Asynchronized" shows the possible result of routing control between controllers out of phase with each other. A sharp transient or instability many result as the hardware jumps from one point in the phase to another as control shifts between controllers. This problem is avoided if the controllers are kept in phase, as show in the "Synchronized" inset. Here, both controllers are synchronized to one another, or to the same signal. An example of this is grid-tied inverter controllers kept in phase with one another by locking onto the same grid frequency.

A Hardware Assisted Monitor, firmware modification, and a second controller are required for these security features. These features maintain the operation of the system, rather than cause downtime in case of fault or failure. These features use a co-processor to instantiate the Hardware Assisted Monitor and a second DSP to instantiate the second controller. Both of these design choices incur an economic and non-recurring engineering cost relatively high to the cost of the CSPR prototype. Using the Hardware Assisted Monitor as a second controller or a failsafe is also a possibility. The present work shows one example of a secure system. It also raises many possibilities for new cybersecure architectures that balance security, cost, and performance.

## **Chapter 5 - Future Work**

## **5.1 Multi-Mission Controls**

The Cybersecure Power Router allows great flexibility in the operation of controls. This flexibility can be extended to allow multi-mission controls. The digital signal processor controls used in the present prototype are redundant. This is not by way of necessity, but of convenience. A different controller could operate in each DSP. These controllers could be optimized for energy management, maximally secure operation, communication network facilitation, grid-reliability, or other objectives or missions. Hardware assisted monitoring within the CSPR could provide sufficient situational awareness to route hardware control to different controllers in different contexts. Consider the following as an example. Controller 1 is optimized for efficient use of energy resources. To provide more resources to power processing, communication is limited in both volume and sophistication. Controller 2 is balanced to manage energy resources and provide more secure communication. Controller 1 is used nominally, and accomplishes the primary mission of efficient use of energy resources. If a communication anomaly or attack is detected, Controller 2 is given system control. The transition between the two controllers is made smooth through the control multiplexing technique shown earlier. When the anomaly or attack is cleared, system control can be returned to Controller 1. Another possibility is booting a new controller onto the DSP used by Controller 1 while the system is operated by Controller 2. This would allow a new controller, say Controller 3, to be instantiated. Controller 3 could be optimized for a different mission, say, to perform more conservative power management or provide forensic data in case of hardware failure. Control could be switched from Controller 2 to Controller 3, and the process

could be repeated. In this way, two DSPs within the Cybersecure Power Router could be used to provide controllers with many different missions.

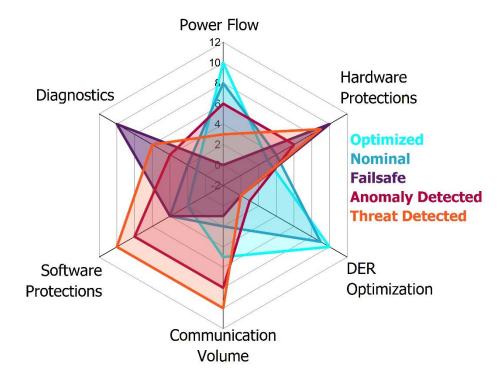


Fig. 40: Radar chart of missions for controls

To reduce cost, controllers could be instantiated in the complex programmable logic device or other hardware of the Cybersecure Power Router. This would remove the need for a second DSP, or possibly both DSPs.

## **5.2 SGPN and CSPR Integration and Completion**

The integration of the Cybersecure Power Router and the Smart Green Power Node are required to realize a secure distributed energy resource pre-production prototype. The Smart Green Power Node contributes hardware designs rated for 2 kW operation, power flow optimization, energy generation prediction, grid arbitrage, and sophisticated controls.

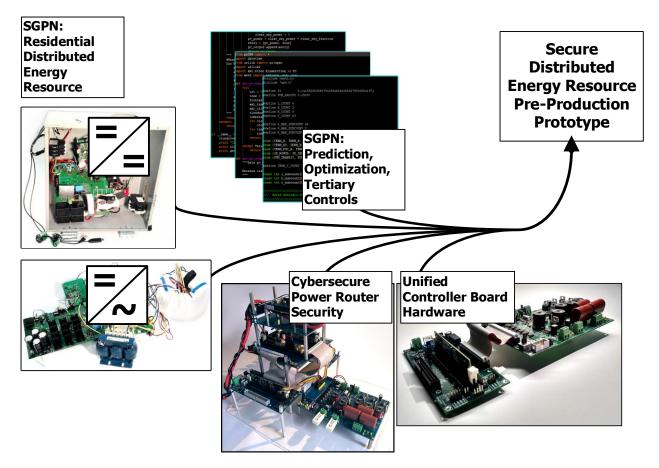


Fig. 41: CSPR and SGPN migration

The Cybersecure Power Router provides tested primary and secondary controls, protected system operation, and enhanced security. Both projects, however, have challenges to their forward development towards a pre-production prototype.

CSPR lacks the power ratings required in the prototype, and is not yet complete as a project. Five milestones remain in the CSPR project: external communication protocols, integrated user/server authentication, exhaustive testing, documentation, and project end. External communication protocols and User/Server authentication are developed, but have not been integrated with the CSPR prototype. Exhaustive testing of the cybersecure inverter is partially accomplished: the hardware has operated over 100 hours at various power levels. Superficial penetration tests on communication during operation were performed, and resulted in no observable changes to system operation. Exhaustive penetration testing, specifically fuzz testing, during operation and hot-patching are suggested in further development. Finally, documentation and project end is also partially accomplished.

The SGPN lacks proper thermal management, optimized PCB layout, and coherent systemlevel design. Many switches are poorly cooled and some devices, such as snubber circuits, receive no forced cooling due to poor planning of thermal management. Electromagnetic interference from switching during typical operation destroyed gate drivers, disrupted serial communication, and deteriorated the integrity of feedback signals. The current SGPN has a poorly defined secondary and tertiary controller. The individual controllers are not able to coordinate power flow in a safe, autonomous way. The coordination and control of power flow through all these devices is manual. The prediction and optimization algorithms provide a simple schedule for charging and discharging the batteries throughout a day. These algorithms lack integration with the coordination and control of the system, and currently provide no improvement to system operation.

Several considerations are necessary to realize a secure DER pre-production prototype. First, an analysis of the design specifications are required to re-evaluate design parameters. The typical operating voltage of the photovoltaic panels, power rating of the dc/dc converters, topology of the dual half bridge, and capabilities of the human-machine interface may warrant re-evaluation. Additionally, the inclusion or exclusion of security features from the CSPR project require consideration. Secondly, the power electronics are suggested to be redesigned with thermal co-design. Forced air convection and an extruded heatsink common to all switching devices may provide a simple and effective means of cooling. The devices may be epoxied to aluminum nitride heat spreaders to provide high thermal conductivity and electrical insulation. Third, radiated and conducted EMI is suggested to be reduced through the reduction of PCB parasitics, and the

inclusion of snubbers and protection circuits during initial design. Finally, a well-defined, systemlevel control scheme is suggested to be designed early in development. A comprehensive, robust scheme is recommended for the integration of: power optimization, prediction, user preference, current protections, battery protections, various modes of operation, and individual controller operation. These schemes are suggested to be well defined before the design of individual converter controls.

### **Chapter 6 - Conclusion**

A security-by-design process identified the key power and data assets and their dependencies for a distributed energy resource. The security-by-design process showed the dependencies between hardware, data, and power assets. Namely, firmware---as source code stored in memory and as a live instantiation as the controller---lies in the center of these dependencies. Defense-indepth was shown by layered security using AES-128 encryption, error detection, hardware assisted monitoring, key management, MD5 hash checking, control multiplexing, heartbeat monitoring, and hardware authentication, and hardware protections. This defense-in-depth protects the integrity, confidentiality, and availability of hardware, data, and power at every layer of design. Communication security includes encryption and error checking of transmitted messages, firmware, and data shared between CSPR modules. Hardware security includes robust controls, shoot-through protection, hardware authentication, galvanic isolation, and hardware failsafe controls of connected resources. Securing the power and data flow through the Cybersecure Power Router primarily means securing the integrity and availability of the hardware controller. The Cybersecure Power Router determines the proper functioning of the controller by means of a hardware assisted monitor and a controller's heartbeat. The Cybersecure Power Router responds to a controller failure by multiplexing control signals, swapping control from a malfunctioning controller to a live one.

The Cybersecure Power Router illustrates the security-by-design process and defense-in-depth method in a single prototype. The Smart Green Power Node was evaluated for present and future use as the hardware of the Cybersecure Power Router prototype. The security features protecting liveness of controllers of the Cybersecure Power Router was researched in depth. Decreased integrity of power flow through the power electronics was shown to correlate with the loss of liveness of the controllers. The use of a heartbeat from the controllers provides a signal sensitive to the liveness of the controller. The Hardware Assisted Monitor uses this heartbeat to provide control signals to hardware from controllers with liveness. The resulting system is resilient to firmware failure and loss of integrity at runtime. This resilience protects against controller overrun and system downtime, such as during firmware patching. Further research can investigate greater flexibility and resiliency in controls, and methods to reduce costs. A pre-production prototype can be realized from the migration of features from both the Cybersecure Power Router and the Smart Green Power Node.

#### References

- [1] O. Ivanchenko, V. Kharchenko, B. Moroz, L. Kabak, and S. Konovalenko, "Risk Assessment of Critical Energy Infrastructure Considering Physical and Cyber Assets: Methodology and Models," in 2018 IEEE 4th International Symposium on Wireless Systems within the International Conferences on Intelligent Data Acquisition and Advanced Computing Systems (IDAACS-SWS), 2018, pp. 225–228.
- [2] C. Konstantinou, M. Maniatakos, F. Saqib, S. Hu, J. Plusquellic, and Y. Jin, "Cyberphysical systems: A security perspective," in 2015 20th IEEE European Test Symposium (ETS), 2015, pp. 1–8.
- [3] A. R. S. Farhan and G. M. M. Mostafa, "A Methodology for Enhancing Software Security During Development Processes," in 2018 21st Saudi Computer Society National Computer Conference (NCC), 2018, pp. 1–6.
- [4] S. Yang, A. Bryant, P. Mawby, D. Xiang, L. Ran, and P. Tavner, "An Industry-Based Survey of Reliability in Power Electronic Converters," *IEEE Trans. Ind. Appl.*, vol. 47, no. 3, pp. 1441–1451, May 2011.
- [5] J. Golovatchev, O. Budde, and Chin-Gi Hong, "Management of product complexity through integrated PLM in a multi-lifecycle environment," in 2009 IEEE International Technology Management Conference (ICE), 2009, pp. 1–9.
- "IEEE Standard for System, Software, and Hardware Verification and Validation," *IEEE Std 1012-2016 Revis. IEEE Std 1012-2012 Inc. IEEE Std 1012-2016Cor1-2017*, pp. 1–260, Sep. 2017.

- [7] J. R. White and M. Doherty, "Hazards in the installation and maintenance of solar panels," in 2017 IEEE IAS Electrical Safety Workshop (ESW), 2017, pp. 1–5.
- [8] Z. Li and M. Shahidehpour, "Defense-in-depth framework for microgrid secure operations against cyberattacks," in 2017 IEEE Power Energy Society General Meeting, 2017, pp. 1–5.
- [9] S. Sridhar, A. Hahn, and M. Govindarasu, "Cyber–Physical System Security for the Electric Power Grid," *Proc. IEEE*, vol. 100, no. 1, pp. 210–224, Jan. 2012.
- [10] A. Mukherjee, S. A. A. Fakoorian, J. Huang, and A. L. Swindlehurst, "Principles of Physical Layer Security in Multiuser Wireless Networks: A Survey," *IEEE Commun. Surv. Tutor.*, vol. 16, no. 3, pp. 1550–1573, Third 2014.
- [11] S. Teleke, M. E. Baran, S. Bhattacharya, and A. Q. Huang, "Rule-Based Control of Battery Energy Storage for Dispatching Intermittent Renewable Sources," *IEEE Trans. Sustain. Energy*, vol. 1, no. 3, pp. 117–124, Oct. 2010.
- [12] E. W. Nahas, D. A. Mansour, H. A. A. el-Ghany, and M. M. Eissa, "Accurate Fault Analysis and Proposed Protection Scheme for Battery Energy Storage System Integrated with DC Microgrids," in 2018 Twentieth International Middle East Power Systems Conference (MEPCON), 2018, pp. 911–917.
- [13] F. Xiao and J. D. McCalley, "Risk-Based Security and Economy Tradeoff Analysis for Real-Time Operation," *IEEE Trans. Power Syst.*, vol. 22, no. 4, pp. 2287–2288, Nov. 2007.
- [14] A. Varghese and A. K. Bose, "Threat modelling of industrial controllers: A firmware security perspective," in 2014 International Conference on Anti-Counterfeiting, Security and Identification (ASID), 2014, pp. 1–4.
- [15] N. R. Saxena and E. J. McCluskey, "Analysis of checksums, extended-precision checksums, and cyclic redundancy checks," *IEEE Trans. Comput.*, vol. 39, no. 7, pp. 969– 975, Jul. 1990.
- [16] A. Cervin, "ANALYSIS OF OVERRUN STRATEGIES IN PERIODIC CONTROL TASKS," *IFAC Proc. Vol.*, vol. 38, no. 1, pp. 219–224, Jan. 2005.
- [17] National Institute of Standards and Technology, "FIPS 197, Advanced Encryption Standard (AES)," *Fed. Inf. Process. Stand. Publ.* 197, p. 51, Nov. 2001.
- [18] "ISO/IEC 18033-3:2010 Information technology -- Security techniques -- Encryption algorithms -- Part 3: Block ciphers." [Online]. Available: https://www.iso.org/standard/54531.html. [Accessed: 22-May-2019].

- [19] V. Saicheur and K. Piromsopa, "An implementation of AES-128 and AES-512 on Apple mobile processor," in 2017 14th International Conference on Electrical Engineering/Electronics, Computer, Telecommunications and Information Technology (ECTI-CON), 2017, pp. 389–392.
- [20] M. P. Priyanka, E. L. Prasad, and A. R. Reddy, "FPGA implementation of image encryption and decryption using AES 128-bit core," in 2016 International Conference on Communication and Electronics Systems (ICCES), 2016, pp. 1–5.
- [21] R. Andriani, S. E. Wijayanti, and F. W. Wibowo, "Comparision Of AES 128, 192 And 256 Bit Algorithm For Encryption And Description File," in 2018 3rd International Conference on Information Technology, Information System and Electrical Engineering (ICITISEE), 2018, pp. 120–124.
- [22] X. Zheng and J. Jin, "Research for the application and safety of MD5 algorithm in password authentication," in 2012 9th International Conference on Fuzzy Systems and Knowledge Discovery, 2012, pp. 2216–2219.
- [23] B. Preneel and P. C. van Oorschot, "On the security of iterated message authentication codes," *IEEE Trans. Inf. Theory*, vol. 45, no. 1, pp. 188–199, Jan. 1999.
- [24] K. Jarvinen, M. Tommiska, and J. Skytta, "Hardware Implementation Analysis of the MD5 Hash Algorithm," in *Proceedings of the 38th Annual Hawaii International Conference on System Sciences*, 2005, pp. 298a–298a.
- [25] H. A. Abyaneh, M. Al-Dabbagh, H. K. Karegar, S. H. H. Sadeghi, and R. A. J. Khan, "A new optimal approach for coordination of overcurrent relays in interconnected power systems," *IEEE Trans. Power Deliv.*, vol. 18, no. 2, pp. 430–435, Apr. 2003.
- [26] S. Yin *et al.*, "Gate driver optimization to mitigate shoot-through in high-speed switching SiC half bridge module," in 2015 IEEE 11th International Conference on Power Electronics and Drive Systems, 2015, pp. 484–491.
- [27] D. Arora, S. Ravi, A. Raghunathan, and N. K. Jha, "Secure embedded processing through hardware-assisted run-time monitoring," in *Design, Automation and Test in Europe*, 2005, pp. 178-183 Vol. 1.
- [28] H. Falaghi and M.- Haghifam, "Distributed Generation Impacts on Electric Distribution Systems Reliability: Sensitivity Analysis," in EUROCON 2005 - The International Conference on "Computer as a Tool," 2005, vol. 2, pp. 1465–1468.
- [29] A. S. Tanenbaum, J. N. Herder, and H. Bos, "Can we make operating systems reliable and secure?," *Computer*, vol. 39, no. 5, pp. 44–51, May 2006.

- [30] J. N. Herder, H. Bos, B. Gras, P. Homburg, and A. S. Tanenbaum, "Fault isolation for device drivers," in 2009 IEEE/IFIP International Conference on Dependable Systems Networks, 2009, pp. 33–42.
- [31] L. Li, M. Lu, and T. Gu, "Constructing runtime models of complex software-intensive systems for analysis of failure mechanism," in 2015 First International Conference on Reliability Systems Engineering (ICRSE), 2015, pp. 1–10.
- [32] D. Hristu-Varsakelis and W. S. Levine, *Handbook of Networked and Embedded Control Systems*. Springer Science & Business Media, 2007.
- [33] B. Alpern and F. B. Schneider, "Recognizing safety and liveness," *Distrib. Comput.*, vol. 2, no. 3, pp. 117–126, Sep. 1987.
- [34] M. Raynal, *Concurrent programming algorithms, principles, and foundations*. Heidelberg; New York: Springer-Verlag, 2013.
- [35] K. Stouffer, V. Pillitteri, S. Lightman, M. Abrams, and A. Hahn, "Guide to Industrial Control Systems (ICS) Security," National Institute of Standards and Technology, NIST SP 800-82r2, Jun. 2015.
- [36] J. Daemen and V. Rijmen, *The Design of Rijndael: AES The Advanced Encryption Standard*. Springer Science & Business Media, 2013.
- [37] C. C. Erway, A. Küpçü, C. Papamanthou, and R. Tamassia, "Dynamic Provable Data Possession," *ACM Trans Inf Syst Secur*, vol. 17, no. 4, pp. 15:1–15:29, Apr. 2015.
- [38] H. Shukla, V. Singh, Y. Choi, J. Kwon, and C. Hahm, "Enhance OS security by restricting privileges of vulnerable application," in *2013 IEEE 2nd Global Conference on Consumer Electronics (GCCE)*, 2013, pp. 207–211.
- [39] Y. Li, J. M. McCune, and A. Perrig, "VIPER: Verifying the Integrity of PERipherals' Firmware," in *Proceedings of the 18th ACM Conference on Computer and Communications Security*, New York, NY, USA, 2011, pp. 3–16.
- [40] A. Ramaswamy, S. Bratus, S. W. Smith, and M. E. Locasto, "Katana: A Hot Patching Framework for ELF Executables," in 2010 International Conference on Availability, *Reliability and Security*, 2010, pp. 507–512.
- [41] Y. Huang, C. Kintala, N. Kolettis, and N. D. Fulton, "Software rejuvenation: analysis, module and applications," in *Twenty-Fifth International Symposium on Fault-Tolerant Computing. Digest of Papers*, 1995, pp. 381–390.
- [42] U. Guin, K. Huang, D. DiMase, J. M. Carulli, M. Tehranipoor, and Y. Makris, "Counterfeit Integrated Circuits: A Rising Threat in the Global Semiconductor Supply Chain," *Proc. IEEE*, vol. 102, no. 8, pp. 1207–1228, Aug. 2014.

- [43] G. Mazzanti, "Distortion limits in international standards vs. reliability of power components: Always on the safe side as to low-order voltage harmonics?," in 2012 IEEE Power and Energy Society General Meeting, 2012, pp. 1–8.
- [44] J. M. Fife, M. Scharf, S. G. Hummel, and R. W. Morris, "Field reliability analysis methods for photovoltaic inverters," in *2010 35th IEEE Photovoltaic Specialists Conference*, 2010, pp. 002767–002772.
- [45] Y. Yuan, Q. Zhu, F. Sun, Q. Wang, and T. Başar, "Resilient control of cyber-physical systems against Denial-of-Service attacks," in 2013 6th International Symposium on Resilient Control Systems (ISRCS), 2013, pp. 54–59.
- [46] Lattice Semiconductor, "MachXO2FamilyDataSheet-948089.pdf," 2016. [Online]. Available: https://www.mouser.com/datasheet/2/225/MachXO2FamilyDataSheet-948089.pdf. [Accessed: 01-Jun-2019].
- [47] Z. Ying *et al.*, "Multiplexing efficiency of high order MIMO in mobile terminal for 5G communication at 15 GHz," in 2016 International Symposium on Antennas and Propagation (ISAP), 2016, pp. 594–595.
- [48] J. Carpenter and R. Melhem, "Deterministic Multiplexing of NoC on Grid CMPs," in 2013 IEEE 21st Annual Symposium on High-Performance Interconnects, 2013, pp. 1–8.
- [49] E. Başar, Ü. Aygölü, E. Panayırcı, and H. V. Poor, "Orthogonal Frequency Division Multiplexing With Index Modulation," *IEEE Trans. Signal Process.*, vol. 61, no. 22, pp. 5536–5549, Nov. 2013.
- [50] Microchip, "1K Microwire Compatible Serial EEPROM," 2008.

# Appendix

# **Appendix A: Hardware and Software Design Details**

The following figures are imbedded PDF objects. To fully view: right click on the figure, select Acrobat Document Object  $\rightarrow$  Open. The PDF will open in a PDF reader.

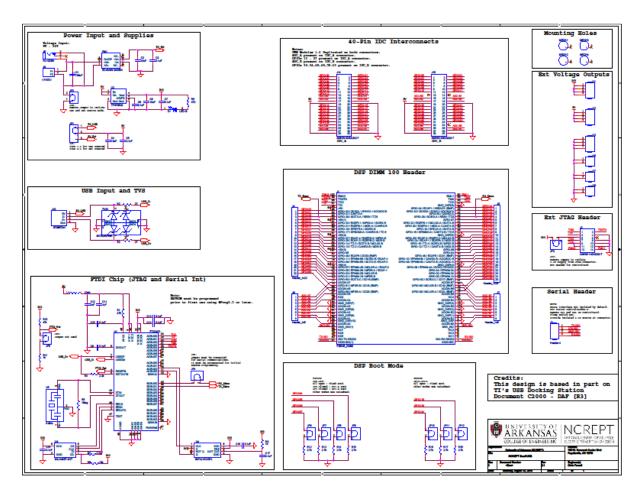


Fig. 42. Digital Signal Processor Unified Controller Board schematic

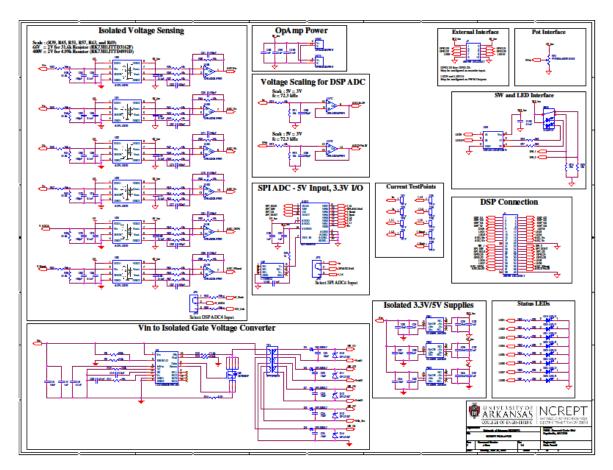


Fig. 43. Power Electronics Evaluation Unified Controller Board schematic

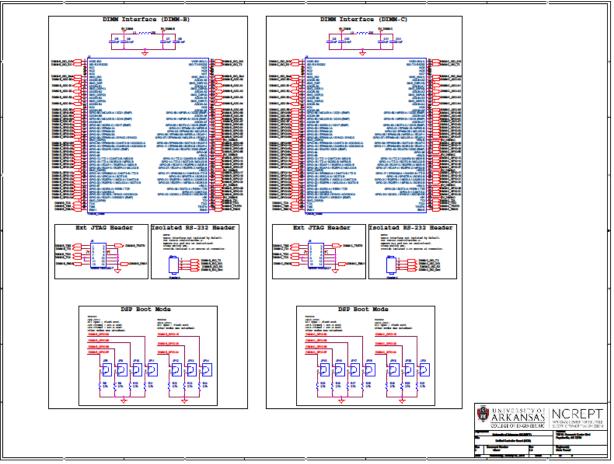


Fig. 44. Complex Programmable Logic Device Unified Controller Board schematic

		4		3			3 1	
	REV	Description	DATE	BY	1			
	A4A	Initial production Release.	11/19/2012	9C	]			
	A5	On the initial production release the processors were to be found incorrect as supplied by TI. Parts while marked AM3359 were actually AM3352. This revision uses the correct parts.	1/2/2013	GC	]	PAGE NO.	SCHEMATIC PAGE	
2		1. Deleted R29-R44 from the LCD lines.		GC		1	COVER PAGE	
11		<ol> <li>Added 47pf capacitors C158-C173 to LCD data lines to ground.</li> <li>Changed schematic revision to A5A.</li> </ol>				2	POWER MANAGEMENT	
11	A5A	4. Changed a few footprints after PCB update for above changes.				3	PROCESSOR 1 OF 3, JTAG HEADER	
		5. Added access point for the battery function of the TP865217C.	2/8/2013			4		
H		<ol> <li>Added Ferrite beads in series with LED power and SV power rail of the USB host connector. Required to pass FCO/CE testing due to noise emissions on that pin.</li> <li>Added power button to enable sleep, waksup, power down and power up features on the system.</li> </ol>				5	PROCESSOR 2 OF 3, UAB PORTS PROCESSOR 3 OF 3	
						-		
		<ol> <li>Added Modification to add 100K ohm resistor to ground to prvent crosstalk when serial cable is not plugged in.</li> </ol>				6	LED, CONFIGURATION AND BUTTON	
	A58	1. Added 100K pulldown on J1 pin 4 to prevent crosstalk when serial cable is not connected into PCB layout. 2. Charged the LED resistors to 4.75K to lower the brightness.	5/21/2013	9C	1	7	DDR3 MEMORY	
c						8	eMMC FLASH	
	ASC	<ol> <li>Charged R46, R47, R48 to 0 ohms.</li> <li>Charged R45 to 22 Ohms.</li> <li>Charge was made due to production failures on some boards due to differences in impedances.</li> </ol>	6/12/2013	oc	1	9	10/100 ETHERNET	
						10	HDMI FRAMER	
	AB	<ol> <li>Moved the enable for the VDD_3V38 regulator to VDD_3V3A rate. Change was made to reduce the delay between the rating up of the 3.3 y rate.</li> <li>Added a AND gate to the SY8_RESET in circuity. There is a small chance that on power up the rRESETOD regulan of the processor may ap high, causing the SY8_RESET is signal to go H before it should. This change reenforces the lead with the PORZD relet signal.</li> <li>Added option are on the relet or the GRN_GODD variant dignal.</li> </ol>	7/25/2013	9C		11 EXP CONN, uSD		
•								
	ABA	<ol> <li>Added optional zero ohm resistor to tie GND_OSC1 to system ground.</li> </ol>	12/13/2013		1			
		2. Changed C106 to a 1uF capacitor. 3. Changed C24 to a 2.2uF capacitor.	12/13/2013	GC				
•		4. Made R8 installed and R9 not installed.			NOTE: PCB Revision for this board is Rev B6			
	8	1.Changed the processor to the AM3358BZCZ100.	1/20/2014	90				
	c	1.Increased the eMMC from 2GB to 4GB.	3/21/2014	GC				
		his schematic is "NOT SUPPORTED" and DUES NOT constitute reference design. Only "community" support is allowed ia resources at BeagleBoard.org/discuss.						
•	P I P H T F	HERE IS NO MAREANTY FOR THIS DESIGN, TO THE EXTENT SEMITTED BY APPLICALE IAN. EXCEPT WERE OTHERWISE STATED IN MEITING THE COPYRIGHT HOLDRES AND/OR OTHER DARTIES ROVIDE THE DESIGN "AS IS" WITHOUT MAREANTY OR ANY KIND, THERE EXPRESSED OR INFLIED, INCLIDING, HUT NOT LIMITSD O, THE INFLIED WAREANTIES OF MENGAMENTABILITY AND FITNESS OR A PARTICULAR VORDORS. THE INTIME RISK AS TO THE UNLITY AND THEROFUNDANCE OF THE DESIGN IS WITH YOU. SHOULD					beagleboard.org	
	Ť	HE DESIGN PROVE DEFECTIVE, YOU ASSUME THE COST OF ALL ECESSARY SERVICING, REPAIR OR CORRECTION.					B         450-5500-001         C           Delix: Friday, March 21, 2014         Bheet 1 of 11	

Fig. 45. BeagleBone Black schematic

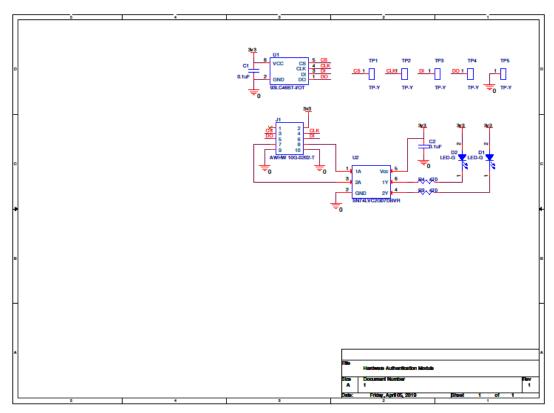
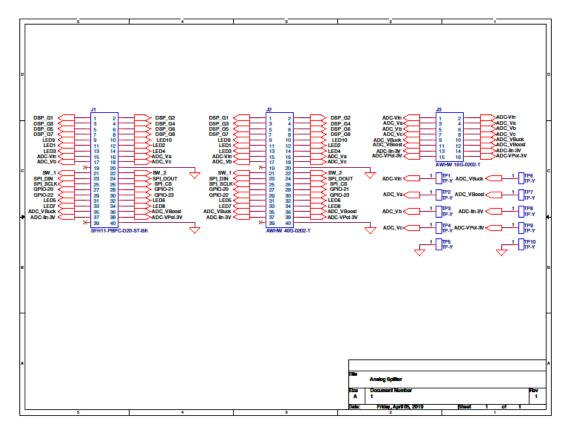


Fig. 46. Hardware Authentication Module schematic





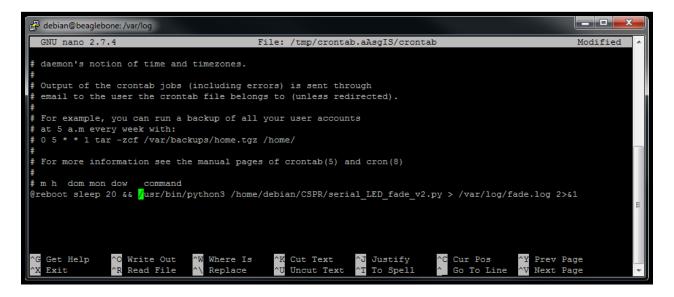


Fig. 48. Crontab configuration on BeagleBone Black to run CPLD UCB LED control script on startup

```
#!/usr/bin/env python3
 2
      import serial
 3
      import time
 4
      import struct
 5
 6
      #comm = input("Enter comm port (ex. \"COM16\")")
 7
      comm = "/dev/tty00"
 8
      ser = serial.Serial(comm, 9600, timeout=0)
 9
    [] for a in range(0,3):
10
          ser.write(b'\x71\x01\x00\x00\x01')#Enable LEDs
11
          time.sleep(0.1)
12
          ser.write(b'\x71\x01\x00\x00')#Set Blink Frequency
13
          time.sleep(0.1)
          ser.write(b'\x71\x01\x02\x2F\xFF')#Set On Time
14
15
          time.sleep(0.1)
16
17
     prefix = b' \times 71 \times 01 \times 03'
18
19
    while True:
20
    Ę
          try:
21
    Ē
               for a in range(0,65534,256):
22
                   b = struct.pack('>H',a)
23
                   packet = b''.join([prefix,b])
24
                   ser.write(packet)
25
26
               for a in range(65534,-1,-256):
    Ē
27
                   b = struct.pack('>H',a)
                   packet = b''.join([prefix,b])
28
29
                   ser.write(packet)
30
               ser.write(b'\x71\x01\x03\x00\x00')
31
               time.sleep(0.5)
    自
32
          except serial.SerialException:
33
               pass
34
```

Fig. 49. Content of LED control Python script running on the BeagleBone Black

## Appendix B: EEPROM\_WRITE\_PASSWORD

/\*

MicrowireEEPROM Example Sketch

Reads and writes a Microwire EEPROM.

Written by Timo Schneider <timos@perlplexity.org> and Joe Moquin \*/

#include <MicrowireEEPROM.h>

// Microwire needs four wires (apart from VCC/GND) DO,DI,CS,CLK // configure them here, note that DO and DI are the pins of the // EEPROM, so DI is an output of the uC, while DO is an input int CS=13; int CLK=12; int DI=7; int DO=2; // EEPROMS have different sizes. Also the number of bits per page varies. // We need to configure the page size in bits (PGS) and address bus width // in bits (ADW). The speed at which the clock is run is configured in // microseconds. //int PGS=16; int ADW=8; int SPD=200; int PGS=16; int ADW=6; int SPD=700; unsigned int password[64] = { 0xABBA,0xABED,0xBABE,0xBADE,0xBEAD,0xBEEF,0xCAFE,0xCEDE, 0xDADA,0xDEAD,0xDEAF,0xDEED,0xFACE,0xFADE,0xFEED,0xFEE0, 0xABBA,0xABED,0xBABE,0xBADE,0xBEAD,0xBEEF,0xCAFE,0xCEDE, 0xDADA,0xDEAD,0xDEAF,0xDEED,0xFACE,0xFADE,0xFEED,0xFEE0, 0xABBA,0xABED,0xBABE,0xBADE,0xBEAD,0xBEEF,0xCAFE,0xCEDE, 0xDADA,0xDEAD,0xDEAF,0xDEED,0xFACE,0xFADE,0xFEED,0xFEE0, 0xABBA,0xABED,0xBABE,0xBADE,0xBEAD,0xBEEF,0xCAFE,0xCEDE, 0xDADA,0xDEAD,0xDEAF,0xDEED,0xFACE,0xFADE,0xFEED,0xFEE0 };

```
// initialize the library
MicrowireEEPROM ME(CS, CLK, DI, DO, PGS, ADW, SPD);
```

```
void setup() {
   Serial.begin(9600);
   set_memory_map();
}
void loop() {
   for (int addr=0; addr < 64; addr++) {
     unsigned int r = ME.read(addr);
     String addr_reading = "Address " + String(addr, HEX) + "(" + String(addr, DEC) + ") DO: "
   + String(r, HEX) + " ";
     Serial.println(addr_reading);
     delay(100);
   }
</pre>
```

}

```
void set_memory_map() {
    ME.writeEnable();
    delay(10);
    for (int addr=0; addr < 64; addr++) {
        ME.write(addr,password[addr]);
        delay(10);
    }
    ME.writeDisable();
    Serial.println("Write complete.");
}</pre>
```

## Appendix C: CSPR\_V7.lpf

#LCMXO2-7000HC 4FG484C WITHIN THE UCB V1.3A **#DESIGNER: CHRIS FARNELL #AUTHOR: JOE MOOUIN #CONTACT: CFARNELL@UARK.EDU** 3/12/2019 **#DATE:** # **#CONTENTS: # BLOCK** # GPIO: IDC-A # # IDC-B IDC-C # # IDC-D # ADC: # ADC1 # ADC2 **# INTERFACE:** # BUTTONS # LEDS **DIP SWITCHES** # **#** COMMUNICATION: # SCI # CLOCKS # #TO DO: XPORT, TI-RX/TX, LAT, EXT CLK/RST **BLOCK RESETPATHS : BLOCK ASYNCPATHS**; BANK 0 VCCIO 3.3 V: BANK 1 VCCIO 2.5 V; BANK 2 VCCIO 3.3 V: BANK 3 VCCIO 3.3 V; BANK 4 VCCIO 3.3 V; BANK 5 VCCIO 3.3 V; #IDC A LOCATE COMP "A0" SITE "A21"; LOCATE COMP "A1" SITE "C19"; LOCATE COMP "A2" SITE "A20"; LOCATE COMP "A3" SITE "D18"; LOCATE COMP "A4" SITE "B19"; LOCATE COMP "A5" SITE "C18"; LOCATE COMP "A6" SITE "F17"; LOCATE COMP "A7" SITE "A18"; #LOCATE COMP "A8" SITE "D17"; #LOCATE COMP "A9" SITE "E17";

LOCATE COMP "A10" SITE "A17"; LOCATE COMP "A11" SITE "C17"; LOCATE COMP "A12" SITE "F16" LOCATE COMP "A13" SITE "E16" : LOCATE COMP "A14" SITE "D16" : LOCATE COMP "A15" SITE "B15" LOCATE COMP "A16" SITE "C16" : LOCATE COMP "A17" SITE "E15" LOCATE COMP "A18" SITE "B14" : LOCATE COMP "A19" SITE "F15" ; #LOCATE COMP "A20" SITE "C15"; #LOCATE COMP "A21" SITE "B13"; #LOCATE COMP "A22" SITE "D15" : #LOCATE COMP "A23" SITE "G15"; LOCATE COMP "A24" SITE "A13"; LOCATE COMP "A25" SITE "E14"; LOCATE COMP "A26" SITE "D14" ; LOCATE COMP "A27" SITE "B12"; **#IDC B** LOCATE COMP "B0" SITE "AA22"; LOCATE COMP "B1" SITE "T19"; LOCATE COMP "B2" SITE "Y22"; LOCATE COMP "B3" SITE "W22" : LOCATE COMP "B4" SITE "W20" ; LOCATE COMP "B5" SITE "V19" : LOCATE COMP "B6" SITE "V21" : LOCATE COMP "B7" SITE "V22"; LOCATE COMP "B8" SITE "U22" : LOCATE COMP "B9" SITE "U19" : LOCATE COMP "B10" SITE "T21" ; LOCATE COMP "B11" SITE "R19"; LOCATE COMP "B12" SITE "U20" LOCATE COMP "B13" SITE "T22" LOCATE COMP "B14" SITE "R20" LOCATE COMP "B15" SITE "R18" LOCATE COMP "B16" SITE "R21" LOCATE COMP "B17" SITE "P19" LOCATE COMP "B18" SITE "T20" LOCATE COMP "B19" SITE "R22" LOCATE COMP "B20" SITE "P20"; LOCATE COMP "B21" SITE "P18"; #LOCATE COMP "B22" SITE "P21" : LOCATE COMP "Usr RX" SITE "P21"; #LOCATE COMP "B23" SITE "N17"; LOCATE COMP "Usr\_TX" SITE "N17"; LOCATE COMP "B24" SITE "N16";

LOCATE COMP "B26" SITE "N20"; LOCATE COMP "B27" SITE "M18"; **#IDC-C** LOCATE COMP "C0" SITE "Y14" : LOCATE COMP "C1" SITE "AB15"; LOCATE COMP "C2" SITE "W12" : LOCATE COMP "C3" SITE "V12"; LOCATE COMP "C4" SITE "Y12"; LOCATE COMP "C5" SITE "V13"; LOCATE COMP "C6" SITE "AA15" ; LOCATE COMP "C7" SITE "Y15"; #LOCATE COMP "C8" SITE "AB16" ; #LOCATE COMP "C9" SITE "AA16"; LOCATE COMP "C10" SITE "T13"; LOCATE COMP "C11" SITE "U13"; LOCATE COMP "C12" SITE "Y16" ; LOCATE COMP "C13" SITE "AB17"; LOCATE COMP "C14" SITE "W14"; LOCATE COMP "C15" SITE "V14"; LOCATE COMP "C16" SITE "Y17"; LOCATE COMP "C17" SITE "AB18"; LOCATE COMP "C18" SITE "W15" : LOCATE COMP "C19" SITE "V15"; #LOCATE COMP "C20" SITE "W16"; #LOCATE COMP "C21" SITE "W17"; #LOCATE COMP "C22" SITE "Y18"; #LOCATE COMP "C23" SITE "AA19"; LOCATE COMP "C24" SITE "AB20"; LOCATE COMP "C25" SITE "AB21" ; LOCATE COMP "C26" SITE "V16"; LOCATE COMP "C27" SITE "U15"; **#IDC D** LOCATE COMP "D0" SITE "C3" : LOCATE COMP "D1" SITE "C2"; LOCATE COMP "D2" SITE "F6" : LOCATE COMP "D3" SITE "F5" : LOCATE COMP "D4" SITE "E4" ; LOCATE COMP "D5" SITE "D3" LOCATE COMP "D6" SITE "G6" : LOCATE COMP "D7" SITE "H7"; LOCATE COMP "D8" SITE "B1" : LOCATE COMP "D9" SITE "C1" : LOCATE COMP "D10" SITE "H6" : LOCATE COMP "D11" SITE "G5"; LOCATE COMP "D12" SITE "E2";

LOCATE COMP "B25" SITE "N21";

71

LOCATE COMP "D13" SITE "D1"; LOCATE COMP "D14" SITE "F4"; LOCATE COMP "D15" SITE "G4" : LOCATE COMP "D16" SITE "F1"; LOCATE COMP "D17" SITE "G3"; LOCATE COMP "D18" SITE "J5"; LOCATE COMP "D19" SITE "J4" ; LOCATE COMP "D20" SITE "G2"; LOCATE COMP "D21" SITE "G1"; LOCATE COMP "D22" SITE "K6" : LOCATE COMP "D23" SITE "K7" LOCATE COMP "D24" SITE "H3" : LOCATE COMP "D25" SITE "H2" ; LOCATE COMP "D26" SITE "K5" ; LOCATE COMP "D27" SITE "L3" : **#BUTTONS SW**[1:4] ACTIVE LOW LOCATE COMP "BTN[1]" SITE "G13"; LOCATE COMP "BTN[2]" SITE "F13"; LOCATE COMP "BTN[3]" SITE "A12"; LOCATE COMP "BTN[4]" SITE "C13"; DEFINE PORT GROUP "BTN" "BTN[1]" "BTN[2]" "BTN[3]" "BTN[4]"; IOBUF GROUP "BTN" IO\_TYPE=LVCMOS33; **#LEDS ACTIVE LOW** LOCATE COMP "LED[1]" SITE "R17"; LOCATE COMP "LED[2]" SITE "T18"; LOCATE COMP "LED[3]" SITE "R16" LOCATE COMP "LED[4]" SITE "T17"; LOCATE COMP "LED[5]" SITE "Y21"; LOCATE COMP "LED[6]" SITE "Y20" LOCATE COMP "LED[7]" SITE "U18"; LOCATE COMP "LED[8]" SITE "U17"; DEFINE PORT GROUP "LED" "LED[1]" "LED[2]" "LED[3]" "LED[4]" "LED[5]" "LED[6]" "LED[7]" "LED[8]"; IOBUF GROUP "LED" IO TYPE=LVCMOS25 PULLMODE=DOWN DRIVE=8 SLEWRATE=SLOW : **#SCI** #LOCATE COMP "SCI TX" SITE "W1";

#LOCATE COMP "SCI RX" SITE "V2"; **#CLOCKS** #LOCATE COMP "XTAL\_CLK" SITE "V3"; #FREQUENCY NET "CLK" 53.200000 MHZ; IOBUF PORT "A0" IO TYPE=LVCMOS33 : IOBUF PORT "A1" IO\_TYPE=LVCMOS33; IOBUF PORT "A2" IO\_TYPE=LVCMOS33 ; IOBUF PORT "A3" IO\_TYPE=LVCMOS33; IOBUF PORT "A4" IO\_TYPE=LVCMOS33; IOBUF PORT "A5" IO\_TYPE=LVCMOS33 ; IOBUF PORT "A6" IO TYPE=LVCMOS33; IOBUF PORT "A7" IO\_TYPE=LVCMOS33; IOBUF PORT "A10" IO TYPE=LVCMOS33 : IOBUF PORT "A11" IO\_TYPE=LVCMOS33 : IOBUF PORT "A12" IO\_TYPE=LVCMOS33; IOBUF PORT "A13" IO TYPE=LVCMOS33; IOBUF PORT "A14" IO TYPE=LVCMOS33 ; IOBUF PORT "A15" IO\_TYPE=LVCMOS33; IOBUF PORT "A16" IO\_TYPE=LVCMOS33; IOBUF PORT "A17" IO\_TYPE=LVCMOS33; IOBUF PORT "A18" IO TYPE=LVCMOS33; IOBUF PORT "A19" IO TYPE=LVCMOS33 ; IOBUF PORT "A24" IO TYPE=LVCMOS33 : IOBUF PORT "A25" IO TYPE=LVCMOS33; IOBUF PORT "A26" IO\_TYPE=LVCMOS33; IOBUF PORT "A27" IO TYPE=LVCMOS33: IOBUF PORT "C0" IO\_TYPE=LVCMOS33; IOBUF PORT "C1" IO TYPE=LVCMOS33; IOBUF PORT "C2" IO\_TYPE=LVCMOS33; IOBUF PORT "C3" IO TYPE=LVCMOS33: IOBUF PORT "C4" IO\_TYPE=LVCMOS33 : IOBUF PORT "C5" IO\_TYPE=LVCMOS33 : IOBUF PORT "C6" IO TYPE=LVCMOS33; IOBUF PORT "C7" IO TYPE=LVCMOS33; IOBUF PORT "C10" IO\_TYPE=LVCMOS33; IOBUF PORT "C11" IO TYPE=LVCMOS33: IOBUF PORT "C12" IO\_TYPE=LVCMOS33; IOBUF PORT "C13" IO\_TYPE=LVCMOS33 : IOBUF PORT "C14" IO TYPE=LVCMOS33; IOBUF PORT "C15" IO\_TYPE=LVCMOS33; IOBUF PORT "C16" IO\_TYPE=LVCMOS33; IOBUF PORT "C17" IO TYPE=LVCMOS33 ; IOBUF PORT "C18" IO TYPE=LVCMOS33; IOBUF PORT "C19" IO TYPE=LVCMOS33; IOBUF PORT "C24" IO\_TYPE=LVCMOS33; IOBUF PORT "C25" IO TYPE=LVCMOS33;

IOBUF PORT "C26" IO\_TYPE=LVCMOS33; IOBUF PORT "C27" IO\_TYPE=LVCMOS33; IOBUF PORT "D23" IO\_TYPE=LVCMOS33; IOBUF PORT "D0" IO TYPE=LVCMOS33; IOBUF PORT "D1" IO\_TYPE=LVCMOS33; IOBUF PORT "D2" IO\_TYPE=LVCMOS33; IOBUF PORT "D3" IO TYPE=LVCMOS33 ; IOBUF PORT "D4" IO\_TYPE=LVCMOS33; IOBUF PORT "D5" IO TYPE=LVCMOS33; IOBUF PORT "D6" IO\_TYPE=LVCMOS33 ; IOBUF PORT "D7" IO TYPE=LVCMOS33; IOBUF PORT "D8" IO\_TYPE=LVCMOS33; IOBUF PORT "D9" IO TYPE=LVCMOS33 ; IOBUF PORT "D10" IO\_TYPE=LVCMOS33; IOBUF PORT "D11" IO\_TYPE=LVCMOS33; IOBUF PORT "D12" IO TYPE=LVCMOS33; IOBUF PORT "D13" IO\_TYPE=LVCMOS33; IOBUF PORT "D14" IO\_TYPE=LVCMOS33; IOBUF PORT "D15" IO\_TYPE=LVCMOS33; IOBUF PORT "D16" IO\_TYPE=LVCMOS33; IOBUF PORT "D17" IO TYPE=LVCMOS33; IOBUF PORT "D18" IO TYPE=LVCMOS33 ; IOBUF PORT "D19" IO\_TYPE=LVCMOS33; IOBUF PORT "D20" IO TYPE=LVCMOS33; IOBUF PORT "D21" IO\_TYPE=LVCMOS33; IOBUF PORT "D22" IO TYPE=LVCMOS33; IOBUF PORT "D24" IO\_TYPE=LVCMOS33 : IOBUF PORT "D25" IO TYPE=LVCMOS33; IOBUF PORT "D26" IO TYPE=LVCMOS33; IOBUF PORT "D27" IO\_TYPE=LVCMOS33;

## Appendix D: hardware\_protections.vhd

LIBRARY ieee; USE ieee.std\_logic\_1164.ALL;

ENTITY censor IS PORT ( pwm\_i : IN STD\_LOGIC\_VECTOR(1 DOWNTO 0); pwm\_o : OUT STD\_LOGIC\_VECTOR(1 DOWNTO 0) ); END censor;

ARCHITECTURE behavior OF censor IS BEGIN WITH pwm\_i SELECT pwm\_o <= "01" WHEN "01", "10" WHEN "10", "00" WHEN OTHERS;

END behavior;

## Appendix E: CSPR\_MODULES.vhdl

\_\_\_\_\_ -- Company: University of Arkansas (NCREPT) -- Engineer: Chris Farnell, Joe Moquin -- Create Date: March 12, 2019 -- Design Name: Generic Components -- Module Name: Various -- Project Name: Cybersecure Power Router -- Target Devices: LCMXO2-7000HC-4FG484C (MachXO2 Eval Board) \_\_\_\_\_ --##################################Generic -----Bus Interface-----LIBRARY IEEE: USE IEEE.std\_logic\_1164.ALL; USE ieee.std\_logic\_unsigned.ALL; USE ieee.numeric\_std.ALL; **ENTITY Bus** Int IS **GENERIC** ( CONSTANT DATA WIDTH : INTEGER := 16; CONSTANT Address\_WIDTH : INTEGER := 16 ); PORT ( clk : IN std logic; rst : IN std\_logic; DataIn : IN std\_logic\_vector(DATA\_WIDTH - 1 DOWNTO 0); DataOut : OUT std\_logic\_vector(DATA\_WIDTH - 1 DOWNTO 0); AddrIn : IN std\_logic\_vector(Address\_WIDTH - 1 DOWNTO 0); WE : IN std logic; RE : IN std logic; Busy : OUT std\_logic; Data : INOUT std logic vector(DATA WIDTH - 1 DOWNTO 0); Addr : OUT std\_logic\_vector(Address\_WIDTH - 1 DOWNTO 0); Xrqst : OUT std logic; XDat : IN std logic; YDat : OUT std logic; BusRqst : OUT std\_logic; BusCtrl : IN std\_logic ); END: ARCHITECTURE behavior OF Bus\_Int IS TYPE state\_type IS (S0, S1, S2, S3, S4, S5, S6, S7, S8, S9, S10);

```
SIGNAL CS, NS : state_type;
      SIGNAL AddrIn_reg_0 : std_logic_vector(DATA_WIDTH - 1 DOWNTO 0) :=
(OTHERS => '0');
      SIGNAL DataIn_reg_o : std_logic_vector(DATA_WIDTH - 1 DOWNTO 0) :=
(OTHERS => '0');
      SIGNAL LD_AddrIn, LD_DataIn, LD_Data : std_logic := '0';
BEGIN
      ----Registers
      Reg_Proc : PROCESS
       BEGIN
             WAIT UNTIL clk'event AND clk = '1';
             IF rst = '0' THEN
                    AddrIn_reg_o \leq (OTHERS => '0');
                    DataIn_reg_o \leq (OTHERS = '0');
                    DataOut \leq (OTHERS = '0');
             ELSE
                    IF (LD_AddrIn = '1') THEN
                           AddrIn_reg_o <= AddrIn;
                    END IF; --Register for reading input address
                    IF (LD_DataIn = '1') THEN
                           DataIn reg o <= DataIn;
                    END IF; --Register for reading input address
                    IF (LD_Data = '1') THEN
                           DataOut <= Data;
                    END IF; --Register for reading input address
             END IF;
      END PROCESS;
       ----End Registers
      ----Next State Logic Bus Interface
      NS Bus_Int : PROCESS (CS, WE, RE, XDat, BusCtrl, AddrIn_reg_o, DataIn_reg_o)
      BEGIN
             ----Default States to remove latches
             Busy <= '1';
             Data <= (OTHERS => 'Z');
             Addr \leq (OTHERS = 'Z');
             XRqst \leq Z';
             YDat <= 'Z';
             BusRqst <= '0';
             NS \leq S0;
             LD_AddrIn \leq '0';
             LD_DataIn \le '0';
             LD Data \leq 0';
             CASE CS IS
                    WHEN S0 = -- Waits until a read or write request is initiated.
                           IF (RE = '1') THEN
                                  NS <= S1;
```

ELSIF (WE = '1') THEN NS <= S3; ELSE  $NS \leq S0;$ END IF: Busy <= '0'; LD\_AddrIn <= '1'; -- Loads the Input Address LD\_DataIn <= '1'; -- Loads the Input Data --Begin Read Process WHEN S1 => -- Request Control of the Bus and wait. IF (BusCtrl = '1') THEN  $NS \le S2;$ ELSE  $NS \le S1;$ END IF; BusRqst <= '1'; WHEN S2 => -- Bus Control granted. Request data. IF (Xdat = '0') THEN -- Active High  $NS \le S2;$ ELSE  $NS \leq S0;$ END IF; Addr <= AddrIn\_reg\_o; XRqst <= '1'; --Active High--Active Low because of pull-ups for internal tristate LD Data <= '1'; --End Read Process --Begin Write Process WHEN S3 => -- Request Control of the Bus and wait. IF (BusCtrl = '1') THEN NS <= S4; ELSE NS <= S3; END IF: BusRqst <= '1'; WHEN S4 => -- Bus Control granted. Write data. Addr <= AddrIn reg o; Data <= DataIn\_reg\_o; YDat <= '1'; --Active High--Active Low because of pull-ups for internal tristate NS  $\leq$  S0; --End Write Process WHEN OTHERS =>  $NS \leq S0;$ 

```
END CASE;
     END PROCESS;
      ----End Next State Logic for Bus Interface
      ----State Sync
      sync_States : PROCESS
     BEGIN
           WAIT UNTIL clk'event AND clk = '1';
           IF rst = '0' THEN
                 CS \leq S0;
           ELSE
                 CS \leq NS;
           END IF:
     END PROCESS:
     ----End State Sync
END behavior:
-----End Bus Interface-----
------Generic FIFO-----
LIBRARY IEEE;
USE IEEE.std_logic_1164.ALL;
USE ieee.std_logic_unsigned.ALL;
USE ieee.numeric std.ALL;
ENTITY STD FIFO IS
     GENERIC (
           DATA WIDTH : INTEGER := 8; -- Width of FIFO
           FIFO_DEPTH : INTEGER := 512; --Depth of FIFO
           FIFO ADDR LEN: INTEGER := 9 -- Required number of bits to represent
FIFO_Depth
     );
     PORT (
           CLK : IN STD_LOGIC; -- Clock input
           RST : IN STD LOGIC; -- Active low reset
           WriteEn : IN STD_LOGIC; -- Write enable signal
           DataIn : IN STD_LOGIC_VECTOR (DATA_WIDTH - 1 DOWNTO 0); -- Data
input bus
           ReadEn : IN STD_LOGIC; -- Read enable signal
           DataOut : OUT STD LOGIC VECTOR (DATA WIDTH - 1 DOWNTO 0); --
Data output bus
           Empty : OUT STD LOGIC; -- FIFO empty flag
           Full : OUT STD LOGIC -- FIFO full flag
     );
END STD_FIFO;
ARCHITECTURE Behavioral OF STD FIFO IS
      TYPE FIFO_Memory IS ARRAY (0 TO FIFO_DEPTH - 1) OF STD_LOGIC_VECTOR
(DATA WIDTH - 1 DOWNTO 0);
     SIGNAL Memory : FIFO_Memory;
     SIGNAL Head : STD LOGIC VECTOR (FIFO ADDR LEN - 1 DOWNTO 0);
```

```
SIGNAL Tail : STD_LOGIC_VECTOR (FIFO_ADDR_LEN - 1 DOWNTO 0);
      SIGNAL Looped : BOOLEAN;
BEGIN
      -- Memory Pointer Process
      fifo_proc : PROCESS (CLK)
      BEGIN
             IF rising_edge(CLK) THEN
                    IF RST = '0' THEN
                           Head \leq (OTHERS => '0');
                           Tail \leq (OTHERS = '0');
                          Looped <= false;
                          Full <= '0';
                           Empty <= '1';
                    ELSE
                          IF (ReadEn = '1') THEN
                                 IF ((Looped = true) OR (Head /= Tail)) THEN
                                        -- Update data output
                                        DataOut <= Memory(CONV_INTEGER(Tail));</pre>
                                        -- Update Tail pointer as needed
                                        IF (Tail = FIFO_DEPTH - 1) THEN
                                               Tail \leq (OTHERS = '0');
                                               Looped <= false;
                                        ELSE
                                               Tail \ll Tail + 1;
                                        END IF;
                                 END IF;
                           END IF;
                           IF (WriteEn = '1') THEN
                                 IF ((Looped = false) OR (Head /= Tail)) THEN
                                        -- Write Data to Memory
                                        Memory(CONV_INTEGER(Head)) <= DataIn;
                                        -- Increment Head pointer as needed
                                        IF (Head = FIFO_DEPTH - 1) THEN
                                               Head \leq (OTHERS => '0');
                                               Looped <= true;
                                        ELSE
                                               Head \leq Head + 1;
                                        END IF;
                                 END IF;
                           END IF:
                           -- Update Empty and Full flags
                          IF (Head = Tail) THEN
                                 IF Looped THEN
                                        Full <= '1';
                                 ELSE
                                        Empty <= '1';
```

END IF; ELSE Empty <= '0': Full <= '0'; END IF: END IF; END IF: END PROCESS; END Behavioral: -----End Generic FIFO----------16-Bit PWM with Phase shift-----LIBRARY IEEE; USE IEEE.STD\_LOGIC\_1164.ALL; USE IEEE.STD\_LOGIC\_ARITH.ALL; USE IEEE.STD\_LOGIC\_UNSIGNED.ALL; USE ieee.numeric std.ALL; ENTITY PWM\_16b IS GENERIC ( Freq in : INTEGER := 25000000; --Clk (25 MHz) Max\_PWM : INTEGER := 65535; --PWM Resolution (2^16-1) Freq Sw : INTEGER := 6104); --PWM Switching Frequency (Should be derived from Main Clock) (25e6/2^12) PORT ( clk: IN STD LOGIC; rst : IN STD\_LOGIC; DC: IN STD LOGIC VECTOR (15 DOWNTO 0); Phase : IN STD\_LOGIC\_VECTOR (15 DOWNTO 0); En : IN STD LOGIC: PWM\_Out : OUT STD\_LOGIC); END PWM 16b; ARCHITECTURE Behavioral OF PWM 16b IS --Constants CONSTANT Max Period : INTEGER := (Freq in/Freq Sw) - 1; CONSTANT PWM Step Inv : INTEGER := Max PWM/Max Period; --Clk cycle step size for Duty cycle CONSTANT PWM Max : INTEGER := Max PWM; CONSTANT PWM\_Min : INTEGER := PWM\_Step\_Inv; --Signals SIGNAL PWM Count, DC Read, Phase Read : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); BEGIN DC\_Update : PROCESS **BEGIN** WAIT UNTIL clk'event AND clk = '1';IF rst = '0' THEN DC Read  $\leq$  (OTHERS = '0');

Phase\_Read  $\leq$  (OTHERS = '0'); ELSE -- For 1.526 kHz --DC Read(15 downto 14)  $\leq$  (others = '0');  $DC_Read(13 \text{ downto } 0) \le DC(15 \text{ downto } 2);$ --shift 2 places for divide by 4 (PWM\_Step\_Inv) Phase\_Read(15 downto 14) <= (others => '0'); Phase\_Read(13 downto 0)<= Phase(15 downto 2); --shift 2 places for divide by 4 (PWM Step Inv) -- For 3.052 kHz DC Read(15 downto 13)  $\leq$  (others = '0'); -- $DC_Read(12 \text{ downto } 0) \le DC(15 \text{ downto } 3);$ --shift 3 places for divide by 8 Phase\_Read(15 downto 13) <= (others => '0'); Phase\_Read(12 downto 0)<= Phase(15 downto 3); --shift 3 places for divide by 8 -- For 6.104 kHz  $DC_Read(15 \text{ DOWNTO } 12) \le (OTHERS \implies '0');$ DC Read(11 DOWNTO 0)  $\leq$  DC(15 DOWNTO 4); --shift 4 places for divide by 16 (PWM\_Step\_Inv) Phase Read(15 DOWNTO 12)  $\leq$  (OTHERS = '0'); Phase Read(11 DOWNTO 0) <= Phase(15 DOWNTO 4); --shift 4 places for divide by 16 (PWM\_Step\_Inv) END IF; END PROCESS; Count Update : PROCESS BEGIN WAIT UNTIL clk'event AND clk = '1';IF rst = '0' THEN PWM Count  $\leq$  (OTHERS = '0'): ELSIF (PWM Count <= (Max Period + Phase Read)) THEN PWM\_Count <= PWM\_Count + 1; ELSE PWM Count <= Phase Read; END IF: END PROCESS; PWM\_Update : PROCESS BEGIN WAIT UNTIL clk'event AND clk = '1'; IF rst = '0' THEN PWM\_Out <= '0'; ELSIF en = 0' THEN PWM Out  $\leq 0'$ ; ELSIF ((PWM Count <= (DC Read + Phase Read)) AND ((PWM Count) > (Phase\_Read))) THEN PWM Out  $\leq 1'$ ;

```
ELSE
                  PWM_Out <= '0';
            END IF:
      END PROCESS;
END Behavioral:
-----End 16-Bit PWM with Phase shift-----
-----16-Bit Shift Register(Parallel-to-Serial)------
LIBRARY IEEE;
USE IEEE.std_logic_1164.ALL;
USE ieee.std_logic_unsigned.ALL;
USE ieee.numeric_std.ALL;
ENTITY Sreg_PS_16 IS
      PORT (
            ld_D, sh_D, rst, clk : IN std_logic;
            Data_In : IN STD_LOGIC_VECTOR(15 DOWNTO 0);
            Data Out : OUT std logic);
END:
ARCHITECTURE BEHAVIOR OF Sreg_PS_16 IS
      SIGNAL temp : STD_LOGIC_VECTOR(15 DOWNTO 0);
BEGIN
      --Data Out \leq temp(15);
      Counter_behav : PROCESS
      BEGIN
            WAIT UNTIL clk'event AND clk = '1';
            IF rst = '0' THEN
                  temp \leq (\text{OTHERS} = 0');
                  Data_Out \leq 0';
            ELSIF ld D = '1' THEN
                  temp <= Data_In;
                  Data_Out \le temp(15);
            ELSIF sh D = '1' THEN
                  temp \leq temp(14 DOWNTO 0) & '0';
                  Data_Out \leq temp(15);
            ELSE
                  Data_Out \le temp(15);
            END IF:
      END PROCESS;
END BEHAVIOR:
-----End of 16-Bit Shift Register(Parallel-to-Serial)------
-----16-Bit Shift Register(Serial-to-Parallel)------
LIBRARY IEEE;
USE IEEE.std_logic_1164.ALL;
USE ieee.std_logic_unsigned.ALL;
USE ieee.numeric std.ALL;
ENTITY Sreg_SP_16 IS
      PORT (
```

```
ld_D, rst, clk : IN std_logic;
            Data_In : IN std_logic;
            Data_Out : OUT STD_LOGIC_VECTOR(15 DOWNTO 0));
END;
ARCHITECTURE BEHAVIOR OF Sreg_SP_16 IS
      SIGNAL temp : STD_LOGIC_VECTOR(15 DOWNTO 0);
BEGIN
      Counter_behav : PROCESS
      BEGIN
            WAIT UNTIL clk'event AND clk = '1';
            IF rst = '0' THEN
                  temp \leq (OTHERS = '0');
                  Data_Out \leq (OTHERS = '0');
            ELSIF ld D = '1' THEN
                  temp <= temp(14 DOWNTO 0) & Data_In;
                  --\text{temp}(0) \leq \text{Data In;}
                  Data_Out <= temp;</pre>
            ELSE
                  Data_Out <= temp;
            END IF;
      END PROCESS;
END BEHAVIOR;
-----End of 16-Bit Shift Register(Parallel-to-Serial)------
----- Standard Counter-----
LIBRARY IEEE;
USE IEEE.std logic 1164.ALL;
USE ieee.std_logic_unsigned.ALL;
USE ieee.numeric std.ALL;
ENTITY Std_Counter IS
      GENERIC (
            Width : INTEGER := 8 --width of counter
      );
      PORT (
            INC, rst, clk : IN std logic;
            Count : OUT STD_LOGIC_VECTOR(Width - 1 DOWNTO 0));
END;
ARCHITECTURE BEHAVIOR OF Std_Counter IS
      SIGNAL temp : STD_LOGIC_VECTOR(Width - 1 DOWNTO 0);
BEGIN
      Counter_behav : PROCESS
      BEGIN
            WAIT UNTIL clk'event AND clk = '1';
            IF rst = '0' THEN
                  temp \leq (OTHERS = '0');
            ELSIF INC = '1' THEN
                  temp \leq temp + 1;
```

ELSE

NULL;

END IF; END PROCESS;

Count <= temp;

END BEHAVIOR;

END DEFIA VIOK,

---

- -- Company: University of Arkansas (NCREPT)
- -- Engineer: Chris Farnell
- --
- -- Create Date: 3Dec2018
- -- Design Name: Bus Interface Common
- -- Module Name: Bus\_Interface\_Common
- -- Project Name: Bus Interface Example
- -- Target Devices: LCMXO2-7000HE-4TG144C (MachXO2 Eval Board)
- -- Tool versions: Lattice Diamond\_x64 Build 3.10.2.115.1

-- Description:

-- This Package was created to allow for Memory Mapping as well as the declaration of various needed constants.

- ---- Register and Memory Map Information:
- -- This section describes the Memory Map used in this project.
- -- This design contains a SPRAM Module which is 16 bits wide and 1024 entries deep.
- -- Register addresses are from X"0000" to X"03FF".
- -- All registers are 16-bits wide.
- -- The SPRAM Module is located in the Bus\_Master portion of the code.
- -- This RAM Module may be accessed externally using either Serial Port interface.
- -- Reserved for future use.
- -- X"0200" X"03FF"
- -- LED Configuration Registers-
- -- Range is X"0100" X"010A"

-- Register Map is found as constants in Bus\_Interface\_Common and shared with all submodules of this program.

-- Revisions:--

- --
- -- Revision 0.01 -

-- File Created; Basic\Classical Operation Implemented

--

-- Additional Comments:

- --
- --

-----

LIBRARY IEEE: USE IEEE.std\_logic\_1164.ALL; USE ieee.std\_logic\_unsigned.ALL; USE ieee.numeric std.ALL; **ENTITY Bus Master IS** PORT ( clk : IN STD\_LOGIC; rst : IN STD\_LOGIC; Data : INOUT STD\_LOGIC\_VECTOR (15 DOWNTO 0); Addr : IN STD\_LOGIC\_VECTOR (15 DOWNTO 0); Xrqst : IN STD LOGIC; XDat : OUT STD\_LOGIC; YDat : IN STD LOGIC; BusRqst : IN STD\_LOGIC\_VECTOR (9 DOWNTO 0); BusCtrl : OUT STD\_LOGIC\_VECTOR (9 DOWNTO 0)); END Bus Master; **ARCHITECTURE Behavioral OF Bus Master IS** TYPE state\_type IS (S0, S1, S2, S3, S4, S5, S6, S7, S8, S9, S10, S11); SIGNAL CS, NS : state\_type; --Signals for Mem1 SIGNAL Mem1 wea : STD LOGIC := '0'; SIGNAL Mem1 rst : STD LOGIC := '0'; SIGNAL Mem1\_addra : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Mem1 dina, Mem1 douta : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0');SIGNAL clk en : STD LOGIC := '1'; --Signals for Registers SIGNAL LD Addr, LD Data, LD BusCtrl : Std Logic := '0'; SIGNAL BusCtrl\_Temp : STD\_LOGIC\_VECTOR(9 DOWNTO 0) := (OTHERS => '0'); --declare SPRAM COMPONENT SPRAM PORT ( Clock : IN std logic; ClockEn : IN std logic; Reset : IN std\_logic; WE : IN std logic; Address : IN std\_logic\_vector(9 DOWNTO 0); Data : IN std logic vector(15 DOWNTO 0); Q : OUT std logic vector(15 DOWNTO 0) ): END COMPONENT; BEGIN --Instantiate SPRAM 16bx1024 Mem1 : SPRAM PORT MAP( Clock => clk,ClockEn => clk en,

```
Reset => Mem1_rst,
       WE \Rightarrow Mem1_wea,
       Address => Mem1_addra(9 DOWNTO 0),
      Data => Mem1_dina,
       Q => Mem1_douta
);
----Registers
Reg_Proc : PROCESS
BEGIN
       WAIT UNTIL clk'event AND clk = '1';
      IF rst = '0' THEN
              Mem1_addra \le (OTHERS \implies '0');
              Mem1_dina \ll (OTHERS \implies '0');
              BusCtrl <= (OTHERS => '0');
      ELSE
              IF (LD Addr = '1') THEN
                     Mem1_addra <= Addr;
              END IF; --Register for reading input address
              IF (LD_Data = '1') THEN
                     Mem1_dina <= Data;
              END IF; --Register for writing input data
              IF (LD_BusCtrl = '1') THEN
                    BusCtrl <= BusCtrl_Temp;</pre>
              END IF;
      END IF;
END PROCESS;
----End Registers
----Next State Logic Bus Control
NS_Bus_Ctrl : PROCESS (CS, BusRqst, XRqst, YDat, Mem1_douta)
BEGIN
       ----Default States to remove latches
      Data \leq (OTHERS => 'Z');
      XDat <= '0';
      BusCtrl Temp \leq (OTHERS = '0');
      LD_BusCtrl <= '0';
      NS \leq S0;
      Mem1_wea \leq 0';
      LD Addr \leq 0';
      LD_Data <= '0';
      clk_en <= '1';
      CASE CS IS
              WHEN S0 = -- Waits until a request is made.
                     IF (BusRqst > 0) THEN
                           NS <= S1;
                     ELSE
                           NS \leq S0;
```

END IF; WHEN S1 => -- Grant Control of the Bus (Priority Encoder) IF (BusRqst(0) = '1') THEN BusCtrl\_Temp(0)  $\leq 1'$ ; ELSIF (BusRqst(1) = '1') THEN BusCtrl\_Temp(1)  $\leq 1'$ ; ELSIF (BusRqst(2) = '1') THEN BusCtrl\_Temp(2)  $\leq 1'$ ; ELSIF (BusRqst(3) = '1') THEN BusCtrl\_Temp(3)  $\leq 1'$ ; ELSIF (BusRqst(4) = '1') THEN BusCtrl\_Temp(4)  $\leq 1'$ ; ELSIF (BusRqst(5) = '1') THEN BusCtrl\_Temp(5)  $\leq 1'$ ; ELSIF (BusRqst(6) = '1') THEN BusCtrl Temp(6)  $\leq 1'$ ; ELSIF (BusRqst(7) = '1') THEN BusCtrl\_Temp(7)  $\leq 1'$ ; ELSIF (BusRqst(8) = '1') THEN BusCtrl\_Temp(8)  $\leq 1'$ ; ELSIF (BusRqst(9) = '1') THEN BusCtrl\_Temp(9)  $\leq 1'$ ; END IF: LD BusCtrl  $\leq 1'$ ;  $NS \le S2;$ WHEN S2 => -- Bus Control granted. Wait until Read or Write Request. IF (XRqst = '1') THEN -- Active High-- Active Low because of pull-ups for internal tristate  $NS \le S3;$ ELSIF (YDat = '1') THEN -- Active High-- Active Low because of pull-ups for internal tristate  $NS \le S5;$ ELSE  $NS \le S2;$ END IF: LD Addr <= '1'; LD\_Data <= '1'; WHEN S3 => --(Read Operation) Send Data NS  $\leq$  S4; WHEN S4 => --(Read Operation) Send Data data <= Mem1\_douta; Xdat <= '1'; --Active High NS <= S6; WHEN S5 = --(Write Operation) Receive Data Mem1\_wea  $\leq 1'$ ; NS <= S6;

```
WHEN S6 =>
                      LD_BusCtrl <= '1';
                      NS <= S0;
                 WHEN OTHERS =>
                      NS \leq S0;
           END CASE;
     END PROCESS:
     ----End Next State Logic for Bus Interface
     ----State Sync
     sync_States : PROCESS
     BEGIN
           WAIT UNTIL clk'event AND clk = '1';
           IF rst = '0' THEN
                 Mem1 rst <= '1'; --reset Memory
                 CS \leq S0;
           ELSE
                 Mem1_rst <= '0';
                 CS \leq NS;
           END IF;
     END PROCESS;
     ----End State Sync
END Behavioral:
#######
LIBRARY IEEE:
USE IEEE.std_logic_1164.ALL;
USE IEEE.std logic unsigned.ALL;
USE IEEE.numeric_std.ALL;
USE IEEE.std_logic_arith.ALL;
ENTITY RS232 Usr Int IS
     GENERIC (
           Baud : INTEGER := 9600; --9,600 bps
           clk in : INTEGER := 24930000); --24.93MHz
     PORT (
           clk: IN STD LOGIC;
           rst : IN STD_LOGIC;
           rs232 rcv : IN STD LOGIC;
           rs232 xmt : OUT STD LOGIC;
           Data : INOUT std_logic_vector(15 DOWNTO 0);
           Addr : OUT std_logic_vector(15 DOWNTO 0);
           Xrqst : OUT std_logic;
           XDat : IN std logic;
           YDat : OUT std logic;
           BusRqst : OUT std_logic;
```

```
BusCtrl : IN std_logic
```

);

END RS232\_Usr\_Int;

ARCHITECTURE Behavioral OF RS232\_Usr\_Int IS

TYPE state\_type IS (S0, S1, S2, S3, S4, S5, S6, S7, S8, S9, S10, S11, S12, S13, S14, S15, S16, S17, S18, S19, S20); SIGNAL CS\_RS232\_R, NS\_RS232\_R, CS\_RS232\_W, NS\_RS232\_W, CS\_FIFO\_Bus,

NS\_FIFO\_Bus : state\_type;

SIGNAL rx\_done, tx\_done : STD\_LOGIC := '0';

SIGNAL temp\_rcv : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0');

SIGNAL i, j : STD\_LOGIC\_VECTOR (15 DOWNTO 0) := (OTHERS => '0');

SIGNAL uartclk : STD\_LOGIC := '0';

SIGNAL u : INTEGER;

SIGNAL rs232\_rcv\_s, rs232\_rcv\_t : STD\_LOGIC := '1';

SIGNAL txbuff : STD\_LOGIC\_VECTOR(9 DOWNTO 0) := (OTHERS => '1'); --buff used to transmit 1 bytes with start and stop bits

--Declare Signals for FIFO Serial Read

SIGNAL STD\_FIFO\_R\_WriteEn, STD\_FIFO\_R\_ReadEn : STD\_LOGIC := '0';

SIGNAL STD\_FIFO\_R\_DataIn, STD\_FIFO\_R\_DataOut : STD\_LOGIC\_VECTOR(7

DOWNTO 0) := (OTHERS => '0');

SIGNAL STD\_FIFO\_R\_Empty, STD\_FIFO\_R\_Full : STD\_LOGIC := '0'; --Declare Signals for FIFO Serial Write

SIGNAL STD\_FIFO\_W\_WriteEn, STD\_FIFO\_W\_ReadEn : STD\_LOGIC := '0';

SIGNAL STD\_FIFO\_W\_DataIn, STD\_FIFO\_W\_DataOut : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0');

SIGNAL STD\_FIFO\_W\_Empty, STD\_FIFO\_W\_Full : STD\_LOGIC := '0'; --Declare Signals for Bus Interface

SIGNAL Bus\_Int1\_WE, Bus\_Int1\_RE, Bus\_Int1\_Busy : STD\_LOGIC := '0';

SIGNAL Bus\_Int1\_DataIn, Bus\_Int1\_DataOut, Bus\_Int1\_AddrIn :

STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0');

--Declare Signals for Registers

SIGNAL LD\_busy, LD\_busy2, LD\_rx, LD\_tx, LD\_temp\_data, LD\_temp2 :

 $STD\_LOGIC := '0';$ 

SIGNAL LD\_Temp\_Addr\_High, LD\_Temp\_Addr\_Low, LD\_Temp\_Data\_High : STD\_LOGIC := '0';

SIGNAL LD\_Temp\_Data\_Low, ld\_temp\_cmd : STD\_LOGIC := '0';

SIGNAL busy, busy\_reg\_o, busy2, busy2\_reg\_o, rx, rx\_reg\_o, tx, tx\_reg\_o : STD\_LOGIC := '0';

SIGNAL temp\_data\_reg\_o, temp\_data : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0');

SIGNAL temp2\_reg\_o, temp2 : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0');

SIGNAL Temp\_Addr\_High\_reg\_o, Temp\_Addr\_High : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0');

SIGNAL Temp\_Addr\_Low\_reg\_o, Temp\_Addr\_Low : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0');

```
SIGNAL Temp Data High reg o, Temp Data High : STD LOGIC VECTOR(7
DOWNTO 0) := (OTHERS => '0');
      SIGNAL Temp_Data_Low_reg_o, Temp_Data_Low : STD_LOGIC_VECTOR(7
DOWNTO 0) := (OTHERS = '0');
      SIGNAL Temp_Cmd_reg_o, Temp_Cmd : STD_LOGIC_VECTOR(7 DOWNTO 0) :=
(OTHERS = '0');
      ----User defined variables
      -- CM is the Clock Divder 25MHz/CM=115,200 Baud
      CONSTANT CM : INTEGER := clk in/Baud;
      -- CN is the read offset for serial input
      CONSTANT CN : INTEGER := CM/2;
      ----End User defined variables
      --declare STD FIFO
      COMPONENT STD_FIFO
            GENERIC (
                   DATA WIDTH : INTEGER; -- Width of FIFO
                   FIFO DEPTH : INTEGER; --
                                                  Depth of FIFO
                   FIFO_ADDR_LEN : INTEGER -- Required number of bits to represent
FIFO_Depth
            );
            PORT (
                   CLK: IN STD LOGIC;
                   RST : IN STD_LOGIC;
                   WriteEn : IN STD LOGIC;
                   DataIn : IN STD_LOGIC_VECTOR (7 DOWNTO 0);
                   ReadEn : IN STD LOGIC:
                   DataOut : OUT STD_LOGIC_VECTOR (7 DOWNTO 0);
                   Empty : OUT STD LOGIC;
                   Full : OUT STD_LOGIC
            ):
      END COMPONENT;
      --declare Bus Interface
      COMPONENT Bus Int
            PORT (
                   clk : IN std_logic;
                   rst : IN std logic;
                   DataIn : IN std_logic_vector(15 DOWNTO 0);
                   DataOut : OUT std logic vector(15 DOWNTO 0);
                   AddrIn : IN std_logic_vector(15 DOWNTO 0);
                   WE : IN std logic;
                   RE : IN std_logic;
                   Busy : OUT std_logic;
                   Data : INOUT std_logic_vector(15 DOWNTO 0);
                   Addr : OUT std logic vector(15 DOWNTO 0);
                   Xrqst : OUT std_logic;
                   XDat : IN std logic;
```

```
YDat : OUT std_logic;
                   BusRqst : OUT std_logic;
                   BusCtrl : IN std_logic
            );
      END COMPONENT;
BEGIN
      --Instantiate STD_FIFO for Reading Serial Data
      STD_FIFO_R : STD_FIFO
      GENERIC MAP
      (
            DATA_WIDTH => 8, -- Width of FIFO
            FIFO_DEPTH => 512, --
                                      Depth of FIFO
            FIFO_ADDR_LEN => 9 -- Required number of bits to represent FIFO_Depth
      )
      PORT MAP
      (
            CLK => clk,
            RST => rst,
            WriteEn => STD_FIFO_R_WriteEn,
            DataIn => STD_FIFO_R_DataIn,
            ReadEn => STD FIFO R ReadEn,
            DataOut => STD FIFO R DataOut,
            Empty => STD_FIFO_R_Empty,
            Full => STD FIFO R Full
      );
      --Instantiate STD FIFO for Writing Serial Data
      STD_FIFO_W : STD_FIFO
      GENERIC MAP
      (
            DATA_WIDTH => 8, -- Width of FIFO
            FIFO DEPTH => 512, --
                                      Depth of FIFO
            FIFO_ADDR_LEN => 9 -- Required number of bits to represent FIFO_Depth
      PORT MAP(
            CLK => clk,
            RST => rst.
            WriteEn => STD_FIFO_W_WriteEn,
            DataIn => STD FIFO W DataIn,
            ReadEn => STD FIFO W ReadEn,
            DataOut => STD_FIFO_W_DataOut,
            Empty => STD_FIFO_W_Empty,
            Full => STD_FIFO_W_Full
      );
      --Instantiate Bus Interface
      Bus_Int1 : Bus_Int PORT MAP(
            clk => clk,
```

```
rst => rst,
      DataIn => Bus_Int1_DataIn,
      DataOut => Bus_Int1_DataOut,
      AddrIn => Bus_Int1_AddrIn,
      WE => Bus_Int1_WE,
      RE => Bus_Int1_RE,
      Busy => Bus_Int1_Busy,
      Data => Data,
      Addr => Addr,
      Xrqst => Xrqst,
      XDat => XDat,
      YDat => YDat,
      BusRqst => BusRqst,
      BusCtrl => BusCtrl
);
----Registers
Reg_Proc : PROCESS
BEGIN
       WAIT UNTIL clk'event AND clk = '1';
      IF rst = '0' THEN
             busy reg o \leq 0';
             busy2_reg_o <= '0';
             rx_reg_o <= '0';
             tx reg o \leq 0';
             temp_data_reg_o <= (OTHERS => '0');
             temp2 reg o \leq (OTHERS = '0');
             Temp_Addr_High_reg_o <= (OTHERS => '0');
             Temp Addr Low reg o \le (OTHERS \implies 0');
             Temp_Data_High_reg_o <= (OTHERS => '0');
             Temp Data Low reg o \le (OTHERS => '0');
             Temp_Cmd_reg_o <= (OTHERS => '0');
      ELSE
             IF (LD busy = '1') THEN
                    busy reg o \leq busy;
             END IF:
             IF (LD busy2 = '1') THEN
                    busy2_reg_o <= busy2;
             END IF:
             IF (LD rx = '1') THEN
                    rx_reg_o <= rx;</pre>
             END IF;
             IF (LD_tx = '1') THEN
                    tx\_reg\_o \le tx;
             END IF:
             IF (LD_temp_data = '1') THEN
                    temp_data_reg_o <= temp_data;
```

END IF: IF (LD\_temp2 = '1') THEN temp2\_reg\_o <= temp2;</pre> END IF; IF (LD\_Temp\_Addr\_High = '1') THEN Temp\_Addr\_High\_reg\_o <= Temp\_Addr\_High; END IF: IF (LD\_Temp\_Addr\_Low = '1') THEN Temp\_Addr\_Low\_reg\_o <= Temp\_Addr\_Low; END IF; IF (LD\_Temp\_Data\_High = '1') THEN Temp\_Data\_High\_reg\_o <= Temp\_Data\_High; END IF: IF (LD\_Temp\_Data\_Low = '1') THEN Temp\_Data\_Low\_reg\_o <= Temp\_Data\_Low; END IF; IF (LD\_Temp\_Cmd = '1') THEN Temp\_Cmd\_reg\_o <= Temp\_Cmd; END IF; END IF; END PROCESS; ----End Registers ----Next State Logic for Serial Interface Read NSL RS232 R : PROCESS (CS RS232 R, rs232 rcv s, rx done, STD FIFO R Full, temp\_rcv) **BEGIN** ----Default States to remove latches busy  $\leq 0'$ ; rx <= '0'; NS RS232 R <= S0;  $LD_busy \ll 0';$ LD\_rx <= '0'; --Signals for FIFO STD FIFO R WriteEn <= '0'; STD\_FIFO\_R\_DataIn <= (OTHERS => '0'); CASE CS RS232 R IS WHEN S0 => -- Waits until data is detected on rs232\_rcv\_s. IF (rs232 rcv s = '1') THEN NS\_RS232\_R <= S0; ELSE NS\_RS232\_R <= S1; END IF: busy  $\leq 0'$ ; -- the busy signal stops the baud generator  $rx \ll 0'$ ; -- signals to stop reading data LD\_rx <= '1'; LD busy  $\leq 1'$ ;

```
WHEN S1 => -- Starts the baud rate generator and reading
                          NS_RS232_R <= S2;
                          busy <= '1'; -- the busy signal starts the baud generator
                          rx \ll 1'; -- signals to start reading data
                          LD_rx <= '1';
                          LD_busy <= '1';
                    WHEN S2 => -- Waits until all data is read
                          IF (rx\_done = '0') THEN
                                 NS_RS232_R <= S2;
                          ELSE
                                 NS_RS232_R <= S3;
                          END IF;
                    WHEN S3 =>
                          IF (STD_FIFO_R_Full = '0') THEN
                                 STD_FIFO_R_DataIn <= temp_rcv;
                                 STD FIFO R WriteEn <= '1';
                          END IF;
                          NS_RS232_R <= S0;
                    WHEN OTHERS =>
                          NS_RS232_R <= S0;
             END CASE;
      END PROCESS:
      ----End Next State Logic for Serial Interface Read
      ----Next State Logic for Serial Interface Write
      NSL_RS232_W : PROCESS (CS_RS232_W, tx_done, STD_FIFO_W_Empty,
STD FIFO W DataOut)
      BEGIN
             ----Default States to remove latches
             tx <= '0';
             NS RS232 W \leq S0;
             temp2 \leq (OTHERS = '0');
             LD_tx <= '0';
             LD_temp2 <= '0';
             Busy2 <= '0';
             LD_Busy2 <= '0';
             --Signals for FIFO
             STD_FIFO_W_ReadEn <= '0';
             CASE CS RS232 W IS
                    WHEN S0 =>
                          IF (STD_FIFO_W_Empty = '1') THEN
                                 NS_RS232_W <= S0;
                          ELSE
                                 NS_RS232_W <= S1;
                                 STD_FIFO_W_ReadEn <= '1';
                          END IF:
                          busy2 \le 0'; -- the busy signal stops the baud generator
```

 $tx \le 0'$ ; -- signals to stop sending data LD\_tx <= '1'; LD\_busy2 <= '1'; WHEN S1 =>temp2 <= STD\_FIFO\_W\_DataOut; LD\_temp2 <= '1'; NS\_RS232\_W <= S2; WHEN S2 =>busy $2 \le 1'$ ; -- the busy signal starts the baud generator  $tx \le '1'$ ; -- signals to start sending data LD tx <= '1';  $LD_busy2 <= '1';$ NS\_RS232\_W <= S3; WHEN S3  $\Rightarrow$ IF  $(tx\_done = '0')$  THEN NS RS232 W <= S3; ELSE NS\_RS232\_W <= S0; END IF; WHEN OTHERS => NS RS232 W  $\leq$  S0; END CASE; END PROCESS; ----End Next State Logic for Serial Interface Write ----Next State Logic for FIFO to Bus NSL FIFO Bus : PROCESS (CS FIFO Bus, STD FIFO R Empty, Temp Cmd reg o, Bus\_Int1\_Busy, STD\_FIFO\_R\_DataOut, Temp\_Addr\_High\_reg\_o, Temp\_Addr\_Low\_reg\_o, Temp Data High reg o, Temp Data Low reg o, Bus Int1 DataOut, temp data reg o) **BEGIN** ----Default States to remove latches NS FIFO Bus  $\leq$  S0; Temp\_Cmd  $\leq$  (OTHERS = '0'); LD Temp Cmd  $\leq 0'$ ; Temp Addr High  $\leq$  (OTHERS = '0'); LD\_Temp\_Addr\_High <= '0'; Temp Addr Low  $\leq$  (OTHERS = '0'); LD\_Temp\_Addr\_Low <= '0'; Bus Int1 AddrIn  $\leq$  (OTHERS = '0'); Bus Int1 RE  $\leq 0'$ ; Bus \_Int1\_DataIn <= (OTHERS => '0'); Bus\_Int1\_WE <= '0'; Temp\_Data  $\leq$  (OTHERS => '0'); LD Temp Data  $\leq 0'$ ; Temp Data High  $\leq$  (OTHERS = '0'); LD\_Temp\_Data\_High <= '0'; Temp Data High  $\leq$  (OTHERS = '0');

Temp Data Low  $\leq$  (OTHERS = '0'); LD\_Temp\_Data\_Low <= '0'; --Signals for FIFO STD FIFO R ReadEn <= '0';  $STD_FIFO_W_DataIn \le (OTHERS \implies '0');$ STD\_FIFO\_W\_WriteEn <= '0'; CASE CS\_FIFO\_Bus IS WHEN S0 => IF (STD FIFO R Empty = '1') THEN --Check to see if commands are in queue NS FIFO Bus  $\leq$  S0; ELSE NS\_FIFO\_Bus <= S1; STD FIFO R ReadEn <= '1'; --Assert Read Signal for **FIFO** END IF; WHEN S1 => --Read Command from FIFO Temp\_Cmd <= STD\_FIFO\_R\_DataOut; LD Temp Cmd  $\leq 1'$ ; NS\_FIFO\_Bus <= S2; WHEN S2 =>IF (STD\_FIFO\_R\_Empty = '1') THEN --Check to see if commands are in queue NS FIFO Bus  $\leq 22$ ; ELSE NS FIFO Bus  $\leq$  S3; STD\_FIFO\_R\_ReadEn <= '1'; --Assert Read Signal for **FIFO** END IF; WHEN S3 => --Read Address High from FIFO Temp\_Addr\_High <= STD\_FIFO\_R\_DataOut; LD\_Temp\_Addr\_High <= '1'; NS\_FIFO\_Bus <= S4; WHEN S4 =>IF (STD\_FIFO\_R\_Empty = '1') THEN --Check to see if commands are in queue NS\_FIFO\_Bus <= S4; ELSE NS FIFO Bus  $\leq$  S5; STD\_FIFO\_R\_ReadEn <= '1'; END IF; WHEN S5 => --Read Address\_Low from FIFO Temp Addr Low <= STD FIFO R DataOut; LD Temp Addr Low <= '1'; NS\_FIFO\_Bus <= S6; WHEN S6 =>

IF (Temp\_Cmd\_reg\_o = X"70") THEN --Check Cmd (Read) NS\_FIFO\_Bus <= S7; ELSIF (Temp\_Cmd\_reg\_o = X"71") THEN --Check Cmd (Write) NS FIFO Bus <= S15; ELSE -- Check Cmd (Invalid Data) NS FIFO Bus  $\leq$  S0; END IF: --Read from Bus and Write to RS232 FIFO WHEN S7 =>Bus\_Int1\_AddrIn(15 DOWNTO 8) <= Temp\_Addr\_High\_reg\_o; --Send Address to Bus Interface for Read Bus\_Int1\_AddrIn(7 DOWNTO 0) <= Temp\_Addr\_Low\_reg\_o; --Send Address to Bus Interface for Read Bus\_Int1\_RE <= '1'; --Read Flag to Bus Interface NS\_FIFO\_Bus <= S8; WHEN S8 => -- Wait until data is ready IF (Bus\_Int1\_Busy = '1') THEN NS\_FIFO\_Bus <= S8; ELSE NS\_FIFO\_Bus <= S9; END IF; Temp Data <= Bus Int1 DataOut; LD\_Temp\_Data <= '1'; WHEN S9 => -- Form First byte of Packet(Start Deliminator) STD\_FIFO\_W\_DataIn <= X"7E"; STD FIFO W WriteEn <= '1'; NS\_FIFO\_Bus <= S10; WHEN S10 => --Form Second byte of Packet(Address High) STD\_FIFO\_W\_DataIn <= Temp\_Addr\_High\_reg\_o; STD FIFO W WriteEn <= '1': NS FIFO Bus  $\leq$  S11; WHEN S11 => --Form Third byte of Packet(Address\_Low) STD FIFO W DataIn <= Temp Addr Low reg o; STD FIFO W WriteEn <= '1'; NS\_FIFO\_Bus <= S12; WHEN S12 = --Form Fourth byte of Packet(Data High) STD\_FIFO\_W\_DataIn <= Temp\_Data\_reg\_o(15 DOWNTO 8); STD FIFO W WriteEn <= '1'; NS FIFO Bus  $\leq$  S13; WHEN S13 => --Form Fifth byte of Packet(Data\_Low) STD\_FIFO\_W\_DataIn <= Temp\_Data\_reg\_o(7 DOWNTO 0); STD\_FIFO\_W\_WriteEn <= '1'; NS FIFO Bus  $\leq$  S0; --End Read from Bus and Write to RS232 FIFO --Write to Bus from RS232 FIFO WHEN S15 =>

IF (STD FIFO R Empty = '1') THEN --Check to see if commands are in queue NS\_FIFO\_Bus <= S15; ELSE NS\_FIFO\_Bus <= S16; STD\_FIFO\_R\_ReadEn <= '1'; END IF: WHEN S16 => --Read Data\_High from FIFO Temp Data High <= STD FIFO R DataOut; LD\_Temp\_Data\_High <= '1'; NS FIFO Bus  $\leq$  S17; WHEN S17 => IF (STD\_FIFO\_R\_Empty = '1') THEN --Check to see if commands are in queue NS\_FIFO\_Bus  $\leq$  S17; ELSE NS\_FIFO\_Bus <= S18; STD\_FIFO\_R\_ReadEn <= '1'; END IF: WHEN S18 => --Read Data\_Low from FIFO Temp Data Low <= STD FIFO R DataOut; LD Temp Data Low <= '1': NS\_FIFO\_Bus <= S19; WHEN S19 => Bus\_Int1\_AddrIn(15 DOWNTO 8) <= Temp\_Addr\_High\_reg\_o; --Send Address to Bus Interface for Write Bus\_Int1\_AddrIn(7 DOWNTO 0) <= Temp\_Addr\_Low\_reg\_o; --Send Address to Bus Interface for Write Bus\_Int1\_DataIn(15 DOWNTO 8) <= Temp\_Data\_High\_reg\_o; --Send Data to Bus Interface for Write Bus\_Int1\_DataIn(7 DOWNTO 0) <= Temp\_Data\_Low\_reg\_o; --Send Data to Bus Interface for Write Bus\_Int1\_WE <= '1'; --Write Flag to Bus Interface NS FIFO Bus <= S20; WHEN S20 => -- Wait until data is ready IF (Bus\_Int1\_Busy = '1') THEN NS\_FIFO\_Bus <= S20; ELSE NS\_FIFO\_Bus <= S0; END IF: --End Write to Bus from RS232 FIFO WHEN OTHERS => NS FIFO Bus  $\leq$  S0; END CASE: END PROCESS: ----End Next State Logic for FIFO to Bus

```
----UART Clock Divider
UART_Clk : PROCESS
BEGIN
      WAIT UNTIL clk'event AND clk = '1';
      --Synchronize async signal
      rs232_rcv_t <= rs232_rcv; --Synchro1 rs232_rcv
      rs232_rcv_s <= rs232_rcv_t; --Synchro2 rs232_rcv
      IF (rst = '0' OR (busy_reg_o = '0' AND busy2_reg_o = '0')) THEN
             uartclk <= '0';
             i <= CONV_STD_LOGIC_VECTOR(CN, 16);
      ELSIF (i = CM) THEN
             uartclk <= '1';
             i <= X"0000";
      ELSE
             i \le i + 1;
             uartclk <= '0';
      END IF:
END PROCESS;
---- End UART Clock Divider
----UART_Read
UART Read : PROCESS
BEGIN
      WAIT UNTIL clk'event AND clk = '1';
      IF rst = '0' OR rx reg o = '0' THEN
             temp_rcv <= x"00";
             j <= x"0000";
             rx_done <= '0';
      ELSIF rx reg o = '1' THEN
             IF uartclk = '1' THEN
                    IF j < X"09" THEN
                          temp_rcv(7) <= rs232_rcv_s;
                          temp_rcv(6 DOWNTO 0) <= temp_rcv(7 DOWNTO 1);
                          i \le i + 1;
                          rx_done <= '0';
                    ELSE
                          j <= X"0000";
                          rx_done <= '1';
                   END IF;
             ELSE
                    rx_done \le 0';
             END IF;
      END IF:
END PROCESS;
----End UART Read
-----UART_Xmit
UART Xmit: PROCESS
```

```
BEGIN
           WAIT UNTIL clk'event AND clk = '1';
           IF (rst = '0' OR tx_reg_o = '0') THEN
                 rs232 xmt <= '1';
                 tx_done <= '0';
                 u <= 0;
                 --structure the 10-bit frame to be sent
                 txbuff(9) <= '1'; --stopbit 2
                 txbuff(8 DOWNTO 1) <= temp2 reg o;
                 txbuff(0) \le 0'; --startbit 2
           ELSE
                 IF uartclk = '1' THEN
                       IF (u < 10) THEN
                             rs232_xmt <= txbuff(0);
                             txbuff(8 DOWNTO 0) <= txbuff(9 DOWNTO 1);</pre>
                             tx done \leq 0';
                             u \le u + 1;
                       ELSE
                             u <= 0;
                             tx_done <= '1';
                       END IF;
                 END IF;
           END IF:
     END PROCESS;
      -----End UART_Xmit
     ----State Sync
     sync_States : PROCESS
     BEGIN
           WAIT UNTIL clk'event AND clk = '1';
           IF rst = '0' THEN
                 CS RS232 R \leq S0;
                 CS_RS232_W <= S0;
                 CS_FIFO_Bus <= S0;
           ELSE
                 CS_RS232_R \le NS_RS232_R;
                 CS RS232 W <= NS_RS232_W;
                 CS_FIFO_Bus <= NS_FIFO_Bus;
           END IF:
     END PROCESS;
     ----End State Sync
END Behavioral:
LIBRARY IEEE;
USE IEEE.STD_LOGIC_1164.ALL;
USE IEEE.STD LOGIC ARITH.ALL;
```

101

USE IEEE.STD LOGIC UNSIGNED.ALL; USE ieee.numeric\_std.ALL; ENTITY LED\_Ctrl IS GENERIC ( Addr\_LED\_En : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0100"; --Enable LED Outputs (LSB) Addr\_LED\_Freq : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0101"; --LED **Blink Frequency** Addr LED PW : STD LOGIC VECTOR(15 DOWNTO 0) := X"0102"; --LED Pulse Width (On-Time) Addr LED1 DC: STD LOGIC VECTOR(15 DOWNTO 0) := X"0103"; --LED1 PWM Duty Cycle Addr LED2 DC: STD LOGIC VECTOR(15 DOWNTO 0) := X"0104"; --LED2 PWM Duty Cycle Addr\_LED3\_DC : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0105"; --LED3 PWM Duty Cycle Addr\_LED4\_DC : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0106"; --LED4 PWM Duty Cycle Addr LED5 DC : STD LOGIC VECTOR(15 DOWNTO 0) := X"0107"; --LED5 PWM Duty Cycle Addr LED6 DC : STD LOGIC VECTOR(15 DOWNTO 0) := X"0108"; --LED6 PWM Duty Cycle Addr\_LED7\_DC : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0109"; --LED7 PWM Duty Cycle Addr\_LED8\_DC : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"010A" --LED8 PWM Duty Cycle ); PORT ( clk: IN STD LOGIC; rst : IN STD LOGIC; Data : INOUT std logic vector(15 DOWNTO 0); Addr: OUT std logic vector(15 DOWNTO 0); Xrqst : OUT std logic; XDat : IN std logic; YDat : OUT std logic; BusRqst : OUT std logic; BusCtrl : IN std\_logic; LED1 Out : OUT STD LOGIC ); END LED Ctrl; ARCHITECTURE Behavioral OF LED Ctrl IS TYPE state\_type IS (S0, S1, S2, S3, S4, S5, S6, S7, S8, S9, S10, S11, S12, S13, S14, S15, S16, S17); SIGNAL CS Bus, NS Bus, CS Blink, NS Blink : state type; --declare Std\_Counter Component COMPONENT Std Counter IS

GENERIC ( Width : INTEGER --width of counter ); PORT ( INC, rst, clk : IN std\_logic; Count : OUT STD\_LOGIC\_VECTOR(Width - 1 DOWNTO 0)); END COMPONENT: --Declare PWM COMPONENT PWM 16b IS GENERIC ( Freq\_in : INTEGER; --Clk Max\_PWM : INTEGER; --PWM Resolution (2^16-1) Freq\_Sw : INTEGER --Switching Freq ); PORT ( clk: IN std logic; rst : IN std\_logic; DC : IN std\_logic\_vector(15 DOWNTO 0); Phase : IN std\_logic\_vector(15 DOWNTO 0); En : IN std\_logic; PWM\_Out : OUT std\_logic ): END COMPONENT: --declare Bus Interface COMPONENT Bus\_Int PORT ( clk : IN std\_logic; rst : IN std logic; DataIn : IN std\_logic\_vector(15 DOWNTO 0); DataOut : OUT std\_logic\_vector(15 DOWNTO 0); AddrIn : IN std\_logic\_vector(15 DOWNTO 0); WE : IN std\_logic; RE : IN std\_logic; Busy : OUT std logic; Data : INOUT std\_logic\_vector(15 DOWNTO 0); Addr : OUT std logic vector(15 DOWNTO 0); Xrqst : OUT std\_logic; XDat : IN std logic; YDat : OUT std\_logic; BusRqst : OUT std\_logic; BusCtrl : IN std\_logic ): END COMPONENT; ----Signals

SIGNAL PWM En, PWM Freq, PWM PW, PWM1 DC, PWM2 DC, PWM3 DC, PWM4\_DC : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0000"; --Set initial Duty Cycles to 0 SIGNAL PWM1\_En, PWM2\_En, PWM3\_En, PWM4\_En : STD\_LOGIC := '0'; CONSTANT PWM1\_Phase : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0000"; --Set Phase shift for PWM1 to 0 CONSTANT PWM2\_Phase : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0000"; --Set Phase shift for PWM2 to 0 CONSTANT PWM3\_Phase : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := X"0000"; --Set Phase shift for PWM3 to 0 CONSTANT PWM4 Phase : STD LOGIC VECTOR(15 DOWNTO 0) := X"0000"; --Set Phase shift for PWM4 to 0 --Max PWM Values CONSTANT PWM\_Max : std\_logic\_vector(15 DOWNTO 0) := X"FFFF"; CONSTANT PWM\_Min : std\_logic\_vector(15 DOWNTO 0) := X"0000"; --Declare Signals for Bus Interface SIGNAL Bus\_Int1\_WE, Bus\_Int1\_RE, Bus\_Int1\_Busy : STD\_LOGIC := '0'; SIGNAL Bus\_Int1\_DataIn, Bus\_Int1\_DataOut, Bus\_Int1\_AddrIn : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Bus\_Cnt\_rst, Bus\_Cnt\_INC : STD\_LOGIC := '0'; SIGNAL Bus Cnt Out : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Delay Cnt rst, Delay Cnt INC : STD LOGIC := '0'; SIGNAL Delay\_Cnt\_Out : STD\_LOGIC\_VECTOR(7 DOWNTO 0) := (OTHERS => '0'); --Signals for Registers SIGNAL LD PWM En, LD PWM Freq, LD PWM PW, LD PWM1 DC, LD\_PWM2\_DC, LD\_PWM3\_DC, LD\_PWM4\_DC : STD\_LOGIC := '0'; --Signals for Clock Divider SIGNAL clk\_temp : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL clk Blink : STD LOGIC := '0': --Signals for blink SIGNAL Freq\_Cnt\_rst, Freq\_Cnt\_Inc : STD\_LOGIC := '0'; SIGNAL Freq\_Cnt\_out : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); BEGIN --instantiate Bus Cnt Bus Cnt: Std Counter GENERIC MAP ( Width  $\Rightarrow 16$ PORT MAP ( clk => clk, rst => Bus Cnt rst, INC => Bus\_Cnt\_INC, Count => Bus Cnt Out

```
);
--instantiate Delay_Cnt
Delay_Cnt : Std_Counter
GENERIC MAP
(
      Width => 8
)
PORT MAP(
      clk => clk,
      rst => Delay_Cnt_rst,
      INC => Delay_Cnt_INC,
      Count => Delay_Cnt_Out
);
-- Instantiate PWM1
PWM1 : PWM_16b
GENERIC MAP
(
      Freq_in => 24930000,
      Max_PWM => 65535,
      Freq_Sw => 6104
)
PORT MAP(
      clk => clk,
      rst => rst,
      DC => PWM1_DC,
      Phase => PWM1 Phase,
      En => PWM1_En,
      PWM_Out => LED1_Out
);
--Instantiate Bus Interface
Bus_Int1 : Bus_Int PORT MAP(
      clk => clk,
      rst => rst,
      DataIn => Bus_Int1_DataIn,
      DataOut => Bus_Int1_DataOut,
      AddrIn => Bus_Int1_AddrIn,
      WE \Rightarrow Bus Int1 WE,
      RE => Bus_Int1_RE,
      Busy => Bus_Int1_Busy,
      Data => Data,
      Addr => Addr,
      Xrqst => Xrqst,
      XDat => XDat,
      YDat => YDat,
      BusRqst => BusRqst,
```

```
BusCtrl => BusCtrl
);
--instantiate Freq_Cnt
Freq_Cnt : Std_Counter
GENERIC MAP
(
      Width => 16
)
PORT MAP(
      clk => clk_Blink,
      rst => Freq_Cnt_rst,
      INC => Freq_Cnt_INC,
      Count => Freq_Cnt_Out
);
----Registers
Reg Proc : PROCESS
BEGIN
      WAIT UNTIL clk'event AND clk = '1';
      IF rst = '0' THEN
             PWM1_DC \le (OTHERS \implies '0');
             PWM2 DC \leq (OTHERS \geq '0');
             PWM3 DC \leq (OTHERS = '0');
             PWM4_DC \le (OTHERS \implies '0');
             PWM Freq \leq (OTHERS = '0');
             PWM_PW \le (OTHERS \implies '0');
             PWM En \ll (OTHERS \implies '0');
      ELSE
            IF (LD PWM1 DC = '1') THEN
                   PWM1_DC <= Bus_Int1_DataOut;
             END IF:
             IF (LD PWM2 DC = '1') THEN
                   PWM2_DC <= Bus_Int1_DataOut;
             END IF;
             IF (LD PWM3 DC = '1') THEN
                   PWM3_DC <= Bus_Int1_DataOut;
             END IF:
             IF (LD_PWM4_DC = '1') THEN
                   PWM4 DC <= Bus Int1 DataOut;
             END IF;
             IF (LD_PWM_Freq = '1') THEN
                   PWM_Freq <= Bus_Int1_DataOut;
             END IF:
             IF (LD PWM PW = '1') THEN
                   PWM PW <= Bus Int1 DataOut;
             END IF:
             IF (LD PWM En = '1') THEN
```

```
PWM En <= Bus Int1 DataOut;
             END IF;
      END IF:
END PROCESS;
----End Registers
----Next State Logic for Bus Interface
NSL_Bus : PROCESS (CS_Bus, Bus_Cnt_Out, Bus_Int1_Busy, Delay_Cnt_Out)
BEGIN
      ----Default States to remove latches
      NS_Bus \le S0;
      Bus_Int1_AddrIn <= (OTHERS => '0');
      Bus_Int1_RE <= '0';
      Bus_Int1_DataIn \leq (OTHERS = '0');
      Bus_Int1_WE <= '0';
      Bus_Cnt_rst <= '1';
      Bus Cnt INC \leq 0';
      LD_PWM1_DC <= '0';
      LD_PWM2_DC <= '0';
      LD_PWM3_DC <= '0';
      LD_PWM4_DC <= '0';
      LD PWM Freq \leq 0';
      LD PWM PW \leq 0':
      LD_PWM_En <= '0';
      Delay Cnt INC \leq 0';
      Delay_Cnt_rst <= '1';
      CASE CS_Bus IS
             WHEN S0 =>
                    Bus_Cnt_rst <= '0'; -- Reset Bus Counter
                    Delay_Cnt_rst <= '0'; -- Reset Delay Counter
                    NS Bus \leq S1;
             WHEN S1 => -- Initial Delay count for sync
                    IF (Delay_Cnt_Out < 40) THEN
                          NS Bus \leq S1;
                    ELSE
                          NS Bus \leq S2;
                    END IF;
                    Delay Cnt INC \leq 1';
             WHEN S2 => --Wait (2^{12}-34) Clk Cycles for 1x per fs
                    IF (Bus_Cnt_Out < 4062) THEN
                          NS_Bus \leq S2;
                    ELSE
                           NS_Bus \leq S3;
                    END IF;
                    Bus_Cnt_INC <= '1';
                    --Read Command Data from Bus
```

```
WHEN S3 =>
                          IF (Bus_Int1_Busy = '1') THEN
                                NS_Bus \leq S3;
                          ELSE
                                NS_Bus \leq S4;
                          END IF;
                          Bus_Cnt_rst <= '0'; -- Reset Bus Counter
                   WHEN S4 =>
                          Bus_Int1_AddrIn <= Addr_LED_En; --Read Data from LED_En
Register
                          Bus_Int1_RE <= '1';
                          NS_Bus \leq S5;
                   WHEN S5 =>
                          IF (Bus Int1 Busy = '1') THEN
                                NS_Bus \leq S5;
                          ELSE
                                LD_PWM_En <= '1';
                                NS_Bus \le S6;
                          END IF;
                   WHEN S6 =>
                          Bus_Int1_AddrIn <= Addr_LED_Freq; --Read Data from LED
Freq Register
                          Bus_Int1_RE <= '1';
                          NS Bus \leq S7;
                   WHEN S7 =>
                          IF (Bus Int1 Busy = '1') THEN
                                NS_Bus <= S7;
                          ELSE
                                LD_PWM_Freq \le '1';
                                NS_Bus \leq S8;
                          END IF;
                   WHEN S8 =>
                          Bus_Int1_AddrIn <= Addr_LED_PW; --Read Data from LED
PulseWidth Register
                          Bus_Int1_RE <= '1';
                          NS Bus \leq S9;
                   WHEN S9 =>
                          IF (Bus Int1 Busy = '1') THEN
                                NS_Bus <= S9;
                          ELSE
                                LD_PWM_PW <= '1';
                                NS_Bus <= S10;
                          END IF;
                   WHEN S10 =>
                          Bus_Int1_AddrIn <= Addr_LED1_DC; --Read Data from
```

LED1\_DC Register

Bus\_Int1\_RE <= '1'; NS\_Bus <= S11; WHEN S11 => IF (Bus\_Int1\_Busy = '1') THEN  $NS_Bus \le S11;$ ELSE LD\_PWM1\_DC <= '1';  $NS_Bus \le S12;$ END IF; WHEN S12 => Bus\_Int1\_AddrIn <= Addr\_LED2\_DC; --Read Data from LED2\_DC Register Bus\_Int1\_RE <= '1'; NS\_Bus <= S13; WHEN S13  $\Rightarrow$ IF (Bus Int1 Busy = '1') THEN  $NS_Bus \le S13;$ ELSE LD\_PWM2\_DC <= '1';  $NS_Bus \le S14;$ END IF; WHEN S14 => Bus\_Int1\_AddrIn <= Addr\_LED3\_DC; --Read Data from LED3\_DC Register Bus\_Int1\_RE <= '1'; NS Bus  $\leq$  S15; WHEN S15 => IF (Bus\_Int1\_Busy = '1') THEN NS\_Bus <= S15; ELSE LD PWM3 DC  $\leq 11'$ ; NS\_Bus <= S16; END IF; WHEN S16 => Bus\_Int1\_AddrIn <= Addr\_LED4\_DC; --Read Data from LED4\_DC Register Bus\_Int1\_RE <= '1'; NS Bus  $\leq$  S17; WHEN S17  $\Rightarrow$ IF (Bus\_Int1\_Busy = '1') THEN NS\_Bus <= S17; ELSE LD\_PWM4\_DC <= '1'; NS Bus  $\leq S2$ ; END IF; WHEN OTHERS =>

```
NS_Bus \le S0;
      END CASE;
END PROCESS;
----End Next State Logic for Bus Interface
----Next State Logic for Blink Update
NSL_Blink : PROCESS (CS_Blink, Freq_Cnt_Out, PWM_Freq, PWM_PW)
BEGIN
      ----Default States to remove latches
      NS Blink \leq S0;
      Freq_Cnt_INC <= '0';</pre>
      Freq_Cnt_rst <= '1';</pre>
      PWM1_En <= '0';
      PWM2_En <= '0';
      PWM3_En <= '0':
      PWM4_En <= '0';
      CASE CS Blink IS
             WHEN S0 =>
                    Freq_Cnt_rst <= '0'; -- Reset Period Counter
                    NS Blink <= S1;
             WHEN S1 => -- Counter for Pulse Width
                    IF (Freq_Cnt_Out < PWM_PW) THEN
                          NS_Blink <= S1;
                    ELSE
                          NS Blink <= S2;
                    END IF;
                    Freq Cnt INC <= '1';
                    PWM1_En <= '1';
                    PWM2 En <= '1';
                    PWM3_En <= '1';
                    PWM4 En <= '1';
             WHEN S2 => --Counter for Period
                    IF (Freq_Cnt_Out < PWM_Freq) THEN
                          NS_Blink <= S2;
                    ELSE
                          NS_Blink <= S0;
                    END IF:
                    Freq_Cnt_INC <= '1';</pre>
             WHEN OTHERS =>
                    NS_Blink \leq S0;
      END CASE;
END PROCESS;
----End Next State Logic for Blink Update
----State Sync
sync_States : PROCESS
BEGIN
      WAIT UNTIL clk'event AND clk = '1';
```

```
IF rst = '0' THEN
                 CS_Bus \le S0;
           ELSE
                 CS_Bus <= NS_Bus;
           END IF:
     END PROCESS;
      ----End State Sync
     ----State Sync for Blink
     sync_Blink : PROCESS
      BEGIN
            WAIT UNTIL clk Blink'event AND clk Blink = '1';
           IF rst = '0' THEN
                 CS_Blink <= S0;
           ELSE
                 CS_Blink <= NS_Blink;
           END IF;
     END PROCESS:
     ----End State Sync
     -- Clock Divider for LED Blink
     Clk_Div_Blink : PROCESS
     BEGIN
           WAIT UNTIL clk'event AND clk = '1';
           clk_temp <= clk_temp + 1;
           clk Blink \leq clk temp(9);
     END PROCESS;
END Behavioral:
LIBRARY IEEE;
USE IEEE.std_logic_1164.ALL;
USE ieee.std_logic_unsigned.ALL;
USE ieee.numeric_std.ALL;
LIBRARY lattice;
USE lattice.components.ALL;
LIBRARY machxo2;
USE machxo2.ALL;
ENTITY Bus_Interface_Top IS
     PORT (
           --RESETn : in STD_LOGIC;-- Global Reset
           -- RS232 Communication
           Usr_RX : IN STD_LOGIC; -- Serial In for User Control
           Usr TX : OUT STD LOGIC; -- Serial Out for User Control
           -- Board LEDs
           LED_1 : OUT STD_LOGIC -- Board LED
     );
```

```
111
```

END Bus\_Interface\_Top; ARCHITECTURE Behavioral OF Bus\_Interface\_Top IS -- Declare Internal Oscillator COMPONENT OSCH GENERIC (NOM\_FREQ : STRING := "8.31"); PORT ( STDBY : IN std\_logic; OSC : OUT std\_logic; SEDSTDBY : OUT std\_logic ); END COMPONENT; -- Declare PLL COMPONENT PLL\_Clk PORT ( ClkI : IN std\_logic; ClkOP : OUT std logic; Lock : OUT std\_logic ); END COMPONENT; -- Declare Bus\_Master **COMPONENT Bus Master** PORT ( clk : IN std\_logic; rst : IN std logic; Data : INOUT std\_logic\_vector(15 DOWNTO 0); Addr : IN std logic vector(15 DOWNTO 0); Xrqst : IN std\_logic; XDat : OUT std logic; YDat : IN std\_logic; BusRqst : IN std\_logic\_vector(9 DOWNTO 0); BusCtrl: OUT std logic vector(9 DOWNTO 0) ); END COMPONENT; -- Declare RS232 Usr Int COMPONENT RS232\_Usr\_Int GENERIC ( Baud : INTEGER; -- Baud Rate clk in : INTEGER -- Input Clk ); PORT ( clk : IN std\_logic; rst : IN std\_logic; rs232 rcv : IN std logic; rs232 xmt : OUT std logic; Data : INOUT std\_logic\_vector(15 DOWNTO 0); Addr : OUT std\_logic\_vector(15 DOWNTO 0);

Xrqst : OUT std logic; XDat : IN std\_logic; YDat : OUT std\_logic; BusRqst : OUT std logic; BusCtrl : IN std\_logic ); END COMPONENT; -- Declare LED Ctrl COMPONENT LED Ctrl IS PORT ( clk: IN STD LOGIC; rst : IN STD\_LOGIC; Data : INOUT std\_logic\_vector(15 DOWNTO 0); Addr : OUT std\_logic\_vector(15 DOWNTO 0); Xrqst : OUT std\_logic; XDat : IN std logic; YDat : OUT std logic; BusRqst : OUT std\_logic; BusCtrl : IN std logic; LED1\_Out : OUT STD\_LOGIC ); END COMPONENT: -- Declare Std\_Counter Component COMPONENT Std Counter IS GENERIC ( Width : INTEGER -- width of counter ); PORT ( INC, rst, clk : IN std logic; Count : OUT STD LOGIC VECTOR(Width - 1 DOWNTO 0)); END COMPONENT; ----Signals -- Declare Signals for Bus Interface SIGNAL Bus Int1 WE, Bus Int1 RE, Bus Int1 Busy : STD LOGIC := '0'; SIGNAL Bus\_Int1\_DataIn, Bus\_Int1\_DataOut, Bus\_Int1\_AddrIn : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); -- Inputs SIGNAL Addr : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Xrqst : STD LOGIC := '0'; SIGNAL YDat : STD LOGIC := '0': SIGNAL BusRqst : STD\_LOGIC\_VECTOR(9 DOWNTO 0) := (OTHERS => '0'); SIGNAL Data : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL XDat : STD LOGIC := '0'; SIGNAL BusCtrl : STD LOGIC VECTOR(9 DOWNTO 0) := (OTHERS => '0'); -- Internal Clock SIGNAL OSC Stdby, OSC Out, OSC SEDSTDBY, clk : STD LOGIC := '0';

-- Reset

```
SIGNAL PLL_Lock, System_rst : STD_LOGIC := '0';
      SIGNAL Reset_Cnt_INC, Reset_Cnt_rst : STD_LOGIC := '0';
      SIGNAL Reset_Cnt_out : STD_LOGIC_VECTOR(7 DOWNTO 0) := (OTHERS => '0');
      -- For inverting LED Outputs
      SIGNAL LED_1n : STD_LOGIC := '0';
BEGIN
      -- Instantiate Internal Oscillator
      Int_OSC : OSCH PORT MAP(
             STDBY => OSC_Stdby,
             OSC => OSC_Out,
             SEDSTDBY => OSC_SEDSTDBY
      );
      -- Instantiate PLL
      PLL_1 : PLL_Clk PORT MAP(
             ClkI => OSC Out,
             ClkOP => clk,
             Lock => Pll_Lock
      );
      -- Instantiate Bus_Master
      BM : Bus_Master PORT MAP(
             clk => clk,
             rst => System_rst,
             Data => Data,
             Addr => Addr,
             Xrqst => Xrqst,
             XDat => XDat,
             YDat => YDat,
             BusRqst => BusRqst,
             BusCtrl => BusCtrl
      );
      -- Instantiate RS232_Usr_Int
      RS232_Usr : RS232_Usr_Int
      GENERIC MAP
      (
             Baud => 9600, -- Baud Rate
             Clk_In => 24930000 --
                                       Input Clk
      )
      PORT MAP(
             clk => clk,
             rst => System_rst,
             rs232\_rcv => Usr\_RX,
             rs232_xmt => Usr_TX,
             Data => Data.
             Addr => Addr,
             Xrqst => Xrqst,
```

```
XDat => XDat,
       YDat => YDat,
      BusRqst => BusRqst(1),
      BusCtrl => BusCtrl(1)
);
-- Instantiate LED_Ctrl
LED_Ctrl1 : LED_Ctrl PORT MAP(
      clk => clk,
      rst => System_rst,
      Data => Data,
      Addr => Addr,
      Xrqst => Xrqst,
      XDat => XDat,
      YDat => YDat,
      BusRqst => BusRqst(0),
      BusCtrl => BusCtrl(0),
      LED1_Out => LED_1n
);
-- Instantiate Reset_Cnt_8
Reset_Cnt : Std_Counter
GENERIC MAP
(
       Width => 8
PORT MAP(
      clk \Rightarrow OSC Out,
      rst => Reset_Cnt_rst,
      INC \Rightarrow Reset Cnt INC,
      Count => Reset_Cnt_Out
);
-- Oscillator
OSC\_Stdby \le '0';
-- Tie unused ports to '0'
BusRqst(9 DOWNTO 2) \leq (OTHERS = '0');
-- Reset Block1
Reset Blk1 : PROCESS
BEGIN
       WAIT UNTIL OSC_Out'event AND OSC_Out = '1';
      IF (PLL_Lock = '0') THEN
              Reset_Cnt_rst <= '0';
      ELSE
             Reset_Cnt_rst <= '1';
      END IF;
END PROCESS;
-- Reset Block
Reset_Blk : PROCESS
```

BEGIN WAIT UNTIL OSC\_Out'event AND OSC\_Out = '1'; IF (Reset\_Cnt\_out < X"7F") THEN System rst  $\leq 0'$ ; Reset\_Cnt\_Inc <= '1'; ELSE System\_rst <= '1'; Reset\_Cnt\_Inc <= '0'; END IF: END PROCESS: -- LED Invert due to Active Low Configuration on Dev Board LED\_Invert : PROCESS BEGIN  $LED_1 \le NOT(LED_1n);$ END PROCESS; END Behavioral; LIBRARY IEEE; USE IEEE.std\_logic\_1164.ALL; USE IEEE.std logic arith.ALL; USE IEEE.std logic unsigned.ALL; USE IEEE.numeric\_std.ALL; LIBRARY machxo2; USE machxo2.ALL; LIBRARY lattice; USE lattice.components.ALL; ENTITY HW AUTH MODULE IS PORT ( MISO : IN STD LOGIC: CLK IN: IN STD LOGIC; CS: OUT STD LOGIC; CLK OUT : OUT STD LOGIC; MOSI: OUT STD LOGIC; HW\_GOOD : OUT STD\_LOGIC ): END HW\_AUTH\_MODULE; ARCHITECTURE Behavior OF HW AUTH MODULE IS TYPE MACHINE IS (START, SB, OPCODE\_H, OPCODE\_L, A5, A4, A3, A2, A1, A0, W, D15, D14, D13, D12, D11, D10, D9, D8, D7, D6, D5, D4, D3, D2, D1, D0, AUTHENTICATE, STDBY); SIGNAL STATE : MACHINE := STDBY; SIGNAL SB VALUE : STD LOGIC := '1'; SIGNAL OP\_READ : STD\_LOGIC\_VECTOR(1 DOWNTO 0) := "10"; SIGNAL EEPROM addr : STD LOGIC VECTOR(5 DOWNTO 0) := "000000";

SIGNAL DATA\_IN : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL DELAY : INTEGER RANGE 0 TO 200\_000 := 200\_000; --1\_500\_000 TYPE T\_ARRAY IS ARRAY(0 TO 63) OF STD\_LOGIC\_VECTOR(15 DOWNTO 0); SIGNAL KEYS : T\_ARRAY; SIGNAL KEY\_INDEX : INTEGER RANGE 0 TO 63 := 0;

# BEGIN

HW\_AUTH\_INTERRUPT : PROCESS (CLK\_IN)

# BEGIN

 $KEYS(0) \le X''ABBA'';$  $KEYS(1) \le X"ABED";$  $KEYS(2) \le X"BABE";$  $KEYS(3) \le X"BADE";$  $KEYS(4) \le X"BEAD";$  $KEYS(5) \le X"BEEF";$  $KEYS(6) \le X"CAFE";$  $KEYS(7) \le X"CEDE";$  $KEYS(8) \le X"DADA";$  $KEYS(9) \leq X"DEAD";$  $KEYS(10) \le X"DEAF";$  $KEYS(11) \le X"DEED";$ KEYS(12) <= X"FACE";  $KEYS(13) \le X"FADE";$  $KEYS(14) \le X"FEED";$ KEYS(15) <= X"FEE0";  $KEYS(16) \le X"ABBA":$  $KEYS(17) \le X"ABED";$  $KEYS(18) \le X"BABE";$  $KEYS(19) \le X"BADE";$  $KEYS(20) \le X"BEAD";$  $KEYS(21) \le X"BEEF";$  $KEYS(22) \le X"CAFE";$  $KEYS(23) \le X"CEDE";$  $KEYS(24) \le X"DADA";$  $KEYS(25) \le X"DEAD";$  $KEYS(26) \le X"DEAF";$ KEYS(27) <= X"DEED"; KEYS(28) <= X"FACE";  $KEYS(29) \le X"FADE";$ KEYS(30) <= X"FEED"; KEYS(31) <= X"FEE0";  $KEYS(32) \le X"ABBA";$  $KEYS(33) \le X"ABED";$  $KEYS(34) \le X"BABE";$  $KEYS(35) \le X"BADE";$  $KEYS(36) \le X"BEAD";$  KEYS(37) <= X"BEEF"; KEYS(38) <= X"CAFE"; KEYS(39) <= X"CEDE";  $KEYS(40) \le X"DADA";$  $KEYS(41) \le X"DEAD";$ KEYS(42) <= X"DEAF";  $KEYS(43) \le X"DEED";$  $KEYS(44) \le X"FACE";$  $KEYS(45) \le X"FADE";$ KEYS(46) <= X"FEED"; KEYS(47) <= X"FEE0";  $KEYS(48) \le X"ABBA";$  $KEYS(49) \le X"ABED";$ KEYS(50) <= X"BABE";  $KEYS(51) \le X"BADE";$  $KEYS(52) \le X"BEAD";$ KEYS(53) <= X"BEEF";  $KEYS(54) \le X"CAFE";$ KEYS(55) <= X"CEDE";  $KEYS(56) \le X"DADA";$  $KEYS(57) \le X"DEAD";$  $KEYS(58) \le X"DEAF";$ KEYS(59) <= X"DEED"; KEYS(60) <= X"FACE"; KEYS(61) <= X"FADE"; KEYS(62) <= X"FEED"; KEYS(63) <= X"FEE0"; IF (CLK\_IN' EVENT) THEN IF (CLK\_IN = '1') THEN CLK OUT  $\leq 1'$ ; END IF: IF (CLK IN = '0') THEN CLK\_OUT <= '0'; CASE STATE IS WHEN START => CS <= '0'; MOSI <= '0':  $STATE \le SB;$ WHEN SB => CS <= '1': MOSI <= SB\_VALUE; STATE <= OPCODE H; WHEN OPCODE\_H => CS <= '1';  $MOSI \le OP\_READ(1);$ STATE <= OPCODE\_L; WHEN OPCODE\_L => CS <= '1';  $MOSI \le OP_READ(0);$ STATE  $\leq A5$ ; WHEN A5 =>CS <= '1'; MOSI <= EEPROM\_addr(5); STATE  $\leq A4;$ WHEN A4 =>CS <= '1'; MOSI <= EEPROM\_addr(4); STATE  $\leq A3$ ; WHEN A3 =>CS <= '1'; MOSI <= EEPROM\_addr(3); STATE  $\leq A2$ ; WHEN A2  $\Rightarrow$ CS <= '1'; MOSI <= EEPROM\_addr(2); STATE  $\leq A1$ ; WHEN A1 =>CS <= '1'; MOSI <= EEPROM\_addr(1); STATE  $\leq A0;$ WHEN A0 =>CS <= '1'; MOSI <= EEPROM\_addr(0); STATE <= W; WHEN W =>CS <= '1'; STATE <= D15; WHEN D15 =>

CS <= '1'; MOSI <= '0';  $DATA_IN(15) \le MISO;$ STATE <= D14; WHEN D14 => CS <= '1'; MOSI <= '0';  $DATA_IN(14) \le MISO;$ STATE <= D13; WHEN D13 => CS <= '1'; MOSI <= '0';  $DATA_IN(13) \le MISO;$ STATE <= D12; WHEN D12 => CS <= '1'; MOSI <= '0';  $DATA_IN(12) \le MISO;$ STATE <= D11; WHEN D11 => CS <= '1'; MOSI <= '0';  $DATA_IN(11) \le MISO;$ STATE <= D10; WHEN D10 => CS <= '1'; MOSI <= '0';  $DATA_IN(10) \le MISO;$ STATE  $\leq D9$ ; WHEN D9 =>CS <= '1'; MOSI <= '0';  $DATA_IN(9) \le MISO;$ STATE  $\leq D8$ ; WHEN D8 =>CS <= '1'; MOSI <= '0';  $DATA_IN(8) \le MISO;$ STATE  $\leq D7$ ;

WHEN D7 =>CS <= '1'; MOSI <= '0';  $DATA_IN(7) \le MISO;$ STATE  $\leq D6$ ; WHEN D6 => CS <= '1'; MOSI <= '0';  $DATA_IN(6) \le MISO;$ STATE  $\leq D5$ ; WHEN D5 =>CS <= '1'; MOSI <= '0'; DATA\_IN(5) <= MISO; STATE  $\leq D4$ ; WHEN D4 =>CS <= '1'; MOSI <= '0';  $DATA_IN(4) \le MISO;$ STATE  $\leq D3$ ; WHEN D3 =>CS <= '1'; MOSI <= '0';  $DATA_IN(3) \le MISO;$ STATE  $\leq D2;$ WHEN D2 =>CS <= '1'; MOSI <= '0';  $DATA_IN(2) \le MISO;$ STATE  $\leq D1$ ; WHEN D1 =>CS <= '1'; MOSI <= '0';  $DATA_IN(1) \le MISO;$ STATE  $\leq D0;$ WHEN D0 =>CS <= '1'; MOSI <= '0';

```
DATA_IN(0) \le MISO;
                              STATE <= AUTHENTICATE;
                        WHEN AUTHENTICATE =>
                             CS <= '0';
                              MOSI <= '0';
                              IF (DATA_IN = KEYS(KEY_INDEX)) THEN
                                  HW_GOOD <= '1';
                              ELSE
                                  HW_GOOD <= '0';
                              END IF;
                              STATE <= STDBY;
                        WHEN STDBY =>
                             CS <= '0';
                              MOSI <= '0';
                              IF (DELAY > 0) THEN
                                  DELAY \le DELAY - 1;
                                  STATE <= STDBY;
                              ELSIF (DELAY = 0) THEN
                                  EEPROM addr <= EEPROM addr +
"000001";
                                  KEY_INDEX <= KEY_INDEX + 1;
                                  DELAY <= 200 000;
                                   STATE <= START;
                              ELSE
                                   DELAY <= 200_000;
                                  STATE <= STDBY;
                              END IF:
                        WHEN OTHERS =>
                              CS <= '0':
                              MOSI <= '0';
                              STATE <= STDBY;
                   END CASE;
              END IF;
         END IF;
    END PROCESS;
END Behavior;
LIBRARY ieee;
USE ieee.std_logic_1164.ALL;
LIBRARY lattice:
USE lattice.components.ALL;
ENTITY HARDWARE ASSISTED SUPERVISOR IS
```

GENERIC (

TIMEOUT : INTEGER :=  $2_{000}$ 

); PORT (

CLK\_IN : IN STD\_LOGIC; --Input Controller 1 I1\_0 : IN STD\_LOGIC; I1\_1 : IN STD\_LOGIC; I1\_2 : IN STD\_LOGIC; I1\_3 : IN STD\_LOGIC; I1\_4 : IN STD\_LOGIC; I1\_5 : IN STD\_LOGIC; I1\_6 : IN STD\_LOGIC; I1\_7 : IN STD\_LOGIC; I1\_8 : IN STD\_LOGIC; I1 9: IN STD LOGIC; I1\_10 : IN STD\_LOGIC; I1\_11 : IN STD\_LOGIC; I1\_12 : IN STD\_LOGIC; I1\_13 : IN STD\_LOGIC; I1 14 : IN STD LOGIC; I1\_15 : IN STD\_LOGIC; I1\_16 : IN STD\_LOGIC; I1 17 : IN STD LOGIC; I1\_18 : IN STD\_LOGIC; I1 19: IN STD LOGIC; I1\_20 : IN STD\_LOGIC; I1 21 : IN STD LOGIC; I1\_22 : IN STD\_LOGIC; I1 23 : IN STD LOGIC; I1\_24 : IN STD\_LOGIC; I1\_25 : IN STD\_LOGIC; I1\_26 : IN STD\_LOGIC; I1 27 : IN STD LOGIC; --Input Controller 2 I2 0: IN STD LOGIC; I2\_1 : IN STD\_LOGIC; I2 2: IN STD LOGIC; I2 3 : IN STD LOGIC; I2\_4 : IN STD\_LOGIC; I2\_5 : IN STD\_LOGIC; I2\_6 : IN STD\_LOGIC; I2 7 : IN STD LOGIC; I2 8: IN STD LOGIC; I2\_9 : IN STD\_LOGIC; I2 10: IN STD LOGIC;

I2\_11 : IN STD\_LOGIC; I2\_12 : IN STD\_LOGIC; I2\_13 : IN STD\_LOGIC; I2\_14 : IN STD\_LOGIC; I2\_15 : IN STD\_LOGIC; I2\_16 : IN STD\_LOGIC; I2\_17 : IN STD\_LOGIC; I2\_18 : IN STD\_LOGIC; I2\_19 : IN STD\_LOGIC; I2\_20 : IN STD\_LOGIC; I2\_21 : IN STD\_LOGIC; I2\_22 : IN STD\_LOGIC; I2\_23 : IN STD\_LOGIC; I2\_24 : IN STD\_LOGIC; I2\_25 : IN STD\_LOGIC; I2 26 : IN STD LOGIC; I2\_27 : IN STD\_LOGIC; --Output Controller O\_0: OUT STD\_LOGIC; O\_1 : OUT STD\_LOGIC; O 2: OUT STD LOGIC; O 3: OUT STD LOGIC; O\_4 : OUT STD\_LOGIC; O 5: OUT STD LOGIC; O\_6 : OUT STD\_LOGIC; O 7: OUT STD LOGIC; O\_8 : OUT STD\_LOGIC; O 9: OUT STD LOGIC; O\_10 : OUT STD\_LOGIC; O 11: OUT STD LOGIC; O 12: OUT STD LOGIC; O\_13 : OUT STD\_LOGIC; O\_14 : OUT STD\_LOGIC; O 15: OUT STD LOGIC; O\_16 : OUT STD\_LOGIC; O 17: OUT STD LOGIC; O\_18 : OUT STD\_LOGIC; O 19: OUT STD LOGIC; O 20: OUT STD LOGIC; O\_21 : OUT STD\_LOGIC; O\_22 : OUT STD\_LOGIC; O\_23 : OUT STD\_LOGIC; O 24 : OUT STD LOGIC; O 25 : OUT STD LOGIC; O\_26 : OUT STD\_LOGIC; O 27 : OUT STD LOGIC; LOCK\_STATE : OUT STD\_LOGIC; C1\_STATE : OUT STD\_LOGIC; C2\_STATE : OUT STD\_LOGIC; NOM\_STATE : OUT STD\_LOGIC;

USR\_IN : IN STD\_LOGIC\_VECTOR(4 DOWNTO 1); HW\_AUTH : IN STD\_LOGIC

);

END HARDWARE\_ASSISTED\_SUPERVISOR;

```
ARCHITECTURE BEHAVIOR OF HARDWARE_ASSISTED_SUPERVISOR IS

TYPE MACHINE IS (LOCKOUT, CONTROL_1, CONTROL_2, NOMINAL);

SIGNAL STATE : MACHINE := LOCKOUT;

SIGNAL CTRL1_CNT : INTEGER RANGE 0 TO 4_000 := 0;

SIGNAL CTRL2_CNT : INTEGER RANGE 0 TO 4_000 := 0;

SIGNAL HRTBT1_LAST : STD_LOGIC := '0';

SIGNAL HRTBT1 : STD_LOGIC;

SIGNAL CTRL1_ISLIVE : BOOLEAN := FALSE;

SIGNAL HRTBT2_LAST : STD_LOGIC := '0';

SIGNAL HRTBT2 : STD_LOGIC;

SIGNAL HRTBT2 : STD_LOGIC := '0';

SIGNAL CTRL2_ISLIVE : BOOLEAN := FALSE;
```

BEGIN

```
CTRL MUX : PROCESS (CLK IN, HW AUTH, USR IN)
BEGIN
     HRTBT1 <= I1 24;
     HRTBT2 <= I2_24;
     --Update liveness timers
     IF (CLK IN'EVENT AND CLK IN = '1') THEN
           IF (HRTBT1 = NOT HRTBT1 LAST) THEN
                 CTRL1 CNT \leq 0;
                 CTRL1 ISLIVE <= TRUE;
                 HRTBT1 LAST <= HRTBT1;
           ELSE
                 IF (CTRL1 CNT < TIMEOUT) THEN
                       CTRL1_CNT \le CTRL1_CNT + 1;
                       CTRL1 ISLIVE <= TRUE;
                 ELSE
                       CTRL1_CNT <= TIMEOUT;
                       CTRL1 ISLIVE <= FALSE;
                 END IF:
           END IF;
           IF (HRTBT2 = NOT HRTBT2 LAST) THEN
                 CTRL2_CNT \ll 0;
                 CTRL2 ISLIVE <= TRUE;
```

```
HRTBT2_LAST <= HRTBT2;
      ELSE
            IF (CTRL2_CNT < TIMEOUT) THEN
                  CTRL2_CNT \leq CTRL2_CNT + 1;
                  CTRL2_ISLIVE <= TRUE;
            ELSE
                  CTRL2_CNT <= TIMEOUT;
                  CTRL2_ISLIVE <= FALSE;
            END IF;
      END IF;
END IF;
--Set states from liveness
IF (CTRL1_ISLIVE AND CTRL2_ISLIVE) THEN
      STATE <= NOMINAL;
ELSE
      IF (CTRL1_ISLIVE) THEN
            STATE <= CONTROL_1;
      ELSIF (CTRL2_ISLIVE) THEN
            STATE <= CONTROL_2;
      ELSE
            STATE <= LOCKOUT;
      END IF;
END IF;
--Set state from hardware authentication module flag
IF (HW_AUTH = '0') THEN
      STATE <= LOCKOUT;
END IF;
--Set states from user input via push buttons
IF (USR_IN(1) = '0') THEN
      STATE <= LOCKOUT;
END IF:
IF (USR_IN(2) = '0') THEN
      STATE <= NOMINAL;
END IF;
IF (USR IN(3) = 0') THEN
      STATE <= CONTROL_1;
END IF:
IF (USR_IN(4) = '0') THEN
      STATE <= CONTROL_2;
END IF;
```

--Route fabric according to set state CASE STATE IS

WHEN LOCKOUT => LOCK\_STATE <= '1'; C1\_STATE <= '0'; C2\_STATE <= '0'; NOM\_STATE  $\leq 0';$ O\_0 <= '0'; O\_1 <= '0'; O\_2 <= '0'; O\_3 <= '0'; O\_4 <= '0'; O\_5 <= '0'; O\_6 <= '0'; O\_7 <= '0'; O\_8 <= '0'; O\_9 <= '0'; O 10 <= '0'; O\_11 <= '0'; O\_12 <= '0'; O\_13 <= '0'; O\_14 <= '0'; O\_15 <= '0'; O\_16 <= '0'; O\_17 <= '0'; O 18 <= '0'; O\_19 <= '0'; O 20 <= '0'; O\_21 <= '0'; O 22 <= '0'; O\_23 <= '0'; O\_24 <= '0'; O\_25 <= '0'; O\_26 <= '0'; O\_27 <= '0'; WHEN CONTROL\_1 => LOCK\_STATE  $\leq 0';$ C1\_STATE <= '1'; C2 STATE  $\leq 0'$ ; NOM\_STATE  $\leq 0';$ O\_0 <= I1\_0; O\_1 <= I1\_1; O\_2 <= I1\_2; O\_3 <= I1\_3;  $O_4 <= I1_4;$ O\_5 <= I1\_5; O\_6 <= I1\_6;

$\begin{array}{l} O_8 <= I1\_8;\\ O_9 <= I1\_9;\\ O_10 <= I1\_10;\\ O_11 <= I1\_11;\\ O_12 <= I1\_12;\\ O_13 <= I1\_13;\\ O_14 <= I1\_14;\\ O_15 <= I1\_15;\\ O_16 <= I1\_16;\\ O_17 <= I1\_17;\\ O_18 <= I1\_18;\\ O_19 <= I1\_19;\\ O_20 <= I1\_20;\\ O_21 <= I1\_21;\\ O_22 <= I1\_22;\\ O_23 <= I1\_23;\\ O_24 <= I1\_24;\\ O_25 <= I1\_25;\\ O_26 <= I1\_26;\\ O_27 <= I1\_27;\\ \end{array}$	O_7 <= I1_7;
$O_10 \le I1_10;$ $O_11 \le I1_11;$ $O_12 \le I1_12;$ $O_13 \le I1_13;$ $O_14 \le I1_14;$ $O_15 \le I1_15;$ $O_16 \le I1_16;$ $O_17 \le I1_17;$ $O_18 \le I1_18;$ $O_19 \le I1_19;$ $O_20 \le I1_20;$ $O_21 \le I1_21;$ $O_22 \le I1_22;$ $O_23 \le I1_23;$ $O_24 \le I1_24;$ $O_25 \le I1_25;$ $O_26 \le I1_26;$	O_8 <= I1_8;
$\begin{array}{l} O_{11} <= I1\_11;\\ O_{12} <= I1\_12;\\ O_{13} <= I1\_13;\\ O_{14} <= I1\_14;\\ O_{15} <= I1\_15;\\ O_{16} <= I1\_16;\\ O_{17} <= I1\_17;\\ O_{18} <= I1\_18;\\ O_{19} <= I1\_19;\\ O_{20} <= I1\_20;\\ O_{21} <= I1\_21;\\ O_{22} <= I1\_22;\\ O_{23} <= I1\_23;\\ O_{24} <= I1\_24;\\ O_{25} <= I1\_25;\\ O_{26} <= I1\_26;\\ \end{array}$	O_9 <= I1_9;
$\begin{array}{l} O_{12} <= I1\_12;\\ O_{13} <= I1\_13;\\ O_{14} <= I1\_14;\\ O_{15} <= I1\_15;\\ O_{16} <= I1\_16;\\ O_{17} <= I1\_17;\\ O_{18} <= I1\_18;\\ O_{19} <= I1\_19;\\ O_{20} <= I1\_20;\\ O_{21} <= I1\_21;\\ O_{22} <= I1\_22;\\ O_{23} <= I1\_22;\\ O_{23} <= I1\_23;\\ O_{24} <= I1\_24;\\ O_{25} <= I1\_25;\\ O_{26} <= I1\_26;\\ \end{array}$	O_10 <= I1_10;
$\begin{array}{l} O_{-13} <= I1\_13;\\ O_{-14} <= I1\_14;\\ O_{-15} <= I1\_15;\\ O_{-16} <= I1\_16;\\ O_{-17} <= I1\_17;\\ O_{-18} <= I1\_18;\\ O_{-19} <= I1\_19;\\ O_{-20} <= I1\_20;\\ O_{-21} <= I1\_21;\\ O_{-22} <= I1\_22;\\ O_{-23} <= I1\_23;\\ O_{-24} <= I1\_24;\\ O_{-25} <= I1\_25;\\ O_{-26} <= I1\_26;\\ \end{array}$	O_11 <= I1_11;
$O_14 \le II_14;$ $O_15 \le II_15;$ $O_16 \le II_16;$ $O_17 \le II_17;$ $O_18 \le II_18;$ $O_19 \le II_19;$ $O_20 \le II_20;$ $O_21 \le II_21;$ $O_22 \le II_22;$ $O_23 \le II_23;$ $O_24 \le II_24;$ $O_25 \le II_25;$ $O_26 \le II_26;$	O_12 <= I1_12;
$\begin{array}{l} O_{15} <= I1\_15;\\ O_{16} <= I1\_16;\\ O_{17} <= I1\_17;\\ O_{18} <= I1\_18;\\ O_{19} <= I1\_19;\\ O_{20} <= I1\_20;\\ O_{21} <= I1\_21;\\ O_{22} <= I1\_22;\\ O_{23} <= I1\_23;\\ O_{24} <= I1\_24;\\ O_{25} <= I1\_25;\\ O_{26} <= I1\_26;\\ \end{array}$	O_13 <= I1_13;
$O_{-16} \le I1_{-16};$ $O_{-17} \le I1_{-17};$ $O_{-18} \le I1_{-18};$ $O_{-19} \le I1_{-19};$ $O_{-20} \le I1_{-20};$ $O_{-21} \le I1_{-21};$ $O_{-22} \le I1_{-22};$ $O_{-23} \le I1_{-23};$ $O_{-24} \le I1_{-24};$ $O_{-25} \le I1_{-25};$ $O_{-26} \le I1_{-26};$	O_14 <= I1_14;
$\begin{array}{l} O_17 <= I1\_17;\\ O_18 <= I1\_18;\\ O_19 <= I1\_19;\\ O_20 <= I1\_20;\\ O_21 <= I1\_21;\\ O_22 <= I1\_22;\\ O_23 <= I1\_23;\\ O_24 <= I1\_24;\\ O_25 <= I1\_25;\\ O_26 <= I1\_26;\\ \end{array}$	O_15 <= I1_15;
$\begin{array}{l} O_{-}18 <= I1\_18;\\ O_{-}19 <= I1\_19;\\ O_{-}20 <= I1\_20;\\ O_{-}21 <= I1\_21;\\ O_{-}22 <= I1\_22;\\ O_{-}23 <= I1\_23;\\ O_{-}24 <= I1\_24;\\ O_{-}25 <= I1\_25;\\ O_{-}26 <= I1\_26;\\ \end{array}$	O_16 <= I1_16;
$O_19 \le I1_19;$ $O_20 \le I1_20;$ $O_21 \le I1_21;$ $O_22 \le I1_22;$ $O_23 \le I1_23;$ $O_24 \le I1_24;$ $O_25 \le I1_25;$ $O_26 \le I1_26;$	O_17 <= I1_17;
$O_20 \le I1_20;$ $O_21 \le I1_21;$ $O_22 \le I1_22;$ $O_23 \le I1_23;$ $O_24 \le I1_24;$ $O_25 \le I1_25;$ $O_26 \le I1_26;$	O_18 <= I1_18;
$O_21 \le I1_21;$ $O_22 \le I1_22;$ $O_23 \le I1_23;$ $O_24 \le I1_24;$ $O_25 \le I1_25;$ $O_26 \le I1_26;$	O_19 <= I1_19;
$O_{22} \le I1_{22};$ $O_{23} \le I1_{23};$ $O_{24} \le I1_{24};$ $O_{25} \le I1_{25};$ $O_{26} \le I1_{26};$	O_20 <= I1_20;
$O_{23} \le I1_{23};$ $O_{24} \le I1_{24};$ $O_{25} \le I1_{25};$ $O_{26} \le I1_{26};$	O_21 <= I1_21;
$O_24 \le I1_24;$ $O_25 \le I1_25;$ $O_26 \le I1_26;$	O_22 <= I1_22;
$O_{25} \le I1_{25};$ $O_{26} \le I1_{26};$	O_23 <= I1_23;
$O_{26} \le I1_{26};$	$O_24 \le I1_24;$
	O_25 <= I1_25;
O_27 <= I1_27;	O_26 <= I1_26;
	O_27 <= I1_27;

WHEN CONTROL\_2 =>

LOCK\_STATE  $\leq 0';$ C1\_STATE <= '0'; C2\_STATE <= '1'; NOM\_STATE  $\leq 0';$  $O_0 \le I2_0;$ O\_1 <= I2\_1; O\_2 <= I2\_2; O\_3 <= I2\_3; O\_4 <= I2\_4; O\_5 <= I2\_5; O\_6 <= I2\_6; O\_7 <= I2\_7; O\_8 <= I2\_8; O\_9 <= I2\_9; O\_10 <= I2\_10; O\_11 <= I2\_11; O\_12 <= I2\_12; O\_13 <= I2\_13; O\_14 <= I2\_14; O\_15 <= I2\_15; O\_16 <= I2\_16; O\_17 <= I2\_17; O\_18 <= I2\_18;

O_19 <= I2_19;
O_20 <= I2_20;
O_21 <= I2_21;
O_22 <= I2_22;
O_23 <= I2_23;
O_24 <= I2_24;
O_25 <= I2_25;
O_26 <= I2_26;
O_27 <= I2_27;

WHEN NOMINAL => LOCK\_STATE <= '0'; C1\_STATE <= '0'; C2\_STATE <= '0'; NOM\_STATE  $\leq 1'$ ; O\_0 <= I1\_0; O\_1 <= I1\_1; O\_2 <= I1\_2; O\_3 <= I1\_3; O\_4 <= I1\_4; O\_5 <= I1\_5; O\_6 <= I1\_6; O\_7 <= I1\_7; O\_8 <= I1\_8; O\_9 <= I1\_9; O\_10 <= I1\_10; O\_11 <= I1\_11; O\_12 <= I1\_12; O\_13 <= I1\_13;  $O_14 \le I1_14;$ O\_15 <= I1\_15; O\_16 <= I1\_16; O\_17 <= I1\_17; O\_18 <= I1\_18; O\_19 <= I1\_19; O\_20 <= I1\_20; O\_21 <= I1\_21; O  $22 \le I1 22;$ O\_23 <= I1\_23; O\_24 <= I1\_24; O\_25 <= I1\_25; O\_26 <= I1\_26; O\_27 <= I1\_27;

WHEN OTHERS => LOCK\_STATE <= '0';

C1_STATE <= '0';
C2 STATE <= '0';
$NOM\_STATE <= '0';$
O_0 <= '0';
O_1 <= '0';
O_2 <= '0';
O_3 <= '0';
O_4 <= '0';
O_5 <= '0';
O_6 <= '0';
O_7 <= '0';
O_8 <= '0';
O_9 <= '0';
O_10 <= '0';
O_11 <= '0';
O_12 <= '0';
O_13 <= '0';
O_14 <= '0';
O_15 <= '0';
O_16 <= '0';
O_17 <= '0';
O_18 <= '0';
O_19 <= '0';
O_20 <= '0';
O_21 <= '0';
O_22 <= '0';
O_23 <= '0';
O_24 <= '0';
O_25 <= '0';
O_26 <= '0';
O_27 <= '0';

END CASE; END PROCESS CTRL\_MUX; END BEHAVIOR;

## Appendix F: top.vhdl

LIBRARY IEEE; USE IEEE.STD\_LOGIC\_1164.ALL; LIBRARY LATTICE; USE LATTICE.COMPONENTS.ALL; LIBRARY WORK; USE WORK.CSPR\_MODULES.ALL;

#### ENTITY CSPR IS

PORT (

--Question: Why aren't we just using STD\_LOGIC\_VECTORs for the IDC ports? --Answer: These IDC ports have both input and output pins. It is much more clear (but verbose) to handle them individually.

--IDC A A0: IN STD LOGIC; A1 : IN STD\_LOGIC; A2 : IN STD\_LOGIC; A3 : IN STD\_LOGIC; A4 : IN STD\_LOGIC; A5 : IN STD LOGIC; A6: IN STD LOGIC; A7 : IN STD\_LOGIC; --A8 IN STD LOGIC; : --A9 IN STD\_LOGIC; : A10: IN STD LOGIC; A11 : IN STD\_LOGIC; A12: IN STD LOGIC; A13 : IN STD\_LOGIC; A14 : IN STD LOGIC; A15 : IN STD LOGIC; A16 : IN STD\_LOGIC; A17: IN STD LOGIC; A18: IN STD LOGIC; A19 : IN STD\_LOGIC; --A20 IN STD LOGIC; : --A21 : IN STD\_LOGIC; IN STD LOGIC; --A22 : --A23 IN STD LOGIC; : A24 : IN STD\_LOGIC; A25 : IN STD\_LOGIC; A26 : IN STD\_LOGIC; A27 : IN STD LOGIC; --IDC B B0 : OUT STD\_LOGIC := 'Z'; B1 : OUT STD LOGIC := 'Z';

B2 : OUT STD_LOGIC := 'Z'; B3 : OUT STD_LOGIC := 'Z'; B4 : OUT STD_LOGIC := 'Z'; B5 : OUT STD_LOGIC := 'Z'; B6 : OUT STD_LOGIC := 'Z'; B7 : OUT STD_LOGIC := 'Z'; B8 : OUT STD_LOGIC := 'Z'; B9 : OUT STD_LOGIC := 'Z'; B10 : OUT STD_LOGIC := 'Z';	
B11 : OUT STD_LOGIC := 'Z'; B12 : OUT STD_LOGIC := 'Z';	
B13 : OUT STD_LOGIC := 'Z';	
$B14 : OUT STD_LOGIC := 'Z';$	
B15 : OUT STD_LOGIC := 'Z';	
B16 : OUT STD_LOGIC := 'Z'; P17 : OUT STD_LOGIC := 'Z';	
B17 : OUT STD_LOGIC := 'Z'; B18 : OUT STD_LOGIC := 'Z';	
$B19:OUT STD_LOGIC := 'Z';$	
$B20: OUT STD\_LOGIC := 'Z';$	
B21 : OUT STD_LOGIC := 'Z';	
B22 :	OUT STD_LOGIC := 'Z';
-B23 :	OUT STD_LOGIC := 'Z';
B24 : OUT STD_LOGIC := 'Z'; B25 : OUT STD_LOGIC := 'Z';	
B26 : OUT STD_LOGIC := 'Z';	
$B27 : OUT STD_LOGIC := 'Z';$	
IDC C	
C0 : IN STD_LOGIC;	
C1 : IN STD_LOGIC;	
C2 : IN STD_LOGIC;	
C3 : IN STD_LOGIC; C4 : IN STD_LOGIC;	
$C_{4}$ : IN STD_LOGIC;	
$C6 : IN STD_LOGIC;$	
C7 : IN STD_LOGIC;	
C8 :	IN STD_LOGIC;
C9 :	IN STD_LOGIC;
C10 : IN STD_LOGIC; C11 : IN STD_LOGIC;	
$C12 : IN STD_LOGIC;$	
$C13 : IN STD_LOGIC;$	
C14 : IN STD_LOGIC;	
C15 : IN STD_LOGIC;	
C16 : IN STD_LOGIC;	
C17 : IN STD_LOGIC;	
C18 : IN STD_LOGIC;	

C19 : IN STD\_LOGIC; --C20 IN STD\_LOGIC; : --C21 IN STD\_LOGIC; : --C22 IN STD LOGIC; : --C23 IN STD\_LOGIC; : C24 : IN STD\_LOGIC; C25 : IN STD\_LOGIC; C26 : IN STD\_LOGIC; C27 : IN STD LOGIC; --IDC D D0: OUT STD\_LOGIC; D1 : OUT STD\_LOGIC; D2 : OUT STD\_LOGIC; D3 : OUT STD\_LOGIC; D4 : OUT STD\_LOGIC; D5 : OUT STD LOGIC; D6 : OUT STD\_LOGIC; D7 : OUT STD\_LOGIC; D8 : OUT STD\_LOGIC; D9 : OUT STD\_LOGIC; D10: OUT STD LOGIC; D11: OUT STD LOGIC; D12 : OUT STD\_LOGIC; D13: OUT STD LOGIC; D14 : OUT STD\_LOGIC; D15: OUT STD LOGIC; D16 : OUT STD\_LOGIC; D17: OUT STD LOGIC; D18 : OUT STD\_LOGIC; D19: OUT STD LOGIC; D20 : OUT STD\_LOGIC; D21 : OUT STD\_LOGIC; D22 : OUT STD\_LOGIC; D23 : IN STD LOGIC; -- EEPROM Master In Slave Out D24 : OUT STD\_LOGIC; D25 : OUT STD LOGIC; D26 : OUT STD\_LOGIC; D27: OUT STD LOGIC; Usr\_RX : IN STD\_LOGIC; Usr\_TX : OUT STD\_LOGIC;

BTN : IN STD\_LOGIC\_VECTOR(4 DOWNTO 1); LED : OUT STD\_LOGIC\_VECTOR(8 DOWNTO 1)

); END CSPR; ARCHITECTURE BEHAVIOR OF CSPR IS COMPONENT OSCH GENERIC (NOM\_FREQ : STRING := "53.2"); PORT ( STDBY : IN STD\_LOGIC; OSC : OUT STD\_LOGIC; SEDSTDBY : OUT STD\_LOGIC ); END COMPONENT; COMPONENT PLL PORT ( CLKI : IN STD\_LOGIC; CLKOP: OUT STD\_LOGIC; CLKOS : OUT STD LOGIC; CLKOS2 : OUT STD LOGIC; CLKOS3 : OUT STD\_LOGIC; LOCK : OUT STD\_LOGIC ); END COMPONENT; COMPONENT HARDWARE\_ASSISTED\_SUPERVISOR PORT ( CLK\_IN : IN STD\_LOGIC; --Input Controller 1 I1\_0 : IN STD\_LOGIC; I1 1: IN STD LOGIC; I1\_2 : IN STD\_LOGIC; I1 3 : IN STD LOGIC: I1 4 : IN STD LOGIC; I1\_5 : IN STD\_LOGIC; I1 6: IN STD LOGIC; I1 7: IN STD LOGIC; I1\_8 : IN STD\_LOGIC; I1 9: IN STD LOGIC; I1\_10 : IN STD\_LOGIC; I1\_11 : IN STD\_LOGIC; I1 12 : IN STD LOGIC; I1\_13 : IN STD\_LOGIC; I1\_14 : IN STD\_LOGIC; I1\_15 : IN STD\_LOGIC; I1 16 : IN STD LOGIC; I1 17 : IN STD LOGIC; I1\_18 : IN STD\_LOGIC; I1 19: IN STD LOGIC;

I1\_20 : IN STD\_LOGIC; I1\_21 : IN STD\_LOGIC; I1\_22 : IN STD\_LOGIC; I1\_23 : IN STD\_LOGIC; I1\_24 : IN STD\_LOGIC; I1\_25 : IN STD\_LOGIC; I1\_26 : IN STD\_LOGIC; I1\_27 : IN STD\_LOGIC; --Input Controller 2 I2\_0 : IN STD\_LOGIC; I2 1: IN STD LOGIC; I2\_2 : IN STD\_LOGIC; I2 3 : IN STD\_LOGIC; I2\_4 : IN STD\_LOGIC; I2\_5 : IN STD\_LOGIC; I2 6: IN STD LOGIC; I2\_7 : IN STD\_LOGIC; I2\_8 : IN STD\_LOGIC; I2\_9 : IN STD\_LOGIC; I2\_10 : IN STD\_LOGIC; I2 11 : IN STD LOGIC; I2\_12 : IN STD\_LOGIC; I2\_13 : IN STD\_LOGIC; I2 14 : IN STD LOGIC; I2\_15 : IN STD\_LOGIC; I2 16 : IN STD LOGIC; I2\_17 : IN STD\_LOGIC; I2 18 : IN STD LOGIC; I2\_19 : IN STD\_LOGIC; I2 20 : IN STD LOGIC; I2\_21 : IN STD\_LOGIC; I2\_22 : IN STD\_LOGIC; I2\_23 : IN STD\_LOGIC; I2 24 : IN STD LOGIC; I2\_25 : IN STD\_LOGIC; I2 26 : IN STD LOGIC; I2\_27 : IN STD\_LOGIC; --Output Controller O 0: OUT STD LOGIC;O\_1 : OUT STD\_LOGIC; O\_2 : OUT STD\_LOGIC; O\_3 : OUT STD\_LOGIC; O 4: OUT STD LOGIC; O 5: OUT STD LOGIC; O\_6 : OUT STD\_LOGIC; O 7: OUT STD LOGIC; O 8: OUT STD LOGIC; O\_9: OUT STD\_LOGIC; O\_10: OUT STD\_LOGIC; O 11: OUT STD LOGIC; O 12: OUT STD LOGIC: O\_13 : OUT STD\_LOGIC; O\_14 : OUT STD\_LOGIC; O\_15 : OUT STD\_LOGIC; O\_16: OUT STD\_LOGIC; O\_17 : OUT STD\_LOGIC; O 18: OUT STD LOGIC; O\_19: OUT STD\_LOGIC; O 20: OUT STD LOGIC: O\_21 : OUT STD\_LOGIC; O\_22 : OUT STD\_LOGIC; O 23: OUT STD LOGIC; O 24 : OUT STD LOGIC; O\_25 : OUT STD\_LOGIC; O 26: OUT STD LOGIC; O\_27: OUT STD\_LOGIC;

LOCK\_STATE : OUT STD\_LOGIC; C1\_STATE : OUT STD\_LOGIC; C2\_STATE : OUT STD\_LOGIC; NOM\_STATE : OUT STD\_LOGIC;

USR\_IN : IN STD\_LOGIC\_VECTOR(4 DOWNTO 1); HW\_AUTH : IN STD\_LOGIC

);

END COMPONENT;

COMPONENT HW\_AUTH\_MODULE

PORT (

MISO : IN STD\_LOGIC; CLK\_IN : IN STD\_LOGIC; CS : OUT STD\_LOGIC; CLK\_OUT : OUT STD\_LOGIC; MOSI : OUT STD\_LOGIC; HW\_GOOD : OUT STD\_LOGIC

END COMPONENT;

--COMPONENT Bus\_Master --PORT ( --clk : IN std\_logic; --rst : IN std\_logic; --Data : INOUT std\_logic\_vector(15 DOWNTO 0);

- --Addr : IN std\_logic\_vector(15 DOWNTO 0);
- --Xrqst : IN std\_logic;
- --XDat : OUT std\_logic;
- --YDat : IN std\_logic;
- --BusRqst : IN std\_logic\_vector(9 DOWNTO 0);
- --BusCtrl : OUT std\_logic\_vector(9 DOWNTO 0)

--);

- --END COMPONENT;
- --COMPONENT RS232\_Usr\_Int
- --GENERIC (
- --Baud : INTEGER; -- Baud Rate
- --clk\_in : INTEGER -- Input Clk
- --);
- --PORT (
- --clk : IN std\_logic;
- --rst : IN std\_logic;
- --rs232\_rcv : IN std\_logic;
- --rs232\_xmt : OUT std\_logic;
- --Data : INOUT std\_logic\_vector(15 DOWNTO 0);
- --Addr : OUT std\_logic\_vector(15 DOWNTO 0);
- --Xrqst : OUT std\_logic;
- --XDat : IN std\_logic;
- --YDat : OUT std\_logic;
- --BusRqst : OUT std\_logic;
- --BusCtrl : IN std\_logic
- --);
- --END COMPONENT;

--COMPONENT LED\_Ctrl

- --PORT (
- --clk : IN STD\_LOGIC;
- --rst : IN STD\_LOGIC;
- --Data : INOUT std\_logic\_vector(15 DOWNTO 0);
- --Addr : OUT std\_logic\_vector(15 DOWNTO 0);
- --Xrqst : OUT std\_logic;
- --XDat : IN std\_logic;
- --YDat : OUT std\_logic;
- --BusRqst : OUT std\_logic;
- --BusCtrl : IN std\_logic;
- --LED1\_Out : OUT STD\_LOGIC

--);

--END COMPONENT;

--COMPONENT Std\_Counter

--GENERIC ( --Width : INTEGER -- width of counter --): --PORT ( --INC, rst, clk : IN std\_logic; : OUT STD\_LOGIC\_VECTOR(Width - 1 DOWNTO 0)); --Count --END COMPONENT; SIGNAL OSC Stdby : STD LOGIC := '0'; SIGNAL PLL\_IN : STD\_LOGIC := '0'; SIGNAL OSC SEDSTDBY : STD LOGIC := '0'; SIGNAL CLK53 2M : STD LOGIC := '0'; SIGNAL CLK24\_93M : STD LOGIC := '0': SIGNAL CLK8\_31M : STD\_LOGIC := '0'; SIGNAL CLK1 5M : STD LOGIC := '0'; SIGNAL PLL\_LOCK : STD\_LOGIC := '0'; SIGNAL MISO0 : STD LOGIC := '0'; SIGNAL CS0 : STD\_LOGIC := '0'; SIGNAL CLK OUT0 : STD LOGIC := '0'; SIGNAL MOSIO : STD LOGIC := '0'; SIGNAL HW\_AUTH\_FLAG : STD\_LOGIC := '0'; SIGNAL A : STD\_LOGIC\_VECTOR(27 DOWNTO 0) := (OTHERS => '0'); SIGNAL B : STD LOGIC VECTOR(27 DOWNTO 0) := (OTHERS => '0'); SIGNAL C : STD\_LOGIC\_VECTOR(27 DOWNTO 0) := (OTHERS => '0'); SIGNAL D : STD LOGIC VECTOR(21 DOWNTO 0) := (OTHERS => 0'); SIGNAL LOCK S : STD LOGIC := '0': SIGNAL C1 S : STD LOGIC := '0'; SIGNAL C2\_S : STD\_LOGIC := '0'; SIGNAL NOM\_S : STD\_LOGIC := '0'; SIGNAL LED\_OUT : STD\_LOGIC\_VECTOR(8 DOWNTO 1); --From Bus Interface Top SIGNAL Bus\_Int1\_WE, Bus\_Int1\_RE, Bus\_Int1\_Busy : STD\_LOGIC := '0'; SIGNAL Bus\_Int1\_DataIn, Bus\_Int1\_DataOut, Bus Int1 AddrIn : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Addr : STD\_LOGIC\_VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL Xrqst : STD\_LOGIC := '0'; SIGNAL YDat : STD LOGIC := '0': SIGNAL BusRqst : STD LOGIC VECTOR(9 DOWNTO 0) := (OTHERS => '0'); SIGNAL Data : STD LOGIC VECTOR(15 DOWNTO 0) := (OTHERS => '0'); SIGNAL XDat : STD\_LOGIC := '0'; SIGNAL BusCtrl : STD LOGIC VECTOR(9 DOWNTO 0) := (OTHERS => '0');

```
SIGNAL System_rst : STD_LOGIC := '0';
SIGNAL Reset_Cnt_INC, Reset_Cnt_rst : STD_LOGIC := '0';
SIGNAL Reset_Cnt_out : STD_LOGIC_VECTOR(7 DOWNTO 0) := (OTHERS => '0');
SIGNAL LED_1n : STD_LOGIC := '0';
```

# BEGIN

```
Int_OSC : OSCH
PORT MAP(
STDBY => OSC_Stdby,
OSC => PLL_IN,
SEDSTDBY => OSC_SEDSTDBY
```

```
);
```

```
PLL0 : PLL
PORT MAP(
CLKI => PLL_IN,
CLKOP => CLK53_2M,
CLKOS => CLK24_93M,
CLKOS2 => CLK8_31M,
CLKOS3 => CLK1_5M,
LOCK => PLL_LOCK
```

);

# CTRL\_MUX : HARDWARE\_ASSISTED\_SUPERVISOR PORT MAP(

```
CLK_IN => CLK53_2M,
I1 0 \Rightarrow A(0),
I1_1 => A(1),
I1_2 => A(2),
I1_3 => A(3),
I1_4 => A(4),
I1_5 => A(5),
I1_6 => A(6),
I1_7 => A(7),
I1_8 => '0', --A(8),
I1_9 => '0', --A(9),
I1 10 => A(10),
I1_1 => A(11),
I1_12 => A(12),
I1_13 => A(13),
I1_14 => A(14),
I1_{15} \Rightarrow A(15),
I1_16 => A(16),
I1_17 => A(17),
I1_{18} => A(18),
```

$$\begin{split} & \text{I1}_{19} => A(19), \\ & \text{I1}_{20} => '0', --A(20), \\ & \text{I1}_{21} => '0', --A(21), \\ & \text{I1}_{22} => '0', --A(23), \\ & \text{I1}_{23} => '0', --A(23), \\ & \text{I1}_{24} => A(24), \\ & \text{I1}_{25} => A(25), \\ & \text{I1}_{26} => A(26), \\ & \text{I1}_{27} => A(27), \\ & \text{I2}_{0} => C(0), \\ & \text{I2}_{1} => C(1), \\ & \text{I2}_{2} => C(2), \\ & \text{I2}_{3} => C(3), \\ & \text{I2}_{4} => C(4), \\ & \text{I2}_{5} => C(5), \\ & \text{I2}_{6} => C(6), \\ & \text{I2}_{7} => C(7), \\ & \text{I2}_{8} => '0', --C(8), \\ & \text{I2}_{9} => '0', --C(8), \\ & \text{I2}_{9} => '0', --C(9), \\ & \text{I2}_{10} => C(10), \\ & \text{I2}_{11} => C(11), \\ & \text{I2}_{12} => C(12), \\ & \text{I2}_{13} => C(13), \\ & \text{I2}_{14} => C(14), \\ & \text{I2}_{15} => C(15), \\ & \text{I2}_{16} => C(16), \\ & \text{I2}_{17} => C(17), \\ & \text{I2}_{18} => C(18), \\ & \text{I2}_{19} => C(19), \\ & \text{I2}_{20} => '0', --C(20), \\ & \text{I2}_{21} => '0', --C(21), \\ & \text{I2}_{22} => C(26), \\ & \text{I2}_{27} => C(27), \\ & \text{O}_{0} => D(0), \\ & \text{O}_{1} => D(1), \\ & \text{O}_{2} => D(2), \\ & \text{O}_{3} => D(3), \\ & \text{O}_{4} => D(4), \\ & \text{O}_{5} => D(5), \\ & \text{O}_{6} => D(6), \\ & \text{O}_{7} => D(7), \\ & \text{O}_{8} => OPEN, \\ \\ \end{split}$$

 $O_9 => OPEN,$  $O_{10} => D(8),$  $O_11 => D(9),$  $O_{12} => D(10),$  $O_13 => D(11),$  $O_14 => D(12),$  $O_15 => D(13),$  $O_16 => D(14),$  $O_17 => D(15),$  $O_18 => D(16),$  $O_19 => D(17),$  $O_20 \Rightarrow OPEN$ ,  $O_21 \Rightarrow OPEN,$  $O_22 \Rightarrow OPEN,$  $O_23 \Rightarrow OPEN,$ O  $24 \Rightarrow D(18)$ ,  $O_25 => D(19),$  $O_26 => D(20),$  $O_27 => D(21),$ LOCK\_STATE => LOCK\_S,  $C1\_STATE => C1\_S$ ,  $C2\_STATE => C2\_S$ , NOM\_STATE => NOM\_S, USR IN => BTN, HW\_AUTH => HW\_AUTH\_FLAG

);

);

--BM : Bus\_Master --PORT MAP( --clk => CLK24\_93M, --rst => System\_rst, --Data => Data, --Addr => Addr, --Xrqst => Xrqst, --XDat => XDat, --YDat => YDat,

```
--BusRqst => BusRqst,
--BusCtrl => BusCtrl
--);
--RS232_Usr : RS232_Usr_Int
--GENERIC MAP(
--Baud => 9600, -- Baud Rate
--Clk_In => 24937500 --Clk_Freq
--)
--PORT MAP(
--clk
       => CLK24_93M,
       => System_rst,
--rst
-rs232_rcv => Usr_RX,
--rs232_xmt => Usr_TX,
--Data
       => Data,
--Addr
        => Addr,
--Xrqst => Xrqst,
--XDat
        => XDat,
--YDat
        \Rightarrow YDat,
--BusRqst => BusRqst(1),
--BusCtrl => BusCtrl(1)
--);
--LED_Ctrl1 : LED_Ctrl
--PORT MAP(
--clk
       => CLK24 93M,
--rst
      => System_rst,
--Data => Data,
--Addr => Addr,
--Xrqst => Xrqst,
--XDat => XDat,
--YDat => YDat,
--BusRqst => BusRqst(0),
--BusCtrl => BusCtrl(0),
--LED1_Out => LED_1n
--);
--Reset Cnt: Std Counter
--GENERIC MAP
--(
--Width => 8
--)
--PORT MAP(
--clk => CLK8_31M,
--rst => Reset_Cnt_rst,
--INC => Reset_Cnt_INC,
```

--Count => Reset\_Cnt\_Out --); --OSC\_Stdby <= '0'; --BusRqst(9 DOWNTO 2) <= (OTHERS => '0'); -- IDC A, INPUT from CONTROLLER 1  $A(0) \le A0;$  $A(1) \le A1;$  $A(2) \le A2;$  $A(3) \le A3;$ A(4) <= A4; $A(5) \le A5;$  $A(6) \le A6;$  $A(7) \le A7;$ A(8) <= '0';--A8; A(9) <= '0';--A9; A(10) <= A10; A(11) <= A11;  $A(12) \le A12;$ A(13) <= A13;  $A(14) \le A14;$ A(15) <= A15; A(16) <= A16; A(17) <= A17;  $A(18) \le A18;$ A(19) <= A19; A(20) <= '0';--A20; A(21) <= '0';--A21; A(22) <= '0';--A22; A(23) <= '0';--A23;  $A(24) \le A24;$ A(25) <= A25; A(26) <= A26; A(27) <= A27; -- IDC C, INPUT from CONTROLLER 2  $C(0) \le C0;$  $C(1) \le C1;$  $C(2) \le C2;$ 

C(3) <= C3;

 $C(4) \le C4;$ 

C(5) <= C5; C(6) <= C6;

C(7) <= C7;

C(8) <= '0';--C8;

C(9) <= '0';C9;
C(10) <= C10;
C(11) <= C11;
C(12) <= C12;
C(13) <= C13;
C(14) <= C14;
C(15) <= C15;
C(16) <= C16;
C(17) <= C17;
C(18) <= C18;
C(19) <= C19;
C(20) <= '0';C20;
C(21) <= '0';C21;
C(22) <= '0';C22;
C(23) <= '0';C23;
C(24) <= C24;
C(25) <= C25;
C(26) <= C26;
C(27) <= C27;

--IDC D, OUTPUT from CONTROLLER 1 or CONTROLLER 2 to POWER ELECTRONICS

D0 <= D(0);  $D1 \le D(1);$  $D2 \le D(2);$  $D3 \le D(3);$  $D4 \le D(4);$ D5 <= D(5);  $D6 \le D(6);$  $D7 \le D(7);$ D8 <= '1'; D9 <= NOT HW\_AUTH\_FLAG; D10 <= D(8); D11 <= D(9); D12 <= D(10); D13 <= D(11); D14 <= D(12);  $D15 \le D(13);$  $D16 \le D(14);$ D17 <= D(15); D18 <= D(16); D19 <= D(17);  $D20 \leq CLK_OUT0;$ D21 <= CS0; D22 <= MOSI0; MISO0 <= D23;

 $D24 \le D(18);$  $D25 \le D(19);$  $D26 \le D(20);$  $D27 \le D(21);$ 

```
\begin{split} LED(1) &= HW\_AUTH\_FLAG;\\ LED(2) &= NOT(HW\_AUTH\_FLAG);\\ LED(3) &= NOT(LED\_1n);\\ LED(4) &= '1';\\ LED(5) &= NOT(LOCK\_S);\\ LED(6) &= NOT(NOM\_S);\\ LED(7) &= NOT(C1\_S);\\ LED(8) &= NOT(C2\_S); \end{split}
```

```
--Reset_Blk1 : PROCESS

--BEGIN

--WAIT UNTIL CLK8_31M'event AND CLK8_31M = '1';

--IF (PLL_Lock = '0') THEN

--Reset_Cnt_rst <= '0';

--ELSE

--Reset_Cnt_rst <= '1';

--END IF;

--END PROCESS;

--Reset_Blk : PROCESS

PECINI
```

```
--BEGIN

--WAIT UNTIL CLK8_31M'event AND CLK8_31M = '1';

--IF (Reset_Cnt_out < X"7F") THEN

--System_rst <= '0';

--Reset_Cnt_Inc <= '1';

--ELSE

--System_rst <= '1';

--Reset_Cnt_Inc <= '0';

--END IF;

--END PROCESS;
```

# END BEHAVIOR;